

# Homework 1-25

ωεβ:=<http://drorbn.net/ld26>



**Strong.**  $\Theta$  vs. a slew of other reasonably-computable invariants (deficits shown):

**Abstract.** I'll start with a review of my recent paper with van der Veen, "A Fast, Strong, Topologically Meaningful, and Fun Knot Invariant" [BV3], and then assign some homework. Much of what I'll say follows earlier work by Rozansky, Kriker, Garoufalidis, and Ohtsuki [Ro1, Ro2, Ro4, Kr, GR, Oh2].



van der Veen

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**A.** With  $T$  an indeterminate, start from a presentation matrix  $A$  for the Alexander module of  $K$ , coming from the Wirtinger presentation of  $\pi_1(K)$ :  $A := I_{2n+1} + \sum_c A_c$ , where

$$A_c = \begin{array}{c|cc} & i+1 & j+1 \\ \hline i & -T^s & T^s-1 \\ j & 0 & -1 \end{array}$$

$$A = \begin{pmatrix} 1 & -T & 0 & 0 & T-1 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & -T & 0 & 0 & T-1 \\ 0 & 0 & 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & T-1 & 0 & 1 & -T & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$\Delta \doteq \det(A)$

**G.** Let  $G = (g_{\alpha\beta}) := A^{-1}$ :

$$G = \begin{pmatrix} 1 & T & 1 & T & 1 & T & 1 \\ 0 & 1 & \frac{1}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T^2}{T^2-T+1} & 1 \\ 0 & 0 & \frac{1}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T^2}{T^2-T+1} & 1 \\ 0 & 0 & \frac{1-T}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T}{T^2-T+1} & \frac{T^2}{T^2-T+1} & 1 \\ 0 & 0 & \frac{1-T}{T^2-T+1} & -\frac{(T-1)T}{T^2-T+1} & \frac{1}{T^2-T+1} & \frac{T}{T^2-T+1} & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$



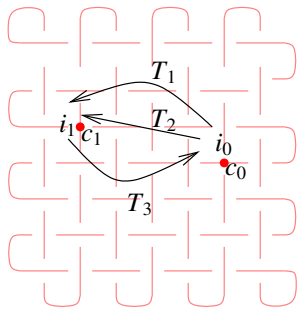
ωεβ/pi

Let  $T_1$  and  $T_2$  be new indeterminates, let  $T_3 = T_1 T_2$ , and let  $G_\nu = (g_{\nu\alpha\beta})$  be  $G$  with  $T \rightarrow T_\nu$ , for  $\nu = 1, 2, 3$ .

$$\theta \sim \Delta_1 \Delta_2 \Delta_3 \sum_{c_0, c_1} g_{1i_0 i_1} g_{2i_0 i_1} g_{3i_1 i_0} + \text{l.o.}$$

$$\Theta = (\Delta, \theta) \in \mathbb{Z}[T^{\pm 1}] \times \mathbb{Z}[T_1^{\pm 1}, T_2^{\pm 1}]$$

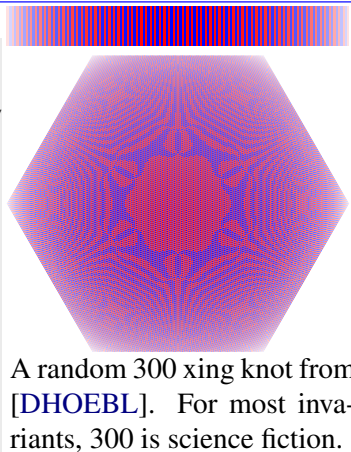
$$\begin{pmatrix} \frac{2}{T} & -1 & 3T \\ T_2 & -T_1 T_2 & \\ -\frac{1}{T_1} & 2 & T_1 \\ \frac{1}{T_1 T_2} & -\frac{1}{T_2} & \end{pmatrix} \rightarrow \begin{pmatrix} \frac{2}{T} & -1 & 3T \\ T_2 & -T_1 T_2 & \\ -\frac{1}{T_1} & 2 & T_1 \\ \frac{1}{T_1 T_2} & -\frac{1}{T_2} & \end{pmatrix}$$



**Fast.**

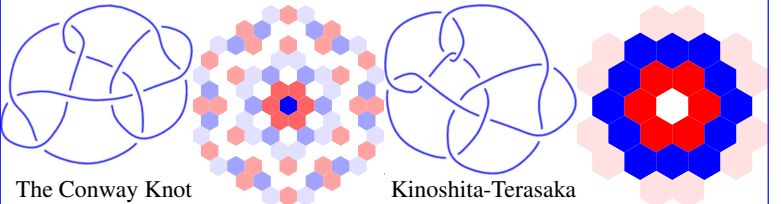
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F1([s_-, e_-, j_-]) := CF
S := (1/2 - B11 + T1 B21 - B21 B31 - (T1-1) B21 B31 + 2 B21 B31 -
(1-T1) B21 B31 - B21 B31 - T1 B21 B31 + B21 B31 +
((T1-1) B21 (T1-1) B21 - T1 B21 + T1 B31) +
(T1-1) B21 (1-T1 B21 + B21 + (T1-2) B21 - (T1-1) (T1+1) B31)/
(T1-1))
F2([s0, i0, j0], [s1, i1, j1]) :=
CF[s1 (T1-1) (T1-1) (T1-1) (T1-1) B21, i0, s0, j0, i1]
((T1-1) B21, i0 - B21, i0) - (T1-1) B21, j0 - B21, j0))
F3([u_-, h_-]) = 1/2 B31 - 1/2;
F4 := T1 T2;
CF[0] := ExpandCollect[0, g_-, F] /. F -> Factor;
Program
theta := 0;
Module[{X, phi, n, A, Delta, G, ev, phi, k1, k2},
{X, phi} = Rot[K]; n = Length[X]; A = IdentityMatrix[2 n + 1];
Cases[X, {s_-, e_-, j_-} >> {A[[{i, j}, {i+1, j+1}]] += {
-T^s T^j-1,
0, -1}}];
Delta = (-Total[phi]-Total[X[[All, 1]]])/2 Det[A];
G = Inverse[A];
ev[phi_] := Factor[phi /. {B21, B31} -> {G[[phi, phi]] /. T -> T1}];
ev[phi_] := Sum[F2[X[[phi]], X[[phi]], {k1, n}, {k2, n}];
ev[phi_] := Sum[F2[X[[phi]], X[[phi]], {k1, n}, {k2, n}];
Factor[phi, {Delta, T -> T1} (Delta, T -> T2) (Delta, T -> T3) phi];
];
    
```



A random 300 xing knot from [DHOEBL]. For most invariants, 300 is science fiction.

$n$	$\leq 10$	$\leq 11$	$\leq 12$	$\leq 13$	$\leq 14$	$\leq 15$
knots	249	801	2,977	12,965	59,937	313,230
$\Delta$	(38)	(250)	(1,204)	(7,326)	(39,741)	(236,326)
$\sigma_{LT}$	(108)	(356)	(1,525)	(7,736)	(40,101)	(230,592)
$J$	(7)	(70)	(482)	(3,434)	(21,250)	(138,591)
$Kh$	(6)	(65)	(452)	(3,226)	(19,754)	(127,261)
$H$	(2)	(31)	(222)	(1,839)	(11,251)	(73,892)
$Vol$	(~6)	(~25)	(~113)	(~1,012)	(~6,353)	(~43,607)
$(Kh, H, Vol)$	(~0)	(~14)	(~84)	(~911)	(~5,917)	(~41,434)
$(\Delta, \rho_1)$	(0)	(14)	(95)	(959)	(6,253)	(42,914)
$(\Delta, \rho_1, \rho_2)$	(0)	(14)	(84)	(911)	(5,926)	(41,469)
$(\rho_1, \rho_2, Kh, H, Vol)$	(0)	(~14)	(~84)	(~911)	(~5,916)	(~41,432)
<b><math>\Theta</math></b>	<b>(0)</b>	<b>(3)</b>	<b>(19)</b>	<b>(194)</b>	<b>(1,118)</b>	<b>(6,758)</b>
$(\Theta, \rho_2)$	(0)	(3)	(10)	(169)	(982)	(6,341)
$(\Theta, \sigma_{LT})$	(0)	(3)	(19)	(194)	(1,118)	(6,758)
$(\Theta, Kh)$	(0)	(3)	(18)	(185)	(1,062)	(6,555)
$(\Theta, H)$	(0)	(3)	(18)	(185)	(1,064)	(6,563)
$(\Theta, Vol)$	(0)	(~3)	(~10)	(~169)	(~973)	(~6,308)
$(\Theta, \rho_2, Kh, H, Vol)$	(0)	(~3)	(~10)	(~169)	(~972)	(~6,304)



**Topologically Meaningful.**  $\theta$  is near  $\Delta$  and we dream that anything  $\Delta$  can do,  $\theta$  does too (sometimes better). The following two conjectures are verified for knots with  $\leq 13$  crossings:

**Conjecture 1.**  $\deg_{T_1} \theta(K) \leq 2g(K)$ .  
**Conjecture 2.** If  $K$  is a fibered knot and  $d$  is the degree of  $\Delta(K)$  (the highest power of  $T$ ), then the coefficient of  $T_1^{2d}$  in  $\theta(K)$ , which is a polynomial in  $T_1$ , is an integer multiple of  $T_1^{2d} \Delta(K)|_{T \rightarrow T_1}$ .  
**Dream.**  $\theta$  has something to say about ribbon knots.

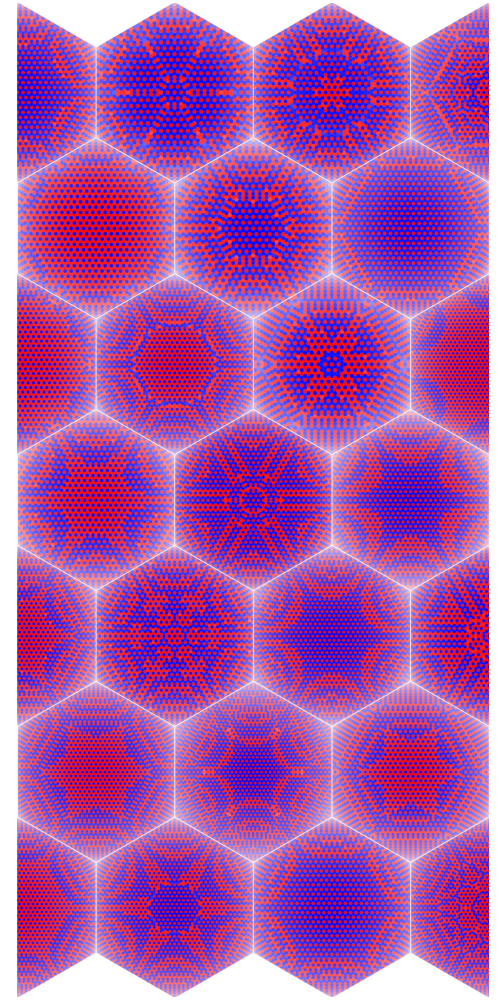
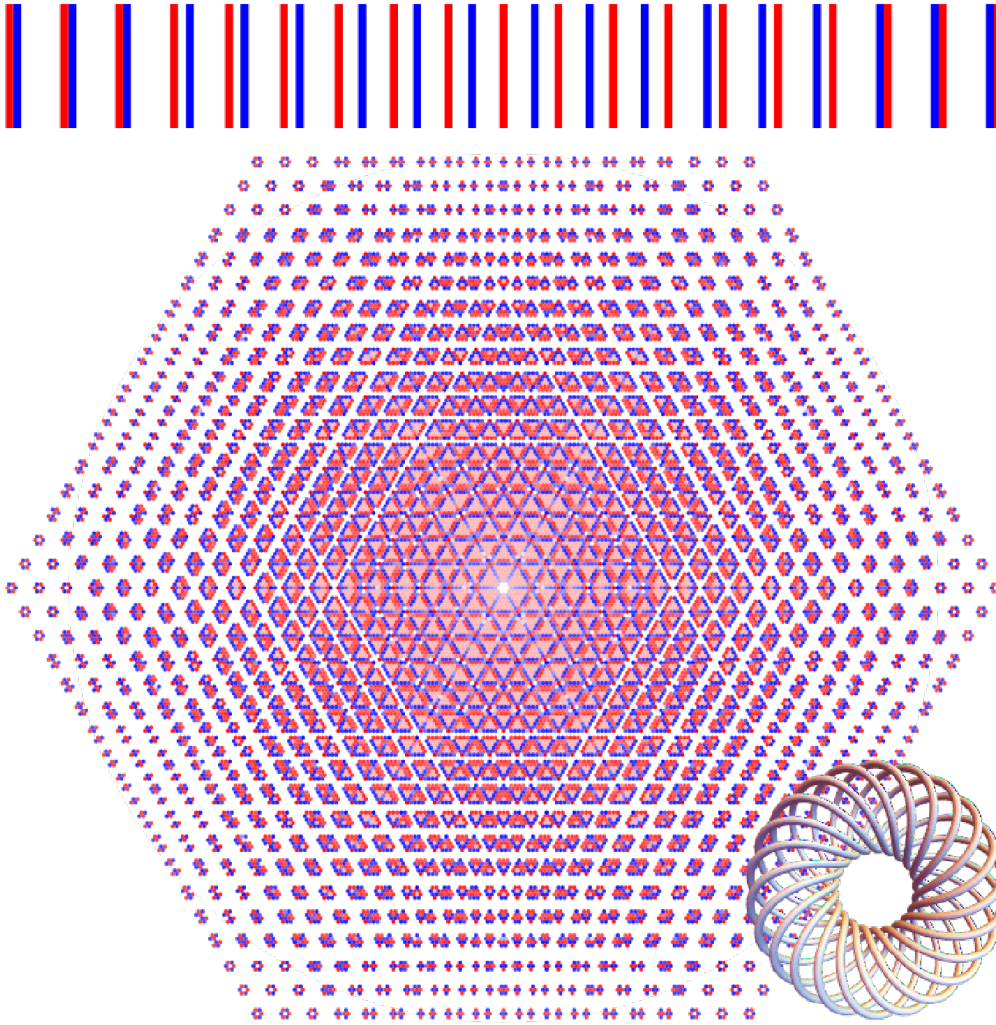
**Fun.**  $\Theta$  on Rolfsen's Table:



The 132-crossing torus knot  $T_{22/7}$ :

(many more at  $\omega\epsilon\beta/\text{TK}$ )

Random knots from [DHOEBL] with 51 – 75 crossings: (many more at  $\omega\epsilon\beta/\text{DK}$ )



**Moral.** We must come to terms with  $\Theta$ !

**Task 1.** Make the “data” formulas human friendly.

**Task 2.** Prove the hexagonal symmetry of  $\theta(K)$ , and that  $\theta(K) = \theta(-K) = -\theta(\bar{K})$ .

That’s harder than it seems! The formulas don’t naively show any of that.  $\Delta$  has a palindromic symmetry first conjectured in Alexander’s original paper [A1] — it is invariant under  $T \rightarrow T^{-1}$ . Proving this took a few years, and the proof starting from the Wirtinger presentation is quite involved (e.g. [CF, Chapter IX]).

**Task 3.** Show that  $\theta$  dominates the Rozansky-Overbay invariant  $\rho_1$  [Ro1, Ro2, Ro4, Ov, BV1]. Precisely, show that  $\rho_1 = -\theta|_{T_1 \rightarrow T, T_2 \rightarrow 1}$ .

This one should be easy with techniques from [BV3, Section 4.2].

**Task 4.** Explain the “Chladni patterns”. Are there “dominant parts” of  $\theta$  that can be computed in isolation?



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**Task 5.** Prove the genus bound of Conjecture 1.

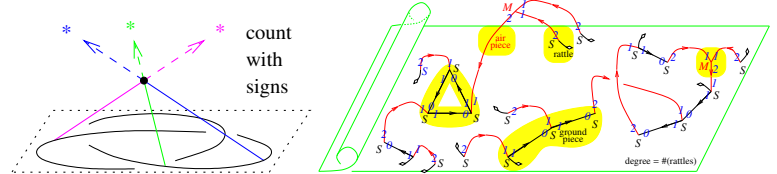
This is probably coming. One can bound the degree of  $\Delta = \det(A)$  in terms of  $g(K)$  using the Seifert presentation of the Alexander module. Pushing further, likely one can bound the degree of

$(g_{\alpha\beta}) = A^{-1}$  in terms of  $g(K)$ , and that’s probably enough.

**Task 6.** Find a 3D interpretation of the  $g_{\alpha\beta}$ ’s.

They must be closely related to the equivariant linking numbers of [KY, GK, GT, Oh3, Le1].

**Task 7.** Find a formula  $\mathcal{F}$  for  $\Theta(K)$  that starts from a Seifert surface  $\Sigma$  of  $K$ . Better if  $\mathcal{F}$  is completely 3D! Assuming Task 13, it is known that  $\Theta$  depends only of invariants of type  $\leq 3$  of  $\Sigma$ . Maybe  $\mathcal{F}$  is about configuration space integrals / chopstick towers? See CS: [Th, Le2, BN1], BF: [CR, BN2]



**Task 8.** Is there an intrinsic theory of finite type invariants for Seifert surfaces? For task 11, does its  $gr$  map to functions on  $H_1$ ?

My current best understanding of finite type invariants for Seifert surfaces goes through thick graphs.



**Task 9.** Prove the the fibered condition of Conjecture 2.

If  $K$  is fibered,  $\deg \Delta(K) = g(K)$  and  $\Delta(K)$  is monic. Indeed,  $K$  is then the mapping cylinder of a diffeomorphism  $f: \Sigma \rightarrow \Sigma$ . The Alexander module of  $K$  is generated by  $H_1(\Sigma)$  with relations

$\{\gamma = T f_* \gamma : \gamma \in H_1(\Sigma)\}$ . Thus the highest monomial in  $\Delta$  is  $T^g \det(f_*)$  and  $\det(f_*) = \pm 1$  as  $f_*$  preserves the intersection pairing. If only we had a formula for  $\theta$  in terms of  $f \dots$

**Task 10.** In general, find a formula for  $\Theta$  corresponding to each known presentation of the Alexander module.

Wirtinger is  $2\{\text{xings}\} \rightarrow \{\text{edges}\}$ . Dehn is  $\{\text{xings}\} \rightarrow \{\text{faces}\}$ . Co-Dehn is  $\{\text{faces}\} \rightarrow \{\text{xings}\}$ . Burau is  $\{\text{braid strands}\} \rightarrow \{\text{braid strands}\}$ . Seifert is  $H_1(\Sigma) \rightarrow H_1(\Sigma)$ , and so is the presentation from Task 9. Grid diagrams lead to  $\{\text{grid number}\} \rightarrow \{\text{grid number}\}$  (may relate to HFK). There's more!

**Task 11.** Write up the integration story.

**Claim** (e.g., [BN4]). Cutting corners, with  $\epsilon^2 = 0$ ,

$$\frac{1}{\Delta_1 \Delta_2 \Delta_3} \exp\left(\epsilon \cdot \frac{\theta}{\Delta_1 \Delta_2 \Delta_3}\right) \sim \oint_{\prod_e \mathbb{R}^6_{p_{1e}, p_{2e}, p_{3e}, x_{1e}, x_{2e}, x_{3e}}} \prod_C e^{L_C},$$

where  $\oint$  denotes perturbed formal Gaussian integration (i.e., “Feynman Diagrams”) and  $L_C$  is

$$\begin{aligned} L[\mathbf{x}_{i,j}, \mathbf{s}_{-}] := & \text{Plus} [ \\ & \sum_{v=1}^3 (\mathbf{x}_{vi} (\mathbf{p}_{vi}^+ - \mathbf{p}_{vi}) + \mathbf{x}_{vj} (\mathbf{p}_{vj}^+ - \mathbf{p}_{vj}) + (\mathbf{T}_v^s - 1) \mathbf{x}_{vi} (\mathbf{p}_{vi}^+ - \mathbf{p}_{vj}^+)), \\ & (\mathbf{T}_1^s - 1) \mathbf{p}_{3j} \mathbf{x}_{1i} (\mathbf{T}_2^s \mathbf{x}_{2i} - \mathbf{x}_{2j}), \\ & \epsilon \in (\mathbf{T}_3^s - 1) \mathbf{p}_{1j} (\mathbf{p}_{2i} - \mathbf{p}_{2j}) \mathbf{x}_{3i} / (\mathbf{T}_2^s - 1), \\ & \epsilon \in (1/2 + \mathbf{T}_2^s \mathbf{p}_{1i} \mathbf{p}_{2j} \mathbf{x}_{1i} \mathbf{x}_{2i} - \mathbf{p}_{1i} \mathbf{p}_{2j} \mathbf{x}_{1i} \mathbf{x}_{2j} - \mathbf{p}_{3i} \mathbf{x}_{3i} - (\mathbf{T}_2^s - 1) \mathbf{p}_{2j} \mathbf{p}_{3i} \mathbf{x}_{2i} \mathbf{x}_{3i} + \\ & (\mathbf{T}_3^s - 1) \mathbf{p}_{2j} \mathbf{p}_{3j} \mathbf{x}_{2i} \mathbf{x}_{3i} + 2 \mathbf{p}_{2j} \mathbf{p}_{3i} \mathbf{x}_{2j} \mathbf{x}_{3i} + \mathbf{p}_{1i} \mathbf{p}_{3j} \mathbf{x}_{1i} \mathbf{x}_{3j} - \mathbf{p}_{2i} \mathbf{p}_{3j} \mathbf{x}_{2i} \mathbf{x}_{3j} - \\ & \mathbf{T}_2^s \mathbf{p}_{2j} \mathbf{p}_{3j} \mathbf{x}_{2i} \mathbf{x}_{3j} + \\ & ((\mathbf{T}_1^s - 1) \mathbf{p}_{1j} \mathbf{x}_{1i} (\mathbf{T}_2^s \mathbf{p}_{2j} \mathbf{x}_{2i} - \mathbf{T}_2^s \mathbf{p}_{2j} \mathbf{x}_{2j} - (\mathbf{T}_2^s + 1) (\mathbf{T}_3^s - 1) \mathbf{p}_{3j} \mathbf{x}_{3i} + \\ & \mathbf{T}_2^s \mathbf{p}_{3j} \mathbf{x}_{3j}) + (\mathbf{T}_3^s - 1) \mathbf{p}_{3j} \mathbf{x}_{3i} \\ & (1 - \mathbf{T}_2^s \mathbf{p}_{1i} \mathbf{x}_{1i} + \mathbf{p}_{2i} \mathbf{x}_{2j} + (\mathbf{T}_2^s - 2) \mathbf{p}_{2j} \mathbf{x}_{2j})) / (\mathbf{T}_2^s - 1) ) ] \end{aligned}$$

In fact, we first found  $L_C$  using the method of undetermined coefficients, and then derived  $F_1$  and  $F_2$  from it.

**Task 12.** Find a similar perturbed Gaussian integral formula for  $\theta$ , but with integration over  $6H_1(\Sigma)$ . The quadratic  $Q$  will be the same as in the Seifert-Alexander formula (but repeated 3 times, for each  $T_v$ ). The perturbation  $P_\epsilon$  will be given by low-degree finite type invariants of curves on  $\Sigma$  (possibly also dependent on the intersection points of such curves, or on other information coming from  $\Sigma$ ).

**Task 13.** Prove that  $\theta$  is equal to the two-loop contribution  $Z^{(2)}$  to the Kontsevich integral  $Z$ .

Composed with the inverse PBW isomorphism  $\chi^{-1}$ ,  $\chi^{-1} \circ Z$  takes values in univalent Jacobi diagrams,  $\mathcal{B} = \{\boxtimes \circ \dots\} / IHX$ . Rozansky conjectured [Ro3, GR] and Kricker proved [Kr] that

$$\log(\chi^{-1} \circ Z) = f_1 \left( \begin{array}{|c|} \hline t \\ \hline \end{array} \right) + f_2 \left( \begin{array}{|c|} \hline t_1 \\ \hline t_2 \\ \hline \end{array} \right) + \text{higher loops},$$

where  $t^k \begin{array}{|c|} \hline t \\ \hline \end{array} := \begin{array}{|c|} \hline | \\ \hline \end{array} \dots n \dots \begin{array}{|c|} \hline | \\ \hline \end{array}$ ,  $f_1 \in \mathbb{Q}[[t]]$ , and  $f_2 \in \mathbb{Q}[[t_1, t_2]]$  satisfy  $f_1 = \frac{1}{2} \log \frac{\sinh(t/2)}{t \Delta(e^t/2)}$  and  $f_2 = Z^{(2)}(e^{t_1}, e^{t_2}) / \Delta(e^{t_1}) \Delta(e^{t_1}) \Delta(e^{t_1+t_2})$  where  $Z^{(2)} \in \mathbb{Z}[T_1^{\pm 1}, T_2^{\pm 1}]$  is the “two loop polynomial”. Ohtsuki [Oh2] studied  $Z^{(2)}$  extensively, and almost certainly,  $Z^{(2)} = \theta$ . Prove that!

**Task 14.** Complete and write up the  $\mathfrak{g}_\epsilon^+$  story.

Let  $\mathfrak{g}$  be a semisimple Lie algebra, let  $\mathfrak{h}$  be its Cartan subalgebra, and let  $\mathfrak{b}^u$  and  $\mathfrak{b}^l$  be its upper and lower Borel subalgebras. Then  $\mathfrak{b}^u$  has a bracket  $\beta$ , and as the dual of  $\mathfrak{b}^l$  it also has a cobracket  $\delta$ , and in fact,  $\mathfrak{g} \oplus \mathfrak{h} \equiv \text{Double}(\mathfrak{b}^u, \beta, \delta)$ . Let  $\mathfrak{g}_\epsilon^+ := \text{Double}(\mathfrak{b}^u, \beta, \epsilon \delta) \pmod{\epsilon^{d+1}}$  it is solvable for any  $d$ . We expect that  $\Theta$  is the universal invariant (in the sense of Lawrence and Ohtsuki [La, Oh1]) corresponding to  $sl_{3,\epsilon}^+$ , computed modulo  $\epsilon^2$  (in fact, that's how we guessed it). See [BN3, BV2].

**Task 15.** Go beyond  $sl_3$  and the first power of  $\epsilon$ !

This sounds very appealing, and you will indeed get stronger and stronger invariants. But they will become less and less computable  $\odot$ .

**Task 16.** Relate the  $\mathfrak{g}_\epsilon^+$  story with (rotational) virtual knots [Kau], with  $\vec{\mathcal{A}}$  [Po], and with quantization of Lie bialgebras [EK1, EK2, En, Se]

We expect that there is a commutative diagram as

Re top on or the right which descends to the bottom one, with  $\oplus$  Re case of  $\mathfrak{g} = \mathfrak{sl}_3$   $\downarrow k=1$ . But  $\odot$

**Task 17.** Find a w-style characterization of  $\Theta$ .

A word about [GR]. Compare with (say) [BD], where  $\Delta$  is characterized on w-knots by the overcrossings / tails commute relation. Similarly it should be possible to characterize  $\Theta$  on rotational virtual knots by some “overcrossings / tails nearly commute” relation.

**Task 18.** Understand Chern-Simons theory with gauge group  $\mathfrak{g}_\epsilon^+$ .

Is there a gauge that leads to the formula  $\mathcal{F}$  of Task 7?

**Task 19.** What happens to representation theory as  $\epsilon \rightarrow 0$ ? Is there any fun in continuous morphisms  $\mathfrak{g}_\epsilon^+ \rightarrow \mathfrak{gl}_{n,\epsilon}^+$ ?

**Task 20.** Does  $\Theta$  extend to knots in  $\mathbb{Z}HS/\mathbb{Q}HS$ ?

$Z$  and  $Z^{(2)}$  do.

**Task 21.** Is there a surgery formula for  $\Theta$ ?



Z and Z<sup>(2)</sup> have.

**Task 22.** Extend  $\Theta$  to tangles and figure out how it behaves under strand doubling.

Z and Z<sup>(2)</sup> extend but their extensions depend on parenthesizations. From Task 14 we expect that  $\Theta$  will extend without the need for parenthesizations, yet with an asymmetry built into the doubling operations.

**Task 23.** Make Kricker / Ohtsuki [Kr, Oh2] more computable!

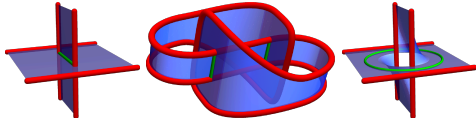
**Task 24.** Find a multi-variable version of  $\theta$  for links, like there is a multi-variable Alexander for links (e.g. [Kaw, Chapter 7]). It is predicted  $g_e^+$  consideration, but not by the loop expansion.

**Task 25.** Find a ribbon condition satisfied by  $\Theta$ .

For a ribbon knot K, one may find a Seifert surface  $\Sigma$  half of whose homology is generated by the components of an unlink embedded in  $\Sigma$ . This makes for a presentation matrix A of the Alexander module of K that has big blocks of zeros, and this leads to the Fox-Milnor condition [FM],  $\Delta \doteq \det(A) \doteq f(T)f(T^{-1})$  for some  $f \in \mathbb{Z}[T^{\pm 1}]$ . If  $\det A$  is constrained for ribbon knots, perhaps so is  $A^{-1}$  and therefore  $\Theta$ ?

**Bonus Task.** Carthago delenda est and every knot polynomial must be categorified.

M. Khovanov & Cato the Elder



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A: Note that tangles & strand doubling are keys to “Algebraic knot Theory” [AKTJ].

better if burger.

## A FAST, STRONG, TOPOLOGICALLY MEANINGFUL, AND FUN KNOT INVARIANT

DROR BAR-NATAN AND ROLAND VAN DER VEEN

ABSTRACT. In this paper we discuss a pair of polynomial knot invariants  $\Theta$  and  $\Delta$  which

- Theoretically and practically fast,  $\Theta$  can be computed in polynomial time. We can compute it in full on random knots with over 300 crossings, and its evaluation at simple rational numbers on random knots with over 600 crossings.
  - Strong. Its separation power is much greater than the hyperbolic volume, the HOMFLY-PT polynomial and Khovanov homology (taken together) on knots with up to 15 crossings while being computable on much larger knots.
  - Topologically meaningful. It likely gives a genus bound, and there are reasons to hope that it would be more.
  - Fun. See below to Figures 1.1–1.4, 3.1, and 6.2.
- $\Delta$  is merely the Alexander polynomial.  $\Theta$  is almost certainly equal to an invariant that was studied extensively by Ohtsuki [Oh2], containing Rozansky, Kricker, and Garoufalidis [Roz1, Roz2, Roz3, Kr, GH]. Yet our formulas, proofs, and programs are much simpler and enable the computation even on very large knots.

## CONTENTS

1. Fun
2. The Main Theorem
3. Implementation and Examples
- 3.1. Implementation
- 3.2. Examples
4. Proof of the Main Theorem, Theorem 1
- 4.1. Proof of Invariance
- 4.2. Proof of Polynomiality
5. Strong and Meaningful
- 5.1. Strong
- 5.2. Meaningful
- 5.2.1. The Knot Genus
- 5.2.2. Filtered Knots
6. Stories, Conjectures, and Dreams
7. Acknowledgement
- References

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This paper is available in electronic form, along with source files and a demo Mathematica notebook at [http://oeß.net/Theta](#) and at [arXiv:2509.18456](#).

1. FUN  
The word “fun” rarely appears in the title of a math paper, so let us start with a brief justification.

$\Theta$  is a pair of polynomials. The first,  $\Delta$ , is old news, the Alexander polynomial [AJ]. It is a one-variable Laurent polynomial in a variable  $T$ . For example,  $\Delta(\bigcirc) = T^{-1} - 1 + T$ . We turn such a polynomial into a list of coefficients (for  $\Delta$ , it is  $(1, -1, 1)$ ), and then to a chain of bars of varying colours: white for the zero coefficients, and red and blue for the positive and negative coefficients (with intensity proportional to the magnitude of the coefficients).

Similarly,  $\Theta$  is a 2-variable Laurent polynomial, in variables  $T_1$  and  $T_2$ . We can turn such a polynomial into a 2D array of coefficients and then using the same rules, into a 2D array of colours, namely, into a picture. To highlight a certain conjectured hexagonal symmetry of the resulting pictures, we apply a shear transformation to the plane before printing. So a monomial  $cT_1^a T_2^b$  gets printed at position  $(a - b/2, \sqrt{3}b/2)$  instead of the more straightforward  $(a, b)$ . On the right is the 2D picture corresponding to the polynomial  $2 + T_1 - T_2 + T_1 T_2 + T_2 - T_1^{-1} + T_1^{-1} T_2^{-1} - T_2^{-1}$ .

Thus  $\Theta$  becomes a pair of pictures: a bar code, and a 2D picture that we call a “hexagonal QR code”. For the knots in the Rolfsen table (with the unknot prepended at the start), they are in Figure 1.1. For some alternating square wave knots, they are in Figure 1.2, and for a random square wave, in Figure 1.3. In addition, the hexagonal QR codes of 15 knots with  $\geq 300$  crossings are in Figure 1.4, and  $\Theta$  of a 132-crossing torus knot is in Figure 3.1. Some further computations and figures, also highlighting the parity of coefficients rather than just their signs, are at [La].

Clearly there are patterns in these figures. There is a hexagonal symmetry and the QR codes are nearly always hexagons (these are independent properties). Much more can be seen in Figure 1.1. In Figure 1.4 there seem to be large-scale patterns perhaps reminiscent of the “Chladni figures” formed by powders atop vibrating plates (on right). We can’t prove any of these things, and the last one, we can’t even formulate properly. Yet they are clearly there, too clear to be the result of chance alone. We plan to have fun over the next few years observing and proving these patterns. We hope that others will join us too.

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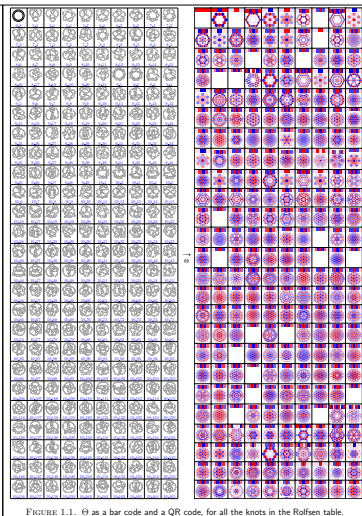
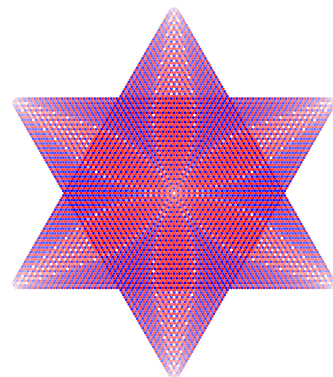


FIGURE 1.1.  $\Theta$  as a bar code and a QR code, for all the knots in the Rolfsen table.



A (2, 41, -41) pretzel for dessert



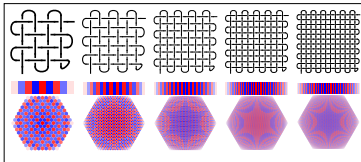


FIGURE 1.2.  $\Theta$  of some square wave knots, as computed by [BV3, WaveKnots.nb]

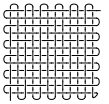


FIGURE 1.3.  $\Theta$  of a randomized wave knot, as computed by [BV3, WaveKnots.nb]. Crossings were chosen to be positive or negative with equal probabilities.

## 2. THE MAIN THEOREM

We start with the definition of  $\Theta$ . Given an oriented  $n$ -crossing knot  $K$ , we draw it in the plane as a long oriented diagram  $D$  in such a way that the two strands intersecting at each crossing are pointing up and down, respectively. We call such a diagram an *upright knot diagram*. An example of an upright knot diagram is shown on the right.



FIGURE 2.1. An example upright knot diagram.

We then label each edge of the diagram with two labels: a running index  $k$  which runs from 1 to  $2n+1$ , and a "rotation number"  $\rho_k$ , the geometric rotation number of that edge.<sup>1</sup>

The signed number of times the tangent to the edge is horizontal and heading right, with caps counted with  $\pm$  signs and caps with  $-$ . This number is well defined because at their ends, all edges are headed up.

In particular, the middle diagram which resembles the Greek letter  $\Theta$  gave the invariant  $\Theta$  name.

**Comment 1.** The computation of  $\Theta$  is a bottleneck for the computation of  $\Theta$ . It requires inverting a  $(2n+1) \times (2n+1)$  matrix whose entries are (degree 1) Laurent polynomials in  $T$ . It's a daunting task yet it takes polynomial time. Even a naive inversion using Gaussian elimination requires only  $n^3$  operations in the ring  $\mathbb{Q}[T]$ . So  $G$  can be computed in practice even if  $n$  is in the hundreds, and everything which follows is not worse.

The polynomials  $F_1(c)$ ,  $F_2(c)$ , and  $F_3(c)$  are not unique, and we are not certain that we have the chosen elements from that ring. They are ugly from a human perspective, but from a computational perspective, having 18 terms (as in the case of  $F_1(c)$ ) isn't really a problem; computers don't care.

Computationally, the worst term in (6) is the middle one, and even it takes merely  $n^2$  operations in the ring  $\mathbb{Q}[T]$  to evaluate.

## 3. IMPLEMENTATION AND EXAMPLES

**3.1. Implementation.** A concise yet reasonably efficient implementation is worth a thousand formulas. It completely removes ambiguities, it is testable, and it allows for experimentation. Here we use the Wolfram Language. The section that follows was generated from a Mathematica [Wol] notebook which is available at [BV3, Theta.nb]. A second implementation of  $\Theta$ , using Python and SageMath (<https://www.sagemath.org/>) is available at <https://www.knottheory.org/Theta/>.

We start by loading the package **KnotTheory** — it is only needed because it has many specific knots predefined. In this Section and in the next,  $\Theta$  and  $\Theta$  mean "human input" while  $\Theta$  means "computer output".

**OneClick[KnotTheory]** Loading KnotTheory version of October 29, 2024, 18:29:52.1861. Read more at <http://knottheory.org/wiki/KnotTheory>.

Next we quietly define the modules **Rot**, **Exp**, used to compute rotation numbers, and **PolysFct**, used to plot polynomials as bar codes and as hexagonal QR codes. Neither is a part of the core of the computation of  $\Theta$ , so neither is shown; yet we do show one usage example for each.

**Rot[RotMatrix[Knot[3, 1]]]**  $\left\{\left\{\left\{1, 1, 4\right\}, \left\{1, 3, 6\right\}, \left\{1, 5, 2\right\}\right\}, \left\{\left\{0, 0, -1, 0, 0\right\}, \left\{0, 0, -1, 0, 0\right\}, \left\{0, 0, -1, 0, 0\right\}\right\}\right\}$

**Exp[Exp[27 + 17T, -1 + 2T, -1 + 2T, 1 + 4T + 17T]]**

**PolysFct[120, Labeled = True]**

We use the reader to compare the above output with the knot diagram in Figure 2.1.

**PolysFct[120, Labeled = True]**

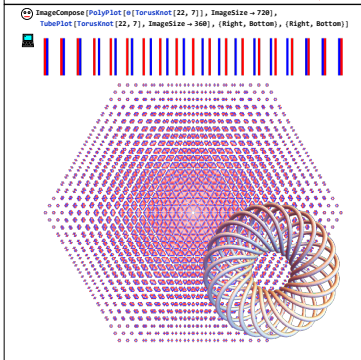


FIGURE 3.1. The 132-crossing torus knot  $T_{22,7}$  and a plot of its  $\Theta$  invariant

$i^*$  (these are  $+1$  and  $+1+2i$  if the labelling is by consecutive integers). Also, by convention  $T^*$  will always refer to the label of the first edge, and  $2n+1$  will always refer to the label of the last. With this in mind, we have that  $A = I + \sum_{k=1}^{2n} A_k$ , with  $A_k$  given by

$$j^* \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \rightarrow \frac{A_k}{\text{row } j^*} \begin{pmatrix} \text{column } j^* & \text{column } j^* + 1 \\ \text{row } j^* & \text{row } j^* + 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad (7)$$

Like in [BV1, Lemma 3], the equalities  $AG = I$  and  $GA = I$  imply that for any crossing  $c = (s, i, j)$  in a knot diagram  $D$ , the Green function  $G = (g_{\alpha\beta})$  of  $D$  satisfies the following

$$g_{\alpha\beta} = g_{\beta\alpha} + T g_{\alpha, \beta+1} + T(1-T) g_{\alpha, \beta+2} + (1-T) g_{\alpha, \beta+3} \quad (13)$$

and

$$g_{\alpha\beta} = g_{\beta\alpha} + T g_{\alpha, \beta-1} + T(1-T) g_{\alpha, \beta-2} + (1-T) g_{\alpha, \beta-3} \quad (14)$$

Next we deal with the case of  $R2^*$ . We use the privileges afforded by us by Lemma 7 to insert  $A$  into the right-hand-side of the move, and like in the case of  $R3$ , we start with pictures annotated with the relevant type (5) and (11)  $g$ -rules, written with the assumption that  $\beta \neq i^*, j^*$ :

As in the case of  $R3$ , we eliminate  $g_{\alpha\beta}$  and  $g_{\beta, \alpha}$  from the equations for the left hand side, and find that for the purpose of determining  $g_{\alpha\beta}$  with  $\beta \neq i^*, j^*$ , they are equivalent to the equations

$$g_{\alpha\beta} = g_{\beta\alpha} + g_{\alpha, \beta+1} + g_{\beta, \alpha+1} \quad (15)$$

and as in the case of  $R3$ , this establishes the invariance of  $g_{\alpha\beta}$  under  $R2^*$  moves.

For the remaining moves,  $R2^*$ ,  $R1$ , and  $R1^*$ , we merely display the  $g$ -rules and leave it to the readers to verify that when the edges  $i^*$  and  $(j^*+1)$  are eliminated, the left hand sides

Figure 2.1 the running index runs from 1 to  $7$ , and the rotation numbers for all edges are 0 (and hence are omitted) except for  $\rho_6$  which is  $-1$ .

Let  $X$  be the set of all crossings in the diagram  $D$ , where we encode each crossing as a triple (sign of the crossing, incoming edge index, outgoing edge index). In our example we have  $X = \{(1, 1, 4), (1, 3, 6), (1, 5, 2)\}$ .

We let  $A$  be the  $(2n+1) \times (2n+1)$  matrix of Laurent polynomials in a variable  $T$ , defined by

$$A = I - \sum_{c=(s,i,j) \in X} (T E_{i,j} + (1-T) E_{j,i} + E_{j,i+1} + E_{i+1,j})$$

where  $E$  is the identity matrix and  $E_{i,j}$  denotes the elementary matrix with a 1 in row  $i$  and column  $j$  and zeros elsewhere.

The definition of  $GF$  above is a tedious task to compute how to best store polynomials in the case of  $F_1$  and  $F_2$ . The programs would run the list sum without it, albeit a bit more slowly.

**GF[F1[c], F2[c], F3[c]]**  $\{F_1[c], F_2[c], F_3[c]\}$

Next, we define that  $T_1 = T_2 = T$  and define the three "Feynman Diagram" polynomials  $F_1, F_2, F_3$  by

$$F_1[c] = T_1 F_1[c] + T_2 F_2[c] + T_3 F_3[c]$$

Next comes the main program computing  $\Theta(K)$ . Fortunately, it matches perfectly with the mathematical description in Section 2. In 1. below we let  $k$  be the length of  $X$ , namely, the number of crossings in  $K$ , and we let the starting value of  $k$  be the  $(2n+1) \times (2n+1)$  identity matrix. Then in line 2, for each crossing in  $X$  we add to  $A$  a  $2 \times 2$  block, in rows  $i$  and  $j$  and columns  $i+1$  and  $j+1$ , as explain in Equation (1). In line 3 we compute the normalized Alexander polynomial  $\Delta$  as in (2). In line 4 we let  $G$  be the inverse of  $A$ . In line 5 we declare that it seems to evaluate, we use formula  $\mathcal{E}$  that may contain symbols of the form  $g_{\alpha\beta}$ , each such symbol is to be replaced by the entry in position  $\alpha, \beta$  of  $G$ , but  $T$  replaced with  $T_1$ . In line 6 we start computing  $\Theta$  by computing the first summand in (6), which in itself, is a sum over the crossings of the knot. In line 7 we add to  $\Theta$  the double sum corresponding to the second term in (6), and in line 8, we add the third summand of (6). Finally, line 9 outputs a pair  $\Delta$ , and the re-normalized version of  $\Theta$ .

"sigma", with  $\delta$  denoting the Kronecker delta:

$$g_{\alpha\beta} = g_{\beta\alpha} + T g_{\alpha, \beta+1} + (1-T) g_{\alpha, \beta+2} + (1-T) g_{\alpha, \beta+3} \quad (8)$$

$$g_{\alpha\beta} = g_{\beta\alpha} + T g_{\alpha, \beta-1} + (1-T) g_{\alpha, \beta-2} + (1-T) g_{\alpha, \beta-3} \quad (9)$$

Furthermore, the system of equations (8) is equivalent to  $AG = I$  and so it fully determines  $g_{\alpha\beta}$  and likewise for the system (9), which is equivalent to  $GA = I$ .

Of course, the same  $g$ -rules also hold for  $G = (g_{\alpha\beta})$  for  $\alpha = 1, 2, 3$ , except with  $T$  replaced with  $T^*$ .

We also need a variant  $g_{\alpha\beta}$  defined whenever  $\alpha$  and  $\beta$  are two distinct points on the edges of a knot diagram  $D$ , away from the crossings. If  $\alpha$  is the edge on which  $\alpha$  lies and  $\beta$  is the edge on which  $\beta$  lies,  $g_{\alpha\beta}$  is defined as follows:

$$g_{\alpha\beta} = \begin{cases} g_{\beta\alpha} & \text{if } \alpha = \beta \\ g_{\alpha\beta} & \text{if } \alpha \neq \beta \text{ and } \alpha < \beta \text{ relative to the orientation of the edge } \alpha = \beta \end{cases} \quad (10)$$

Of course, we can define  $g_{\alpha\beta}$  from  $g_{\beta\alpha}$  in a similar way. It is clear that  $g$  and  $g$  contain the same information and are easily computable from each other. The variant  $g$  is, strictly speaking, not a matrix and so  $g$  is a bit more suitable for computations. Yet  $g$  is a bit better behaved when we try to track, as below, the changes in  $g$  and  $g$  under Reidemeister moves. Reidemeister moves sometimes merge two edges into one or break an edge into two. In such cases the points  $\alpha$  and  $\beta$  can be "pulled" along with the move so as to retain their ordering along the overall parametrization of the knot, yet somewhat edge labels lose this information. From the perspective of traffic functions,  $g$  is somewhat more natural than  $g$ , as it makes sense to inject traffic and to count traffic anywhere along an edge, provided the injection point and the counting point are distinct.

The following discussion and lemma further exemplify the advantage of  $g$  over  $g$ :

**Discussion 6.** We introduce "we shall see" as on the right into knot diagrams, whose only function (as we will see) is to cut edges into parts that may carry different labels. When dealing with upright knot diagrams as in Figure 2.1, we may carry different labels when the tangent to the knot is pointing up, so that the rotation numbers  $\rho_k$  remain well defined on all edges. In the presence of nil vertices the matrix  $A$  becomes a bit larger (by as many null vertices as added to a knot diagram). The rule (7) for the creation of the matrix  $A$  gets an amendment for null vertices:

$$j^* \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \rightarrow \frac{A_k}{\text{row } j^*} \begin{pmatrix} \text{column } j^* & \text{column } j^* + 1 \\ \text{row } j^* & \text{row } j^* + 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad (7)$$

and the summation for  $A = I + \sum_{k=1}^{2n} A_k$  is extended to include summands for all the null vertices. The matrix  $G = A^{-1}$  and the function  $g_{\alpha\beta}$  are defined as before. The  $g$ -rules of (8) and (9) now become

$$g_{\alpha\beta} = g_{\beta\alpha} + g_{\alpha, \beta+1} \quad (11)$$

$$g_{\alpha\beta} = g_{\beta\alpha} + g_{\alpha, \beta-1} \quad (12)$$

and it remains true that the system of equations (8)-(11) (as well as (8)-(12)) fully determines  $g_{\alpha\beta}$ . The variant  $g_{\alpha\beta}$  is also defined as before, except now  $\alpha$  and  $\beta$  need to also be away from the null vertices.

Figure 4.3: The upright Reidemeister moves. The  $R1$  and  $R3$  moves are already upright and remain the same as in Figure 4.2. The crossings in the  $R2$  moves of Figure 4.2 are related to be upright. We also show two further moves. The null vertex move  $IV$  for adding and removing null vertices, and the swirl move  $Sw$  which implies that any two ways of turning a crossing upright are the same. We sometimes indicate rotation numbers symbolically rather than using complicated signs.

become equivalent to the right hand sides:

We can now move on to the main part of the proof of our Main Theorem, Theorem 1. We need to show the invariance of  $\Theta$  under the "upright Reidemeister" moves of Figure 4.3.

**Proposition 10.** The moves in Figure 4.3 are sufficient. If any upright knot diagrams (with null vertices) represent the same knot, they can be connected by a sequence of moves as in the figure.

**Proof Sketch.** There is an obvious well-defined map

upright knot diagram  $\rightarrow$  oriented knot diagram

relations as in Figure 4.3  $\rightarrow$  relations as in Figure 4.2

We then have to construct an inverse to that map. To do that we have to choose how to turn each crossing in an oriented knot diagram to be upright. The different ways of doing so differ by instances of the Sw relation (if deeper spins need to be swirled away, null vertices may be inserted using  $IV$  and the spins can be undone one rotation at a time). A more detailed version of the proof is in [BV3].

**Proposition 11.** The quantity  $\Theta$  is invariant under  $R1$ .

Alternatively,  $A = I + \sum_{k=1}^{2n} A_k$ , where  $A_k$  is a matrix of zeros except for the blacks as follows:

$$j^* \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \rightarrow \frac{A_k}{\text{row } j^*} \begin{pmatrix} \text{column } j^* & \text{column } j^* + 1 \\ \text{row } j^* & \text{row } j^* + 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad (1)$$

We note that the determinant of  $A$  is equal up to a unit to the normalized Alexander polynomial  $\Delta$  of  $K$ .<sup>2</sup> In fact, we have

$$\Delta(K) = \Delta(A) = T^{n-\text{ord}(D)/2} \det(A), \quad (2)$$

where  $\text{ord}(D) = \sum_{s=1}^n \rho_s$  is the total rotation number of  $D$  and where  $\text{ord}(D) = \sum_{s=1}^n \rho_s$  is the writhe of  $D$ , namely the sum of the signs  $s_i$  of all the crossings  $c_i$  in  $D$ .

We let  $G = (g_{\alpha\beta}) = A^{-1}$ , and thinking of it as a function  $g_{\alpha\beta}$  of a pair of edges  $\alpha$  and  $\beta$ , we call it the Green function of the diagram  $D$ . When inspired by physics (e.g. Fact 33 and [BNS]) we sometimes call it "the 2-point function", and when thinking of car traffic (e.g. Comment 3 and [BV1, BNT]) we sometimes call it "the traffic function". As an example, here are  $A$  and  $G$  for the knot diagram  $D$  of Figure 2.1:

$$A = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \quad T = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

Let  $T_1$  and  $T_2$  be indeterminates and let  $T_3 = T_1 T_2$ . Let  $\Delta_{\alpha\beta} = \Delta_{\alpha\beta, T_1, T_2}$  and  $G_{\alpha\beta} = G_{\alpha\beta, T_1, T_2}$  be  $A$  and  $G$  subject to the substitution  $T \rightarrow T_1 T_2$ , where  $\alpha = 1, 2, 3$ . Given crossing  $c = (s, i, j)$ ,  $c_1 = (s, i, j)$ , and  $c_2 = (s, i, j)$  in  $X$  and an edge label  $k$ , let

$$F_1(c) = \frac{1}{2} (T_1 - T_2) g_{\alpha\beta} + T_1 g_{\alpha, \beta+1} + T_2 g_{\alpha, \beta+2} + (T_1 - T_2) g_{\alpha, \beta+3} \quad (3)$$

$$F_2(c) = \frac{1}{2} (T_1 - T_2) g_{\alpha\beta} + T_1 g_{\alpha, \beta-1} + T_2 g_{\alpha, \beta-2} + (T_1 - T_2) g_{\alpha, \beta-3} \quad (4)$$

$$F_3(c) = (g_{\alpha\beta} - 1/2) g_{\alpha\beta} \quad (5)$$

These formulas are uninspiring, yet they are easy to compute (given  $G$ ), and they work:

<sup>2</sup>The informed reader will notice that  $A$  is a presentation matrix for the Alexander module of  $K$ , obtained by using Fox calculus on the Wirtinger presentation of the fundamental group of the complement of  $K$ .

**Exp[Exp[27 + 17T, -1 + 2T, -1 + 2T, 1 + 4T + 17T]]**

**PolysFct[120, Labeled = True]**

**PolysFct[120, Labeled = True]**

**3.2. Examples.** On to examples! Starting with the trefoil knot.

**Exp[Exp[27 + 17T, -1 + 2T, -1 + 2T, 1 + 4T + 17T]]**

**PolysFct[120, Labeled = True]**

Next are the Conway knot  $11_{45}$  and the Kinoshita-Terasaka knot  $11_{45}$ . The two are mutants and famously hard to separate: they have the same  $h$ -invariant (as evidenced by their mirror Alexander bar codes below), and they have the same hyperbolic volume, HOMFLY-PT polynomial, and Khovanov homology. Yet their  $\Theta$  invariants are different. The Conway knot is  $2$ , while the genus of the Kinoshita-Terasaka knot is  $2$ . This agrees with the apparent higher complexity of the QR code of the Conway polynomial and with Conjecture 18 below.

**PolysFct[120, Labeled = True]**

**PolysFct[120, Labeled = True]**

Torus knots have particularly nice-looking  $\Theta$  invariants. Here are the torus knots  $T_{12,7}$ ,  $T_{12,5}$ , and  $T_{12,3}$ .

Figure 4.1: The modified Green function  $g_{\alpha\beta}$  is invariant under Reidemeister moves performed away from where it is measured.

**Lemma 7.** Inserting a null vertex does not change  $g_{\alpha\beta}$  provided it is inserted away from the points  $\alpha$  and  $\beta$ .

**Proof.** Let  $D$  be an upright knot diagram having an edge labelled  $i$  and let  $D'$  be obtained from it by adding a null vertex within edge  $i$ , naming the two resulting half-edges  $j$  and  $k$  (in order). Let  $A$  be the Green function for  $D$ , and similarly  $A'$  for  $D'$ . We claim that

$$g_{\alpha\beta} = g'_{\alpha\beta} \quad \text{if } \alpha \neq j, k \text{ and } \beta \neq j, k$$

Indeed, all we have to do is verify that the above-defined  $g_{\alpha\beta}$  satisfies all the  $g$ -rules (8)-(11), and that is easy. The lemma now follows easily from the definition of  $g$  in Equation (10).

**Remark 8.** The statement of our Main Theorem, Theorem 1, does not change in the presence of null vertices: There are no " $T^*$ " terms for them, and their only effect on the definition of  $\Theta$  is in Equation (5) where the edge labels that appear within  $c_1$ ,  $c_2$ , and  $c_3$  within the  $F_3$  sum.

The following theorem was not named in [BV1] yet it was stated there as the first part of the first proof of [BV1, Theorem 1].

**Theorem 9.** The variant Green function  $g_{\alpha\beta}$  is a "relative invariant", meaning that once points  $\alpha$  and  $\beta$  are fixed within a knot diagram  $D$ , the value of  $g_{\alpha\beta}$  does not change if Reidemeister moves are performed away from the points  $\alpha$  and  $\beta$  (an illustration appears in Figure 4.1). It follows that the same is also true for  $\Theta$  for  $\alpha, \beta \neq i^*, j^*$ .

We note that  $g_{\alpha\beta}$  is nearly the same as  $g_{\alpha\beta}$ , if  $\alpha$  is on  $\alpha$  and  $\beta$  is on  $\beta$ . So Theorem 9 also says that  $g_{\alpha\beta}$  is invariant under Reidemeister moves away from  $\alpha$  and  $\beta$ , except for edge-removing moves and  $\pm 1$  contributions that arise if  $\alpha$  and  $\beta$  correspond to edges that get merged or broken by the Reidemeister moves.

By Theorem 9,  $g_{\alpha\beta}$  is a relative invariant under Reidemeister moves away from  $\alpha$  and  $\beta$ , except for edge-removing moves and  $\pm 1$  contributions that arise if  $\alpha$  and  $\beta$  correspond to edges that get merged or broken by the Reidemeister moves.

This statement does not make sense for  $g_{\alpha\beta}$  inverting a null vertex changes the dimensions of the matrix  $G = (g_{\alpha\beta})$ .

Figure 4.4: The two sides  $IV$  and  $Sw$  of the  $R3$  move. The left side  $IV$  consists of 3 distinguished crossings  $c_1 = (1, i, k)$ ,  $c_2 = (1, i, k)$ ,  $c_3 = (1, i, k)$ , and a collection of further crossings  $c_4 = (s, i, n)$ ,  $c_5 = (s, i, n)$ , where  $s$  is the set of crossings not participating in the  $R3$  move. The right side  $IV$  consists of  $c_1 = (1, i, k)$ ,  $c_2 = (1, i, k)$ ,  $c_3 = (1, i, k)$ , and the same set  $S$  of further crossings.

**Proof.** Let  $D_1$  and  $D_2$  be two knot diagrams that differ only by an  $R3$  move, and label their relevant edges and crossings as in Figure 4.4. Let  $g_{\alpha\beta}$  and  $g'_{\alpha\beta}$  be their corresponding Green functions. Let  $F_1(c)$ ,  $F_2(c)$ , and  $F_3(c)$  be defined from  $g_{\alpha\beta}$  as in (3)-(5), and similarly make  $F'_1(c)$ ,  $F'_2(c)$ , and  $F'_3(c)$  using  $g'_{\alpha\beta}$ .

We claim that  $A' = A$ ,  $B' = B$ , and  $C' = C$ .

To show that  $A' = A$ , we need to rewrite the entries of  $A'$  with polynomials in  $g_{\alpha\beta}$  in terms of  $g'_{\alpha\beta}$  and  $\beta$  may belong to the set  $\{i^*, j^*, k^*\}$  in which it may be that  $g'_{\alpha\beta} \neq g_{\alpha\beta}$ . For the  $g$ -rules of Equations (8) and (9) allow us to rewrite the offending  $g'_{\alpha\beta}$  in terms of  $g_{\alpha\beta}$  with subscripts in  $\{i^*, j^*, k^*\}$  (in terms of other  $g$ -rules whose subscripts are in  $\{i, j, k, n\}$ ), where  $g'_{\alpha\beta} = g_{\alpha\beta}$ . It is enough to show that

$$g'_{\alpha\beta} = g_{\alpha\beta} \quad \text{if } \alpha \neq i^*, j^*, k^* \text{ and } \beta \neq i^*, j^*, k^* \quad (18)$$

where the symbol " $\rightarrow$ " means "apply the rules". This is a finite computation that can in principle be carried out by hand, but each  $A^k$  is a sum of  $3+9=12$  polynomials in the  $g$ 's or the  $g'$ 's, these polynomials are rather unpleasant (see (3) and (4)), and applying the relevant  $g$ -rules is a bit further to the complexity. Luckily, we can delegate this tedious calculation to an entity that works accurately and doesn't complain.

First, we implement the Kronecker  $\delta$ -function, the  $g$ -rules for a crossing  $c = (s, i, j)$ , and the  $g$ -rules for a set of crossings  $c = \{c_1, c_2, c_3\}$ .

**Exp[Exp[27 + 17T, -1 + 2T, -1 + 2T, 1 + 4T + 17T]]**

We then let  $X$  be the crossings in the left-hand-side of the  $R3$  move, as in Figure 4.4, we let  $A$  be the  $A'$  terms of (17), and we let  $B$  be the result of applying the  $g$ -rules for the

**Theorem 1** (The Main Theorem, proof in Section 4). The following are knot invariants:

$$\Theta(K) = \frac{1}{\Delta(K)} \left( F_1(c) + \sum_{c \in X} F_2(c) + \sum_{c \in X} F_3(c) \right) \quad (9)$$

Furthermore,  $\Theta$  is a Laurent polynomial in  $T_1$  and  $T_2$ , with integer coefficients.

Some comments are now in order:

**Comment 2.** The entries of  $G_{\alpha\beta}$  are rational functions with denominators  $\Delta_{\alpha\beta}$ , and so  $\Theta$  is rational in the ring of rational functions  $\mathbb{Q}(T_1, T_2)$ . The point of  $\Theta$  is to clear these denominators by multiplying by  $\Delta_{\alpha\beta} \Delta_{\beta\alpha}$  so as to get an invariant valued in Laurent polynomials in  $T_1$  and  $T_2$ . (There remains a potential denominator of the form  $(T_2 - 1)^2$  coming from the explicit denominators in Equations (3) and (4). It will be shown to cancel in Section 4.2.)

**Comment 3.** We note following [BV1] that  $g_{\alpha\beta}$  can be interpreted as measuring "car traffic": assuming that the diagram  $D$  is a stream of traffic is injected onto the start of edge  $\alpha$  and a "traffic count" is placed near the end of edge  $\beta$ , and where cars always obey the following traffic rules:

- Car travel on the "end" of the knot, always in a direction consistent with the orientation of the knot.
- When a car reaches a crossing of edges  $\alpha = \pm 1$  on the overstrand, it continues right and continues on the other side.
- When a car reaches a crossing of edges  $\alpha = \pm 1</$

