

Dror Bar-Natan: Talks: LesDiablerets-1508:

$\omega\beta := \text{http://www.math.toronto.edu/~drorbn/Talks/LesDiablerets-1508/}$

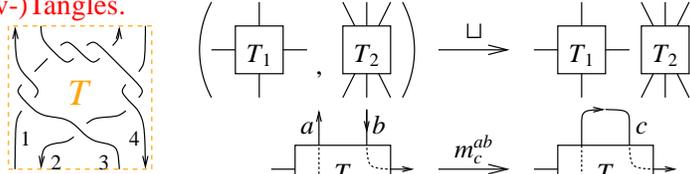
Work in Progress on

# Polynomial Time Knot Polynomials, A

**Abstrant.** The value of things is inversely correlated with their computational complexity. "Real time" machines, such as our brains, only run linear time algorithms, and there's still a lot we don't know. Anything we learn about things doable in linear time is truly valuable. Polynomial time we can in-practice run, even if we have to wait; these things are still valuable. Exponential time we can play with, but just a little, and exponential things must be beautiful or philosophically compelling to deserve attention. Values further diminish and the aesthetic-or-philosophical bar further rises as we go further slower, or un-computable, or ZFC-style intrinsically infinite, or large-cardinalish, or beyond.

I will explain some things I know about polynomial time knot polynomials and explain where there's more, within reach.

## (v-)Tangles.



## Why Tangles?

- Finitely presented. (meta-associativity:  $m_a^{ab} // m_a^{ac} = m_b^{bc} // m_a^{ab}$ )
  - Divide and conquer proofs and computations.
  - "Algebraic Knot Theory": If  $K$  is ribbon,  $z(K) \in \{cl_2(\zeta) : cl_1(\zeta) = 1\}$ .
- (Genus and crossing number are also definable properties).  $cl_1$ : trivial  $cl_2$ : ribbon  $\mathcal{T}_{2n}$   $K \in \mathcal{T}_1$
- A blackboard aside on genus? Faster is better, leaner is meaner!

**Theorem 1.**  $\exists!$  an invariant  $z_0$ : {pure framed  $S$ -component tangles}  $\rightarrow \Gamma_0(S) := R \times M_{S \times S}(R)$ , where  $R = R_S = \mathbb{Z}((T_a)_{a \in S})$  is the ring of rational functions in  $S$  variables, intertwining

$$\left( \begin{array}{c|c} \omega_1 & S_1 \\ \hline S_1 & A_1 \end{array}, \begin{array}{c|c} \omega_2 & S_2 \\ \hline S_2 & A_2 \end{array} \right) \sqcup \rightarrow \begin{array}{c|cc} \omega_1 \omega_2 & S_1 & S_2 \\ \hline S_1 & A_1 & 0 \\ S_2 & 0 & A_2 \end{array}$$

$$\begin{array}{c|ccc} \omega & a & b & S \\ \hline a & \alpha & \beta & \theta \\ b & \gamma & \delta & \epsilon \\ S & \phi & \psi & \Xi \end{array} \xrightarrow{m_c^{ab}} \begin{array}{c|cc} \mu\omega & c & S \\ \hline c & \gamma + \alpha\delta/\mu & \epsilon + \delta\theta/\mu \\ S & \phi + \alpha\psi/\mu & \Xi + \psi\theta/\mu \end{array}$$

and satisfying  $(|a; a \nearrow b, b \nearrow a) \xrightarrow{z_0} \left( \begin{array}{c|c} 1 & a \\ \hline a & 1 \end{array}; \begin{array}{c|cc} 1 & a & b \\ \hline b & 0 & T_a^{\pm 1} \end{array} \right)$

**In Addition** • The matrix part is just a stitching formula for Burau/Gassner [LD, KLW, CT].

- $K \mapsto \omega$  is Alexander, mod units.
- $L \mapsto (\omega, A) \mapsto \omega \det'(A - I)/(1 - T')$  is the MVA, mod units.
- The fastest Alexander algorithm I know.
- There are also formulas for strand deletion, reversal, and doubling.
- Every step along the computation is the invariant of something.
- Extends to and more naturally defined on v/w-tangles.
- Fits in one column, including propaganda & implementation.



M. Polyak & T. Ohtsuki @ Heian Shrine, Kyoto

## Implementation key idea:

```

(\omega, A = (\alpha_{ab})) \leftrightarrow
(\omega, \lambda = \sum \alpha_{ab} t_a h_b)

F := \Gamma[\omega_1, \lambda_1] \Gamma[\omega_2, \lambda_2] := \Gamma[\omega_1 \# \omega_2, \lambda_1 \# \lambda_2];
m_{a,b \to c}[\Gamma[\omega, \lambda]] := Module[(\alpha, \beta, \gamma, \delta, \theta, \epsilon, \phi, \psi, \Xi, \mu),
  (\alpha \beta \theta) := (\partial_{t_a, h_a} \lambda \partial_{t_b, h_b} \lambda \partial_{t_c, \lambda}) / (. (t | h)_{a|b} \to 0;
  (\gamma \delta \epsilon) := (\partial_{t_a, h_a} \lambda \partial_{t_b, h_b} \lambda \partial_{t_c, \lambda}) / (. (t | h)_{a|b} \to 0;
  (\phi \psi \Xi) := (\partial_{h_a, \lambda} \partial_{h_b, \lambda} \lambda) / (. (t | h)_{a|b} \to 0;
  \Gamma[(\mu = 1 - \beta) \omega, \{t_c, 1\}. (\gamma + \alpha\delta/\mu \epsilon + \delta\theta/\mu) (\phi + \alpha\psi/\mu \Xi + \psi\theta/\mu) \cdot (h_c, 1)]
  / (. {T_a \to T_c, T_b \to T_c} // FCollect];
RP_{a,b} := \Gamma[1, \{t_a, t_b\}. (0 \ 1 - T_a^{-1}) \cdot (h_a, h_b)];
RM_{a,b} := RP_{ab} / . T_a \to 1 / T_b;
  
```

## Meta-Associativity

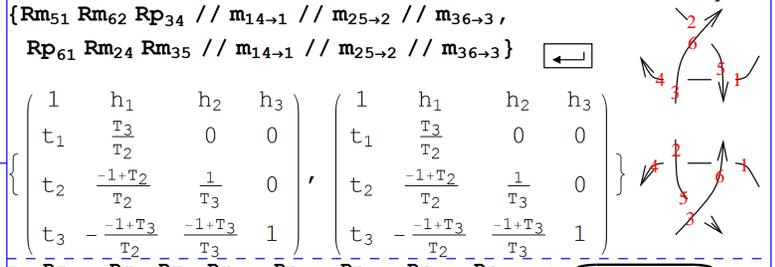
$$\zeta = \Gamma[\omega, \{t_1, t_2, t_3, t_s\}. \begin{pmatrix} \alpha_{11} & \alpha_{12} & \alpha_{13} & \theta_1 \\ \alpha_{21} & \alpha_{22} & \alpha_{23} & \theta_2 \\ \alpha_{31} & \alpha_{32} & \alpha_{33} & \theta_3 \\ \phi_1 & \phi_2 & \phi_3 & \Xi \end{pmatrix} \cdot \{h_1, h_2, h_3, h_s\}];$$

$$(\zeta // m_{12 \to 1} // m_{13 \to 1}) = (\zeta // m_{23 \to 2} // m_{12 \to 1})$$

True

R3

... divide and conquer!



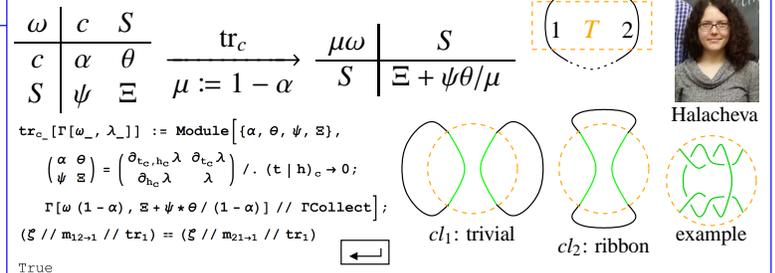
$$z = RM_{12,1} RM_{27} RM_{83} RM_{4,11} RP_{16,5} RP_{6,13} RP_{14,9} RP_{10,15};$$

$$Do[z = z // m_{1k \to 1}, \{k, 2, 16\}];$$

$$z = \begin{pmatrix} 11 - \frac{1}{T_1^3} + \frac{4}{T_1^2} - \frac{8}{T_1} - 8 T_1 + 4 T_1^2 - T_1^3 & h_1 \\ \dots & \dots \\ \dots & 1 \end{pmatrix}$$

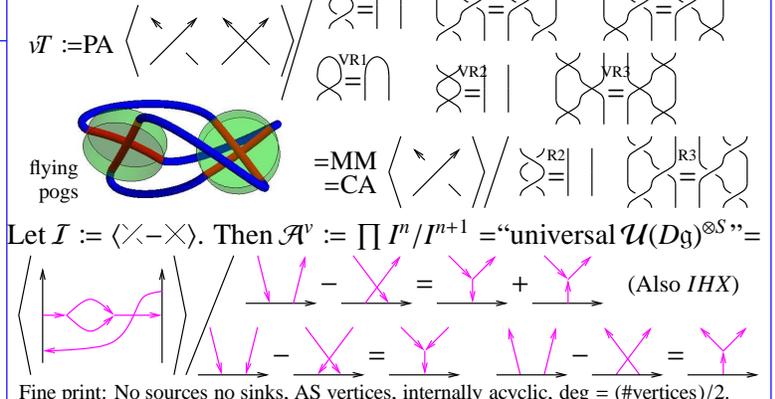
## Closed Components.

The Halacheva trace  $tr_c$  satisfies  $m_c^{ab} // tr_c = m_c^{ba} // tr_c$  and computes the MVA for all links in the atlas, but its domain is not understood:



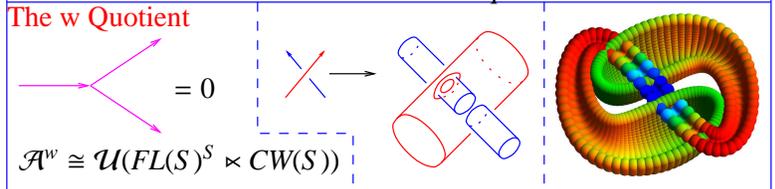
**Weaknesses.** •  $m_c^{ab}$  and  $tr_c$  are non-linear. • The product  $\omega A$  is always Laurent, but my current proof takes induction with exponentially many conditions. • I still don't understand  $tr_c$ , "unitarity", the algebra for ribbon knots. **Where does it come from?**

## v-Tangles.



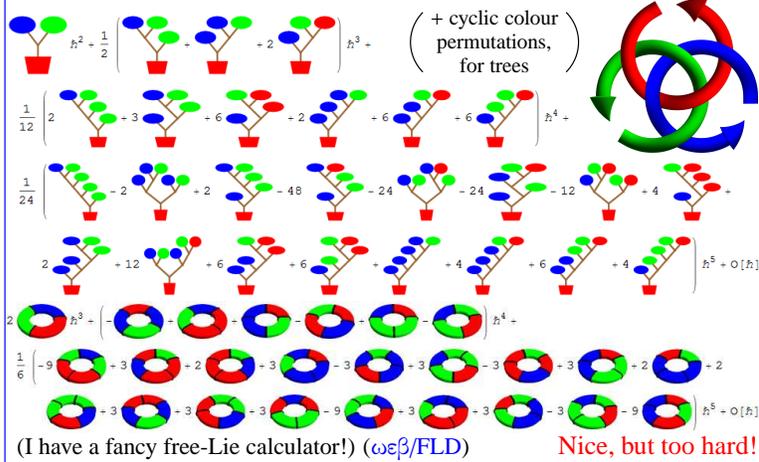
**Likely Theorem.** [EK, En] There exists a homomorphic expansion (universal finite type invariant)  $Z: \mathcal{vT} \rightarrow \mathcal{A}^v$ . (issues suppressed)

## The w Quotient



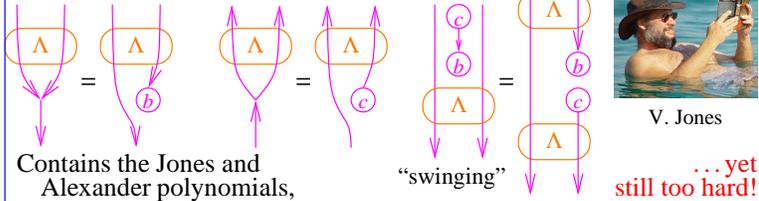
**Theorem 2 [BND].**  $\exists!$  a homomorphic expansion, aka a homomorphic universal finite type invariant  $Z^w$  of pure w-tangles.  $z^w := \log Z^w$  takes values in  $FL(S)^S \times CW(S)$ .

$z$  is computable.  $z$  of the Borromean tangle, to degree 5 [BN]:



**Proposition [BN].** Modulo all relations that universally hold for the 2D non-Abelian Lie algebra and after some changes-of-variable,  $z^w$  reduces to  $z_0$ .

Back to v – the 2D “Jones Quotient”.



**The OneCo Quotient.** Likely related to [ADO]  $= 0$ , only one co-bracket is allowed. Everything should work, and everything is being worked!

**References.**  
 [ADO] Y. Akutsu, T. Deguchi, and T. Ohtsuki, *Invariants of Colored Links*, J. of Knot Theory and its Ramifications **1-2** (1992) 161–184.  
 [BN] D. Bar-Natan, *Balloons and Hoops and their Universal Finite Type Invariant, BF Theory, and an Ultimate Alexander Invariant*,  $\omega\epsilon\beta$ /KBH, arXiv:1308.1721.  
 [BND] D. Bar-Natan and Z. Dancso, *Finite Type Invariants of W-Knotted Objects I-II*,  $\omega\epsilon\beta$ /WKO1,  $\omega\epsilon\beta$ /WKO2, arXiv:1405.1956, arXiv:1405.1955.  
 [BNS] D. Bar-Natan and S. Selmani, *Meta-Monoids, Meta-Bicrossed Products, and the Alexander Polynomial*, J. of Knot Theory and its Ramifications **22-10** (2013), arXiv:1302.5689.  
 [CT] D. Cimasoni and V. Turaev, *A Lagrangian Representation of Tangles*, Topology **44** (2005) 747–767, arXiv:math.GT/0406269.  
 [En] B. Enriquez, *A Cohomological Construction of Quantization Functors of Lie Bialgebras*, Adv. in Math. **197-2** (2005) 430–479, arXiv:math/0212325.  
 [EK] P. Etingof and D. Kazhdan, *Quantization of Lie Bialgebras, I*, Selecta Mathematica **2** (1996) 1–41, arXiv:q-alg/9506005.  
 [GST] R. E. Gompf, M. Scharlemann, and A. Thompson, *Fibered Knots and Potential Counterexamples to the Property 2R and Slice-Ribbon Conjectures*, Geom. and Top. **14** (2010) 2305–2347, arXiv:1103.1601.  
 [KLW] P. Kirk, C. Livingston, and Z. Wang, *The Gassner Representation for String Links*, Comm. Cont. Math. **3** (2001) 87–136, arXiv:math/9806035.  
 [LD] J. Y. Le Dimet, *Enlacements d’Intervalles et Représentation de Gassner*, Comment. Math. Helv. **67** (1992) 306–315.

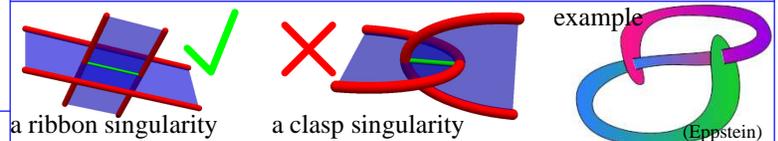
**Definition.** (Compare [BNS, BN]) A **The Abstract Context** meta-monoid is a functor  $M: (\text{finite sets, injections}) \rightarrow (\text{sets})$  (think “ $M(S)$  is quantum  $G^S$ ”, for  $G$  a group) along with natural operations  $*$ :  $M(S_1) \times M(S_2) \rightarrow M(S_1 \sqcup S_2)$  whenever  $S_1 \cap S_2 = \emptyset$  and  $m_c^{ab}: M(S) \rightarrow M((S \setminus \{a, b\}) \sqcup \{c\})$  whenever  $a \neq b \in S$  and  $c \notin S \setminus \{a, b\}$ , such that  
 meta-associativity:  $m_a^{ab} // m_a^{ac} = m_b^{bc} // m_a^{ab}$   
 meta-locality:  $m_c^{ab} // m_f^{de} = m_f^{de} // m_c^{ab}$   
 and, with  $\epsilon_b = M(S \hookrightarrow S \sqcup \{b\})$ ,  
 meta-unit:  $\epsilon_b // m_a^{ab} = Id = \epsilon_b // m_a^{ba}$ .

**Claim.** Pure virtual tangles  $PVT$  form a meta-monoid.  
**Theorem.**  $S \mapsto \Gamma_0(S)$  is a meta-monoid and  $z_0: PVT \rightarrow \Gamma_0$  is a morphism of meta-monoids.  
**Strong Conviction.** There exists an extension of  $\Gamma_0$  to a bigger meta-monoid  $\Gamma_{01}(S) = \Gamma_0(S) \times \Gamma_1(S)$ , along with an extension of  $z_0$  to  $z_{01}: PVT \rightarrow \Gamma_{01}$ , with  
 $\Gamma_1(S) = V \oplus V^{\otimes 2} \oplus V^{\otimes 3} \oplus S^2(V)^{\otimes 2}$  (with  $V := R_S \langle S \rangle$ ).

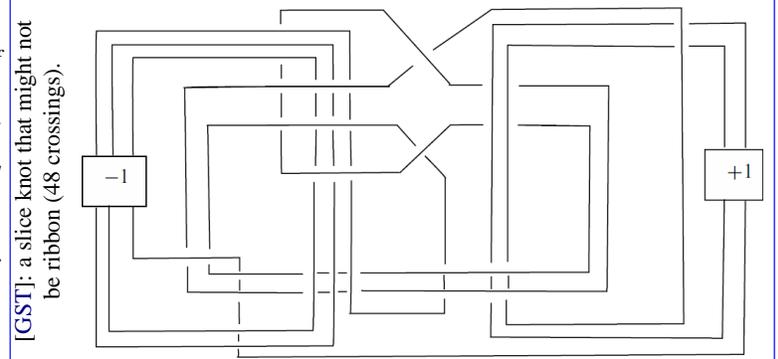
**Furthermore,** upon reducing to a single variable everything is polynomial size and polynomial time.  
**Furthermore,**  $\Gamma_{01}$  is given using a “meta-2-cocycle  $\rho_c^{ab}$  over  $\Gamma_0$ ”: In addition to  $m_c^{ab} \rightarrow m_{0c}^{ab}$ , there are  $R_S$ -linear  $m_{1c}^{ab}: \Gamma_1(S \sqcup \{a, b\}) \rightarrow \Gamma_1(S \sqcup \{c\})$ , a meta-right-action  $\alpha^{ab}: \Gamma_1(S) \times \Gamma_0(S) \rightarrow \Gamma_1(S)$   $R_S$ -linear in the first variable, and a first order differential operator (over  $R_S$ )  $\rho_c^{ab}: \Gamma_0(S \sqcup \{a, b\}) \rightarrow \Gamma_1(S \sqcup \{c\})$  such that

$$(\zeta_0, \zeta_1) // m_c^{ab} = (\zeta_0 // m_{0c}^{ab}, (\zeta_1, \zeta_0) // \alpha^{ab} // m_{1c}^{ab} + \zeta_0 // \rho_c^{ab})$$

**What’s done?** The braid part, with still-ugly formulas.  
**What’s missing?** A lot of concept- and detail-sensitive work towards  $m_{1c}^{ab}$ ,  $\alpha^{ab}$ , and  $\rho_c^{ab}$ . The “ribbon element”.



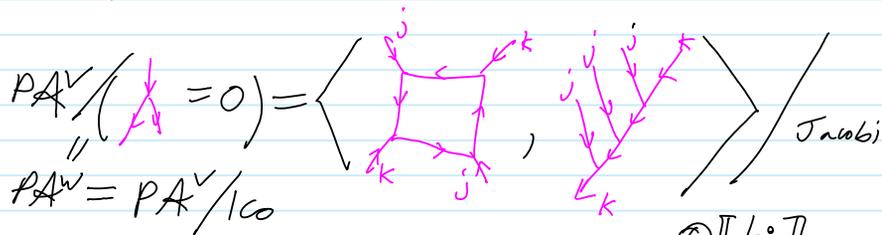
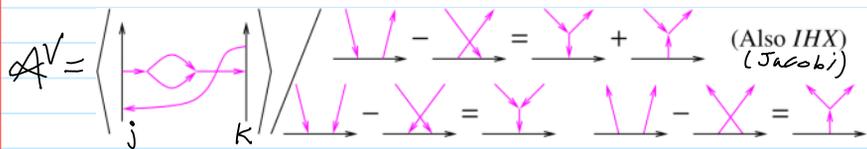
**A bit about ribbon knots.** A “ribbon knot” is a knot that can be presented as the boundary of a disk that has “ribbon singularities”, but no “clasp singularities”. A “slice knot” is a knot in  $S^3 = \partial B^4$  which is the boundary of a non-singular disk in  $B^4$ . Every ribbon knot is clearly slice, yet,  
**Conjecture.** Some slice knots are not ribbon.  
**Fox-Milnor.** The Alexander polynomial of a ribbon knot is always of the form  $A(t) = f(t)f(1/t)$ . (also for slice)



“God created the knots, all else in topology is the work of mortals.”  
 Leopold Kronecker (modified)  
[www.katlas.org](http://www.katlas.org) The Knot Atlas  
 Inverse Can Edit



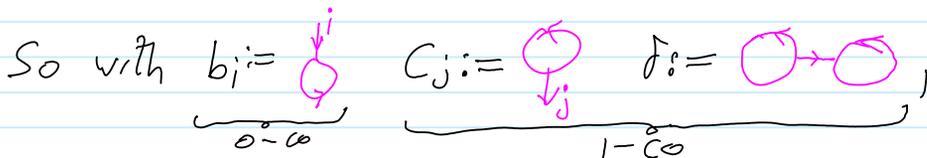
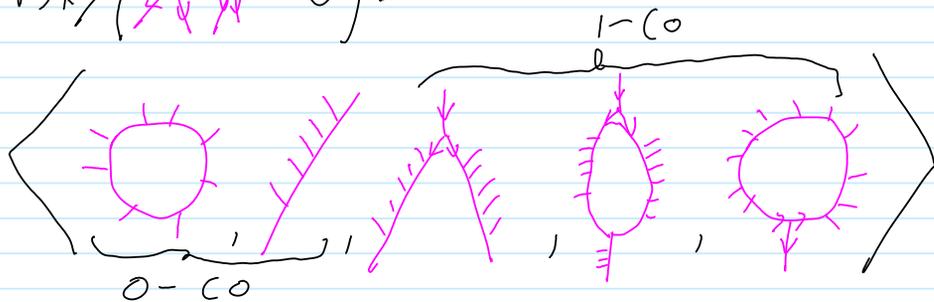
**Help Needed!**  
 I'm slow and feeble-minded.



So  $PA^w(\uparrow_s) / \left( \begin{matrix} i & j \\ \downarrow & \downarrow \\ k & k \end{matrix} = i \downarrow_k^j - j \downarrow_k^i \right) = \hat{R}_s \oplus M_{s \times s}(\hat{R}_s)$

and the rest is (hard!) calculations, which lead to a simple **rational-function** result.

$PA^v / (\text{crossings} = 0) =$



$(PA^v / 2co) / 2D \subset$

$\hat{R}_s \oplus M_{s \times s}(\hat{R}_s) \oplus \hat{R}_s \downarrow_i^j \oplus \hat{R}_s \downarrow_k^j \oplus \hat{R}_s \downarrow_i^j \oplus \hat{R}_s \downarrow_j^i \oplus \hat{R}_s \downarrow_j^k$   
 $= V_s + V_s^{\otimes 2} + V_s + V_s^{\otimes 2} + V_s^{\otimes 3} + (S^2(V_s))^{\otimes 2}$

[The product law is awful, but experience shows that things simplify....]

Stitching is clearly possible, but I still don't have explicit formulas.

Proposition The element  $R_{ij}$  given below solves the YB equation

$R_{12}R_{13}R_{23} = R_{23}R_{13}R_{12}$

in  $A^v / 2co / 2D$ :

$R_{jk} = e^{j-k} e^{\rho}$ , with

$\rho = -\phi_2(b_j) \begin{matrix} j & k \\ | & | \\ \hline c \end{matrix}$

$+ \frac{\phi_2(b_j)}{b_j} \begin{matrix} j & k \\ | & | \\ \hline c \end{matrix}$

$+ \frac{\phi_1(b_j)\phi_2(b_k)}{b_k \phi_1(b_k)} \begin{matrix} j & k \\ | & | \\ \hline c \end{matrix}$

$- \frac{\phi_2(b_j)}{b_j^2} \rho \begin{matrix} j & k \\ | & | \\ \hline c \end{matrix}$

$- \frac{\phi_1(b_j)\phi_2(b_k)}{b_j b_k \phi_1(b_k)} \rho \begin{matrix} j & k \\ | & | \\ \hline c \end{matrix}$

where  $\phi_1(x) = e^{-x} - 1$

and  $\phi_2(x) = \frac{(x+2)e^{-x} - 2 + x}{2x}$

Loading, initializing variables, setting default degree to 6.

Meaningless calculations.

```
(The Mathematica packages FreeLie' and AwCalculus' are at œfβ/WKO4).
path = "C:/drorbn/AcademicPensive/";
SetDirectory[path <> "2015-08/LesDiablerets-1508"];
Get[path <> "Projects/WKO4/FreeLie.m"];
Get[path <> "Projects/WKO4/AwCalculus.m"];
x = LW@"x"; y = LW@"y"; u = LW@"u";
$SeriesShowDegree = 6;
```

```
FreeLie` implements / extends
{*, +, **, $SeriesShowDegree, (<), ∫, =, ad, Ad, adSeries, AllCyclicWords,
AllLyndonWords, AllWords, Arbitrator, ASeries, AW, b, BCH, BooleanSequence,
BracketForm, BS, CC, Crop, cw, CW, CWS, CWSeries, D, Deg, DegreeScale,
DerivationSeries, div, DK, DKS, DKSeries, EulerE, Exp, Inverse, j, J, JA,
LieDerivation, LieMorphism, LieSeries, LS, LW, LyndonFactorization, Morphism,
New, RandomCWSeries, Randomizer, RandomLieSeries, RC, SeriesSolve, Support, t,
tb, TopBracketForm, tr, UndeterminedCoefficients, aMap, Γ, ℓ, Λ, σ, ħ, ←, →}.
```

FreeLie` is in the public domain. Dror Bar-Natan is committed to support it within reason until July 15, 2022. This is version 150814.

```
AwCalculus` implements / extends
{*, **, =, dA, dc, deg, dm, dS, dΔ, dσ, El, Es, hA, hm, hS, hΔ, hσ,
ho, RandomElSeries, RandomEsSeries, tA, tha, tm, tS, tΔ, tσ, Γ, Λ}.
```

AwCalculus` is in the public domain. Dror Bar-Natan is committed to support it within reason until July 15, 2022. This is version 150814.

**BCH[x, y] (\* Can raise degree to 22 \*)**

$$\begin{aligned} &LS \left[ \overline{x+y}, \frac{\overline{xy}}{2}, \frac{1}{12} \overline{xx\overline{xy}} + \frac{1}{12} \overline{x\overline{xy}y}, \frac{1}{24} \overline{xx\overline{xy}y}, \right. \\ &\quad - \frac{1}{720} \overline{xxx\overline{xy}} + \frac{1}{180} \overline{xx\overline{xy}y} + \frac{1}{180} \overline{x\overline{xy}yy} + \frac{1}{120} \overline{x\overline{xy}y} + \\ &\quad \frac{1}{360} \overline{xx\overline{xy}y} - \frac{1}{720} \overline{x\overline{xy}yy}, - \frac{xxx\overline{xy}}{1440} + \frac{1}{360} \overline{xx\overline{xy}y} + \\ &\quad \left. \frac{1}{240} \overline{xx\overline{xy}y} + \frac{1}{720} \overline{xx\overline{xy}y} - \frac{x\overline{xy}yy}{1440}, \dots \right] \end{aligned}$$

**KV Direct.**

{F = LS[{x, y}, Fs], G = LS[{x, y}, Gs]}; Fs["y"] = 1/2;

SeriesSolve[{F, G},

$$\hbar^{-1} (LS[x+y] - BCH[y, x] \equiv F - G - Ad[-x][F] + Ad[y][G]) \wedge$$

$$\begin{aligned} &div_x[F] + div_y[G] \equiv \\ &\frac{1}{2} tr_u [adSeries[\frac{ad}{e_{ad-1}}, x][u] + adSeries[\frac{ad}{e_{ad-1}}, y][u] - \\ &\quad adSeries[\frac{ad}{e_{ad-1}}, BCH[x, y]][u]]; \end{aligned}$$

{F, G} (\* Can raise degree to 13 \*)

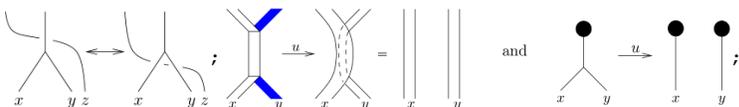
$$\begin{aligned} &LS \left[ \frac{\overline{y}}{2}, \frac{\overline{xy}}{6}, \frac{1}{24} \overline{x\overline{xy}y}, - \frac{1}{180} \overline{xx\overline{xy}} + \frac{1}{80} \overline{x\overline{xy}y} + \frac{1}{360} \overline{x\overline{xy}y}, \right. \\ &\quad - \frac{1}{720} \overline{xx\overline{xy}y} + \frac{1}{240} \overline{x\overline{xy}yy} + \frac{1}{240} \overline{x\overline{xy}y} + \frac{1}{720} \overline{xx\overline{xy}y} - \\ &\quad \frac{\overline{xy}yy}{1440}, \frac{xxx\overline{xy}}{5040} - \frac{xxx\overline{xy}}{1344} + \frac{13xx\overline{xy}y}{15120} + \frac{1}{840} \overline{xx\overline{xy}y} + \\ &\quad \frac{xx\overline{xy}y}{3360} + \frac{xx\overline{xy}yy}{6720} + \frac{\overline{xy}yy}{1260} + \frac{\overline{xy}yy}{1680} - \frac{\overline{xy}yy}{10080}, \dots \left. \right], \\ &LS \left[ 0, \frac{\overline{xy}}{12}, \frac{1}{24} \overline{x\overline{xy}y}, - \frac{1}{360} \overline{xx\overline{xy}} + \frac{1}{120} \overline{x\overline{xy}y} + \frac{1}{180} \overline{x\overline{xy}y}, \right. \\ &\quad - \frac{1}{720} \overline{xx\overline{xy}y} + \frac{1}{240} \overline{x\overline{xy}yy} + \frac{1}{240} \overline{x\overline{xy}y} + \frac{1}{720} \overline{xx\overline{xy}y} - \\ &\quad \frac{\overline{xy}yy}{1440}, \frac{xxx\overline{xy}}{10080} - \frac{xxx\overline{xy}}{2016} + \frac{xx\overline{xy}y}{1890} + \frac{xx\overline{xy}y}{1120} + \frac{xx\overline{xy}y}{5040} + \\ &\quad \left. \frac{x\overline{xy}yy}{2520} + \frac{1}{840} \overline{x\overline{xy}y} + \frac{x\overline{xy}y}{1260} - \frac{\overline{xy}yy}{5040}, \dots \right] \end{aligned}$$

{b[F, G], tr\_x[F]}

$$\begin{aligned} &LS \left[ 0, 0, - \frac{1}{24} \overline{x\overline{xy}y}, - \frac{1}{48} \overline{x\overline{xy}yy}, \frac{1}{720} \overline{xx\overline{xy}y} - \frac{1}{240} \overline{x\overline{xy}yy} - \right. \\ &\quad \frac{\overline{xy}yy}{1440} - \frac{1}{720} \overline{xx\overline{xy}y} - \frac{1}{360} \overline{x\overline{xy}yy}, \frac{xx\overline{xy}y}{1440} - \\ &\quad \left. \frac{1}{480} \overline{xx\overline{xy}yy} - \frac{1}{288} \overline{x\overline{xy}yy} - \frac{7xx\overline{xy}y}{2880} + \frac{\overline{xy}yy}{2880}, \dots \right], \\ &CWS \left[ - \frac{\overline{y}}{6}, \frac{\overline{xy}}{24}, \frac{\overline{xy}}{180} + \frac{\overline{xy}}{80} - \frac{\overline{xy}}{360}, - \frac{\overline{xy}}{180} + \frac{\overline{xy}}{240} - \frac{\overline{xy}}{1440}, \right. \\ &\quad - \frac{\overline{xy}}{5040} + \frac{\overline{xy}}{6720} - \frac{\overline{xy}}{1120} + \frac{2\overline{xy}}{945} - \frac{\overline{xy}}{336} + \frac{\overline{xy}}{6720} + \frac{\overline{xy}}{10080}, \\ &\quad \frac{\overline{xy}}{3360} - \frac{\overline{xy}}{1344} - \frac{\overline{xy}}{2240} + \frac{\overline{xy}}{2016} + \frac{13\overline{xy}}{10080} + \frac{\overline{xy}}{1680} - \\ &\quad \left. \frac{\overline{xy}}{3780} - \frac{\overline{xy}}{840} + \frac{\overline{xy}}{5040} + \frac{\overline{xy}}{2240} + \frac{\overline{xy}}{6720} + \frac{\overline{xy}}{60480}, \dots \right] \end{aligned}$$

(Also implemented: ∂λ and derivations in general, tb, e<sup>∂λ</sup> and morphisms in general, div, j, Drinfel'd-Kohno, etc.)

The [BND] "vertex" equations.



$$\begin{aligned} &\alpha = LS[\{x, y\}, \alpha s]; \beta = LS[\{x, y\}, \beta s]; \\ &\gamma = CWS[\{x, y\}, \gamma s]; \\ &V = Es[\langle x \rightarrow \alpha, y \rightarrow \beta \rangle, \gamma]; \\ &\kappa = CWS[\{x\}, \kappa s]; Cap = Es[\langle x \rightarrow LS[0], \kappa \rangle]; \\ &Rs[a_-, b_-] := Es[\langle a \rightarrow LS[0], b \rightarrow LS[LW@a] \rangle, CWS[0]]; \\ &R4Eqn = V ** (Rs[x, z] // dΔ[x, x, y]) ≡ Rs[y, z] ** Rs[x, z] ** V; \\ &UnitarityEqn = \\ &\quad (V ** (V // dA) ≡ Es[\langle x \rightarrow LS[0], y \rightarrow LS[0] \rangle, CWS[0]]); \\ &CapEqn = ((V ** (Cap // dΔ[x, x, y]) // dc[x] // dc[y]) ≡ \\ &\quad (Cap (Cap // dσ[x, y]) // dc[x] // dc[y])); \\ &\beta s["x"] = 1/2; \beta s["y"] = 0; \\ &SeriesSolve[\{\alpha, \beta, \gamma, \kappa\}, \\ &\quad (\hbar^{-1} R4Eqn) \wedge UnitarityEqn \wedge CapEqn]; \\ &\{V, \kappa\} \end{aligned}$$

SeriesSolve::ArbitrarilySetting: In degree 1 arbitrarily setting {κs[x] → 0}.  
SeriesSolve::ArbitrarilySetting: In degree 3 arbitrarily setting {αs[x, y] → 0}.  
SeriesSolve::ArbitrarilySetting: In degree 5 arbitrarily setting {αs[x, x, y] → 0}.  
General::stop:  
Further output of SeriesSolve::ArbitrarilySetting will be suppressed during this calculation. >>

$$\begin{aligned} &Es \left[ \left( \overline{x} \rightarrow LS \left[ 0, - \frac{\overline{xy}}{24}, 0, \frac{7xx\overline{xy}}{5760} - \frac{7x\overline{xy}y}{5760} + \frac{\overline{xy}yy}{1440}, 0, \right. \right. \right. \\ &\quad - \frac{31xxx\overline{xy}}{967680} + \frac{31xxx\overline{xy}}{483840} - \frac{83xx\overline{xy}y}{967680} - \frac{31x\overline{xy}yy}{725760} - \frac{31xx\overline{xy}y}{645120} + \\ &\quad \left. \frac{13x\overline{xy}yy}{241920} + \frac{101\overline{xy}yy}{1451520} + \frac{527xx\overline{xy}y}{5806080} - \frac{\overline{xy}yy}{60480}, \dots \right), \\ &\quad \overline{y} \rightarrow LS \left[ \frac{\overline{xy}}{2}, - \frac{\overline{xy}}{12}, 0, \frac{xx\overline{xy}}{5760} - \frac{1}{720} \overline{x\overline{xy}y} + \frac{1}{720} \overline{x\overline{xy}y}, - \frac{xx\overline{xy}}{7680} + \right. \\ &\quad \frac{xx\overline{xy}}{3840} - \frac{xx\overline{xy}}{6912} - \frac{xxx\overline{xy}}{645120} + \frac{23xx\overline{xy}}{483840} - \frac{13xx\overline{xy}}{161280} - \frac{xx\overline{xy}}{22680} - \\ &\quad \left. \frac{41xx\overline{xy}}{580608} + \frac{xx\overline{xy}}{15120} + \frac{\overline{xy}yy}{12096} + \frac{71xx\overline{xy}}{483840} - \frac{\overline{xy}yy}{30240}, \dots \right], \\ &CWS \left[ 0, - \frac{\overline{xy}}{48}, 0, \frac{\overline{xy}}{2880} + \frac{\overline{xy}}{2880} + \frac{\overline{xy}}{5760} + \frac{\overline{xy}}{2880}, 0, \right. \\ &\quad - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \\ &\quad \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{362880} - \frac{\overline{xy}}{120960} - \frac{\overline{xy}}{241920} - \frac{\overline{xy}}{120960}, \dots \left. \right], \\ &CWS \left[ 0, - \frac{\overline{xy}}{96}, 0, \frac{\overline{xy}}{11520}, 0, - \frac{\overline{xy}}{725760}, \dots \right] \end{aligned}$$





# The Main Course

$B^{(m)} = (\text{PaB}^{(m)}, \mathbf{S} : \text{PaB}^{(m)} \rightarrow \text{PaP}, d_i, s_i, \square, \sigma) :$

same-skeleton linear combinations allowed

$d_0$  (crossing) = (strand) ;  $d_3$  (crossing) = (strand)

$d_2$  (strand) = (crossing) ;  $s_2$  (crossing) = (crossing)

$a$  = (strand) ;  $\sigma$  = (crossing)

$\square$  (box A) = (box B)

and  $\square$  (crossing) = (crossing)

$\square$  (strand) = (strand)

$C^{(m)} = (\text{PaCD}^{(m)}, \mathbf{S} : \text{PaCD}^{(m)} \rightarrow \text{PaP}, d_i, s_i, \square, \tilde{R}) :$

same-skeleton linear combinations allowed

$d_2$  (crossing) = (strand) = (strand) + (strand) + (strand) + (strand)

$d_0$  (strand) = (strand) ;  $s_1$  (strand) = (strand) ;  $s_1$  (crossing) = 0

$a$  = (strand) ;  $X$  = (crossing) ;  $H$  = (strand) ;  $\tilde{R} = X \exp \frac{H}{2}$

$\square$  (local) +  $\square$  (global) =  $\square$  (A)  $\square$  (B) =  $\square$  (A)  $\square$  (B)

$\square$  (A)  $\square$  (B)  $\square$  (C) =  $\square$  (A)  $\square$  (B)  $\square$  (C)

$\square$  (crossing) =  $\square$  (crossing)

$\square$  (strand) =  $\square$  (strand)

$\square$  (crossing) =  $\square$  (strand) +  $\square$  (strand)

ASSO

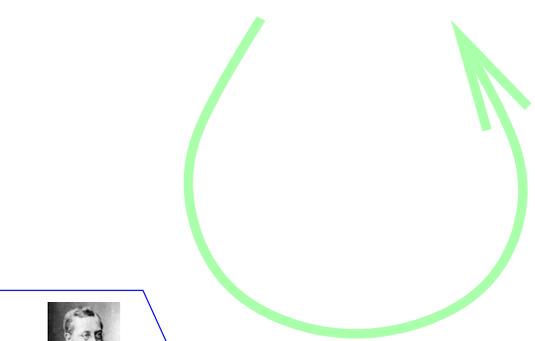
$d_4 \Phi \cdot d_2 \Phi \cdot d_0 \Phi = d_1 \Phi \cdot d_3 \Phi$

$d_1 \exp \left( \pm \frac{1}{2} t^{12} \right) =$

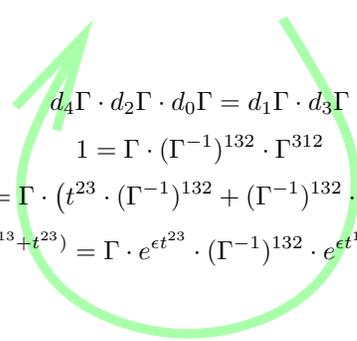
$\Phi \cdot \exp \left( \pm \frac{1}{2} t^{23} \right) \cdot (\Phi^{-1})^{132} \cdot \exp \left( \pm \frac{1}{2} t^{13} \right) \cdot \Phi^{312}$

$s_1 \Phi = s_2 \Phi = s_3 \Phi = 1$

$\square \Phi = \Phi \otimes \Phi$



PaP



$d_4 \Gamma \cdot d_2 \Gamma \cdot d_0 \Gamma = d_1 \Gamma \cdot d_3 \Gamma$

$1 = \Gamma \cdot (\Gamma^{-1})^{132} \cdot \Gamma^{312}$

$d_1 t^{12} = \Gamma \cdot (t^{23} \cdot (\Gamma^{-1})^{132} + (\Gamma^{-1})^{132} \cdot t^{13}) \cdot \Gamma^{312}$

$e^{\epsilon(t^{13} + t^{23})} = \Gamma \cdot e^{\epsilon t^{23}} \cdot (\Gamma^{-1})^{132} \cdot e^{\epsilon t^{13}} \cdot \Gamma^{312}$

GT

GRT

I have a nifty  
Free Lie calculator.  
Wanna play?

**Abstract.** I will describe a **semi-rigorous** reduction of perturbative BF theory (Cattaneo-Rossi [CR]) to computable combinatorics, in the case of ribbon 2-links. Also, I will explain how and why my approach may or may not work in the non-ribbon case. **Weak** this result is, and at least partially already known (Watanabe [Wa]). Yet in the ribbon case, the resulting invariant is a universal finite type invariant, a gadget that significantly generalizes and clarifies the Alexander polynomial and that is closely related to the Kashiwara-Vergne problem. I cannot rule out the possibility that the corresponding gadget in the non-ribbon case will be as interesting.

**The BF Feynman Rules.** For an edge  $e$ , let  $\Phi_e$  be its direction, in  $S^3$  or  $S^1$ . Let  $\omega_3$  and  $\omega_1$  be volume forms on  $S^3$  and  $S^1$ . Then for a 2-link  $(f_i)_{i \in T}$ ,



$$\zeta = \log \sum_{\text{diagrams } D} \frac{D}{|\text{Aut}(D)|} \int_{\mathbb{R}^2} \cdots \int_{\mathbb{R}^2} \int_{\mathbb{R}^4} \cdots \int_{\mathbb{R}^4} \prod_{e \in D} \Phi_e^* \omega_3 \prod_{e \in D} \Phi_e^* \omega_1$$

is an invariant in  $CW(FL(T)) \rightarrow CW(T)$ , "cyclic words in  $T$ ".

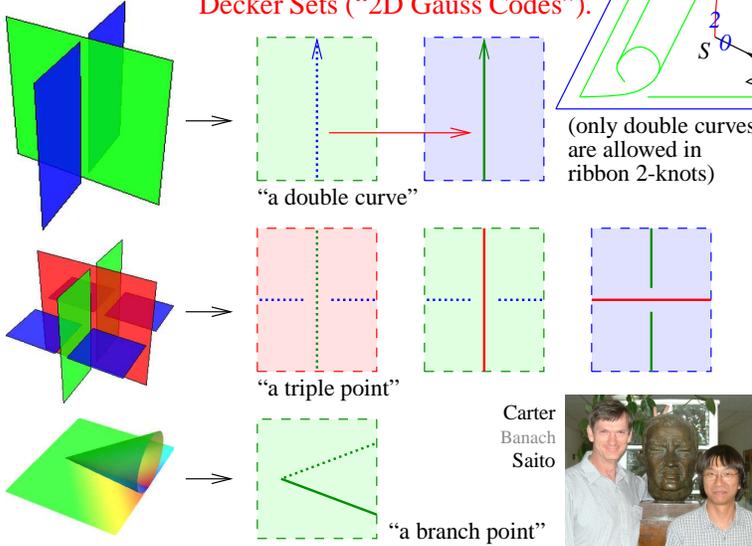
**BF Following [CR].**  $A \in \Omega^1(M = \mathbb{R}^4, \mathfrak{g})$ ,  $B \in \Omega^2(M, \mathfrak{g}^*)$ ,

$$S(A, B) := \int_M \langle B, F_A \rangle.$$

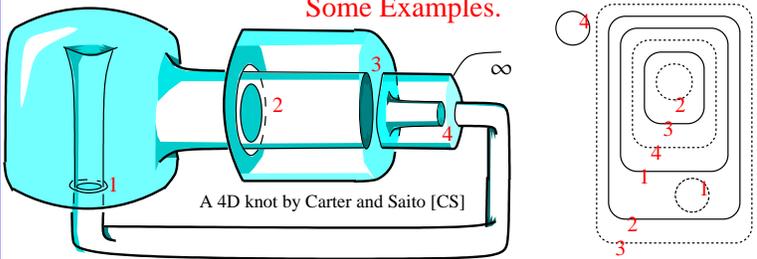
With  $f: (S = \mathbb{R}^2) \rightarrow M$ ,  $\xi \in \Omega^0(S, \mathfrak{g})$ ,  $\beta \in \Omega^1(S, \mathfrak{g}^*)$ , set

$$O(A, B, f) := \int \mathcal{D}\xi \mathcal{D}\beta \exp\left(\frac{i}{\hbar} \int_S \langle \xi, d_{f^*A}\beta + f^*B \rangle\right).$$

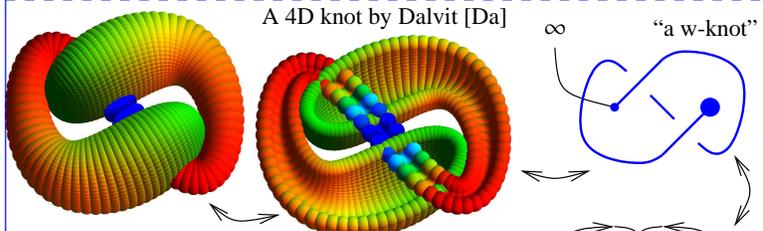
**Decker Sets ("2D Gauss Codes").**



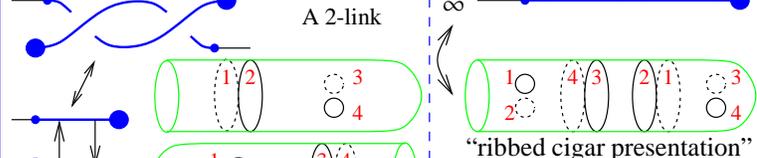
**Some Examples.**



A 4D knot by Dalvit [Da]

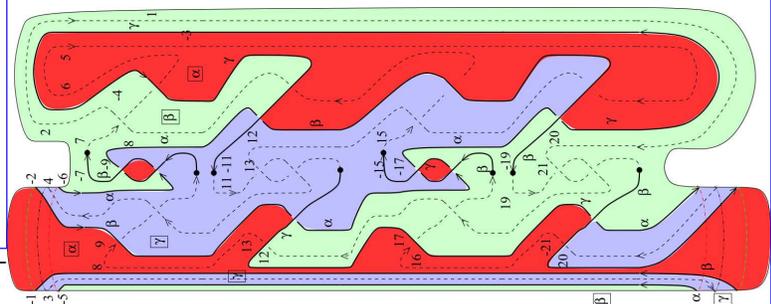
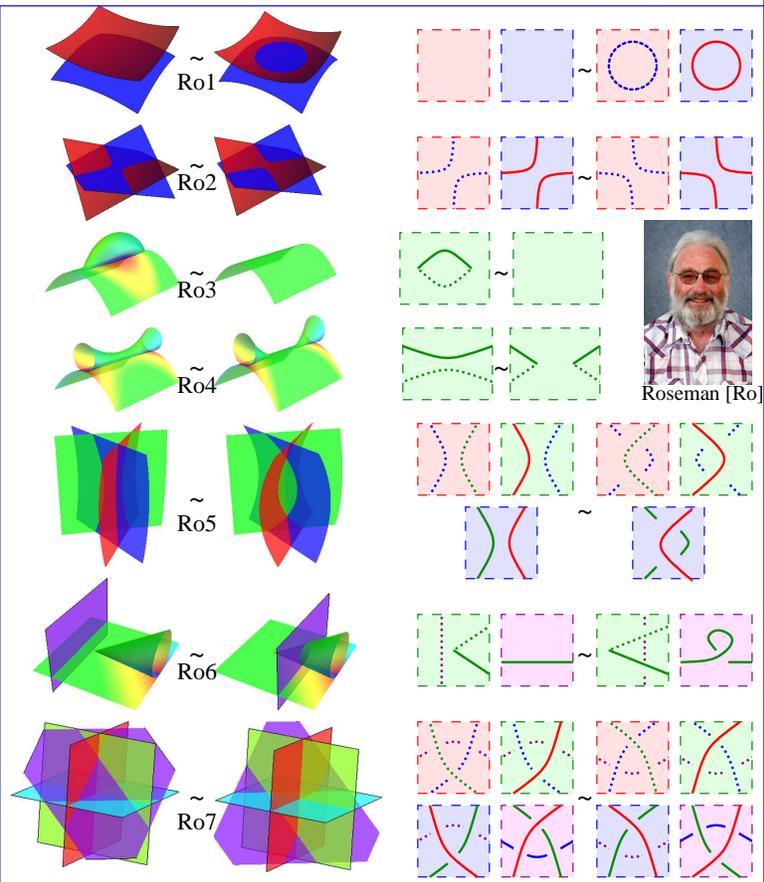
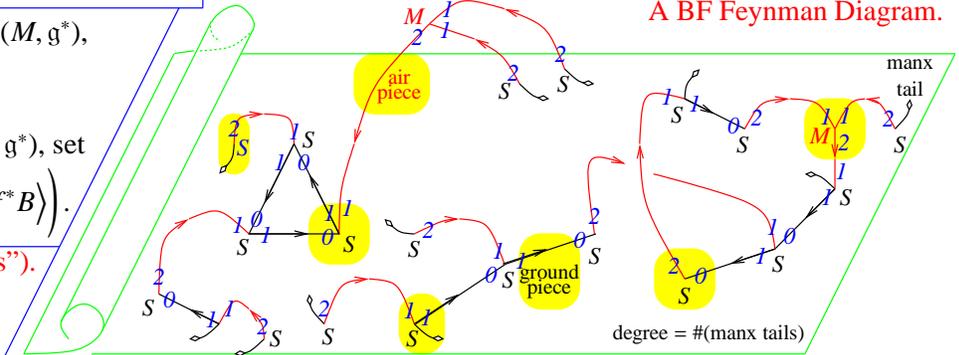


A 2-link



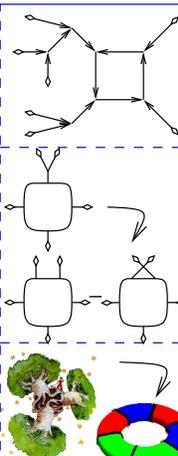
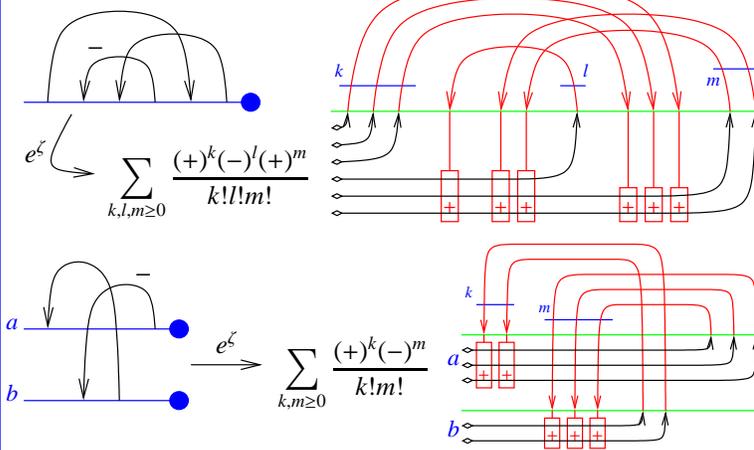
A 2-twist spun trefoil by Carter-Kamada-Saito [CKS].

**A BF Feynman Diagram.**

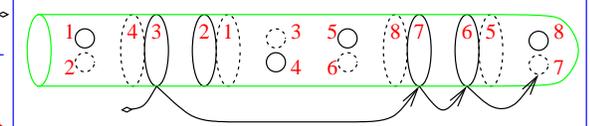


## A Partial Reduction of BF Theory to Combinatorics, 2

**Theorem 1 (with Cattaneo (credit, no blame)).** In the ribbon case,  $e^\zeta$  can be computed as follows:



**Sketch of Proof.** In 4D axial gauge, only “drop down” red propagators, hence in the ribbon case, no  $M$ -trivalent vertices.  $S$  integrals are  $\pm 1$  iff “ground pieces” run on nested curves as below, and exponentials arise when several propagators compete for the same double curve. And then the combinatorics is obvious...



**Theorem 2.** Using Gauss diagrams to represent knots and  $T$ -component pure tangles, the above formulas define an invariant in  $CW(FL(T)) \rightarrow CW(T)$ , “cyclic words in  $T$ ”.

- Agrees with BN-Dancso [BND] and with [BN2].
- In-practice computable!
- Vanishes on braids.
- Extends to w.
- Contains Alexander.
- The “missing factor” in Levine’s factorization [Le] (the rest of [Le] also fits, hence contains the MVA).
- Related to / extends Farber’s [Fa]?
- Should be summed and categorified.

**References.**

[Ar] V. I. Arnold, *Topological Invariants of Plane Curves and Caustics*, University Lecture Series 5, American Mathematical Society 1994.

[BN1] D. Bar-Natan, *Bracelets and the Goussarov filtration of the space of knots, Invariants of knots and 3-manifolds (Kyoto 2001)*, Geometry and Topology Monographs 4 1–12, arXiv:math.GT/0111267.

[BN2] D. Bar-Natan, *Balloons and Hoops and their Universal Finite Type Invariant, BF Theory, and an Ultimate Alexander Invariant*, <http://www.math.toronto.edu/~drorbn/papers/KBH/>, arXiv:1308.1721.

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[Da] E. Dalvit, <http://science.unitn.it/~dalvit/>.

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[Fa] M. Farber, *Noncommutative Rational Functions and Boundary Links*, *Math. Ann.* **293** (1992) 543–568.

[Le] J. Levine, *A Factorization of the Conway Polynomial*, *Comment. Math. Helv.* **74** (1999) 27–53, arXiv:q-alg/9711007.

[Ro] D. Roseman, *Reidemeister-Type Moves for Surfaces in Four-Dimensional Space*, *Knot Theory*, Banach Center Publications **42** (1998) 347–380.

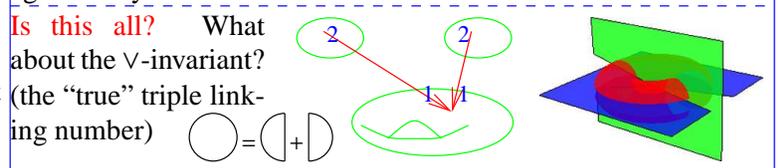
[Wa] T. Watanabe, *Configuration Space Integrals for Long  $n$ -Knots, the Alexander Polynomial and Knot Space Cohomology*, *Alg. and Geom. Top.* **7** (2007) 47–92, arXiv:math/0609742.

**Continuing Joost Slingerland...**

<http://youtu.be/YCA0VIEvHge>

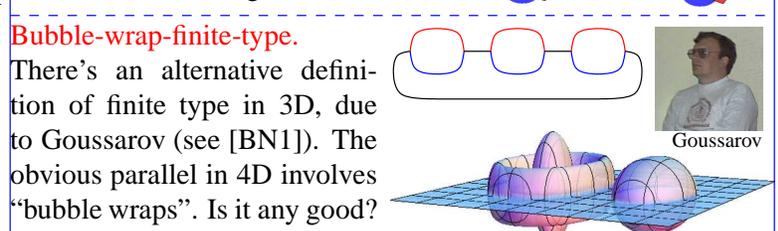
<http://youtu.be/mHyT0cf990>

**Chem-Simons.** When the domain of BF is restricted to ribbon knots, and the target of CS is restricted to trees and wheels, they agree. Why?



**Gnots.** In 3D, a generic immersion of  $S^1$  is an embedding, a knot. In 4D, a generic immersion of a surface has finitely-many double points (a gnot?). Perhaps we should be studying these?

**Finite type.** What are finite-type invariants for 2-knots? What would be “chord diagrams”?



**Shielded tangles.** In 3D, one can’t zoom in and compute “the Chern-Simons invariant of a tangle”. Yet there are well-defined invariants of “shielded tangles”, and rules for their compositions. What would the 4D analog be?



Will the relationship with the Kashiwara-Vergne problem [BND] necessarily arise here?

**Plane curves.** Shouldn’t we understand integral / finite type invariants of plane curves, in the style of Arnold’s  $J^+$ ,  $J^-$ , and  $St$  [Ar], a bit better?

	$a(\times)$	$a(\times)$	$a(\times)$	$\infty$	$\bigcirc$	$\bigcirc$	$\bigcirc$	$\bigcirc$	$\dots$
St	1	0	0	0	0	1	2	3	$\dots$
$J^+$	0	2	0	0	0	-2	-4	-6	$\dots$
$J^-$	0	0	-2	-1	0	-3	-6	-9	$\dots$

“God created the knots, all else in topology is the work of mortals.”

Leopold Kronecker (modified)

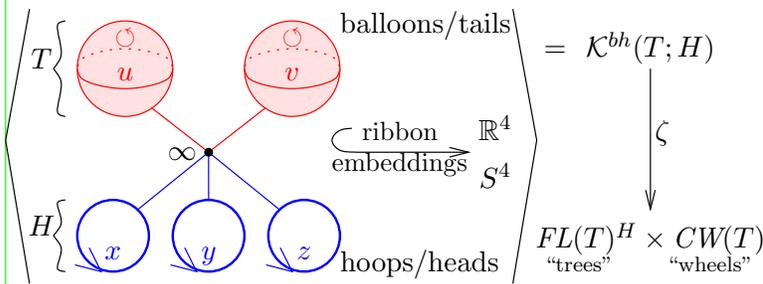
[www.katlas.org](http://www.katlas.org)

The Knot Atlas  
Inventor: Carl Tait

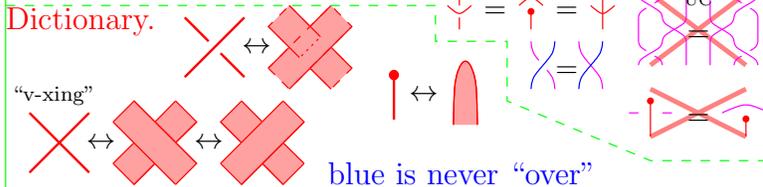
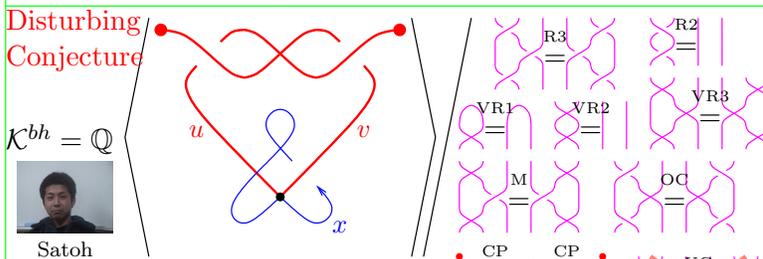


# Finite Type Invariants of Ribbon Knotted Balloons and Hoops

**Abstract.** On my September 17 Geneva talk ( $\omega/\text{sep}$ ) I described a certain trees-and-wheels-valued invariant  $\zeta$  of ribbon knotted loops and 2-spheres in 4-space, and my October 8 Geneva talk ( $\omega/\text{oct}$ ) describes its reduction to the Alexander polynomial. Today I will explain how that same invariant arises completely naturally within the theory of finite type invariants of ribbon knotted loops and 2-spheres in 4-space.



My goal is to tell you why such an invariant is expected, yet not to derive the computable formulas.



**Expansions**  
the semi-virtual  $\otimes := \diagup - \diagdown$  i.e.  $\diagup - \diagdown$  or  $\diagdown - \diagup$

Let  $\mathcal{I}^n := \langle \text{pictures with } \geq n \text{ semi-virts} \rangle \subset \mathcal{K}^{bh}$ .  
We seek an "expansion"

$$Z: \mathcal{K}^{bh} \rightarrow \text{gr } \mathcal{K}^{bh} = \widehat{\bigoplus} \mathcal{I}^n / \mathcal{I}^{n+1} =: \mathcal{A}^{bh}$$

satisfying "property U": if  $\gamma \in \mathcal{I}^n$ , then

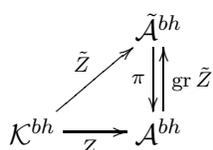
$$Z(\gamma) = (0, \dots, 0, \gamma / \mathcal{I}^{n+1}, *, *, \dots).$$



X.-S. Lin

**Why?** • Just because, and this is vastly more general.  
•  $(\mathcal{K}^{bh} / \mathcal{I}^{n+1})^*$  is "finite-type/polynomial invariants".  
• The Taylor example: Take  $\mathcal{K} = C^\infty(\mathbb{R}^n)$ ,  $\mathcal{I} = \{f \in \mathcal{K} : f(0) = 0\}$ . Then  $\mathcal{I}^n = \{f : f \text{ vanishes like } |x|^n\}$  so  $\mathcal{I}^n / \mathcal{I}^{n+1}$  is homogeneous polynomials of degree  $n$  and  $Z$  is a "Taylor expansion"! (So Taylor expansions are vastly more general than you'd think).

**Plan.** We'll construct a graded  $\tilde{\mathcal{A}}^{bh}$ , a surjective graded  $\pi: \tilde{\mathcal{A}}^{bh} \rightarrow \mathcal{A}^{bh}$ , and a filtered  $\tilde{Z}: \mathcal{K}^{bh} \rightarrow \tilde{\mathcal{A}}^{bh}$  so that  $\pi \circ \text{gr } \tilde{Z} = \text{Id}$  (property U: if  $\text{deg } D = n$ ,  $\tilde{Z}(\pi(D)) = \pi(D) + (\text{deg } \geq n)$ ). Hence •  $\pi$  is an isomorphism. •  $Z := \tilde{Z} \circ \pi$  is an expansion.



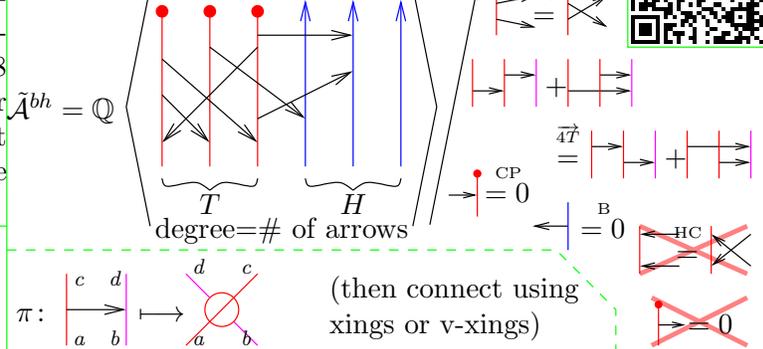
"God created the knots, all else in topology is the work of mortals."

Leopold Kronecker (modified)

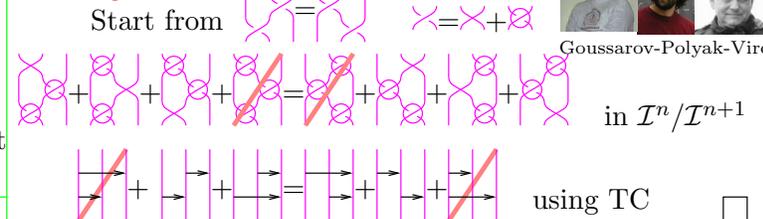
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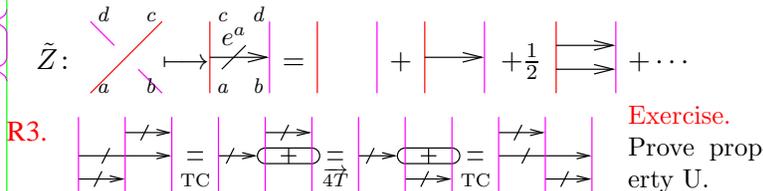
## Action 1.



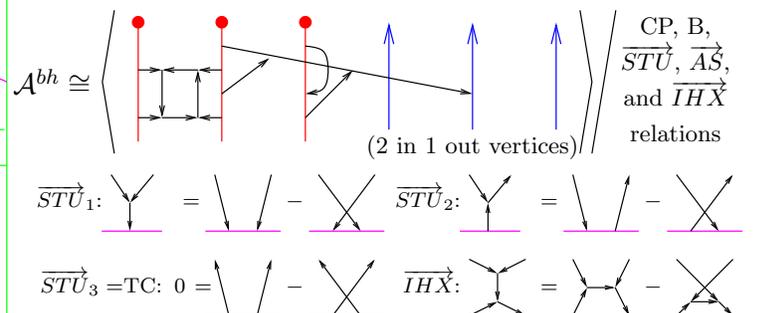
## Deriving 4T.



## Action 2.



## The Bracket-Rise Theorem.



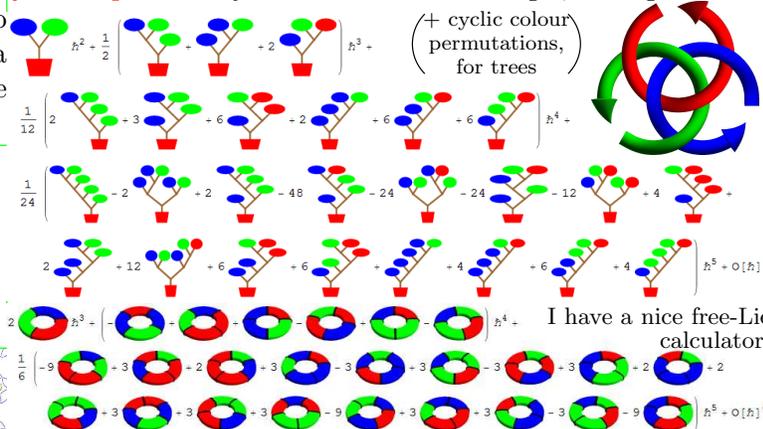
## Proof.



**Corollaries.** (1) Related to Lie algebras! (2) Only trees and wheels persist.

**Theorem.**  $\mathcal{A}^{bh}$  is a bi-algebra. The space of its primitives is  $FL(T)^H \times CW(T)$ , and  $\zeta = \log Z$ .

$\zeta$  is computable!  $\zeta$  of the Borromean tangle, to degree 5:



I have a nice free-Lie calculator!



# The Kashiwara-Vergne Problem and Topology

**Abstract.** I will describe a general machine, a close cousin of Taylor's theorem, whose inputs are topics in topology and whose outputs are problems in algebra. There are many inputs the machine can take, and many outputs it produces, but I will concentrate on just one input/output pair. When fed with a certain class of knotted 2-dimensional objects in 4-dimensional space, it outputs the Kashiwara-Vergne Problem (1978  $\omega/KV$ , solved Alekseev-Meinrenken 2006  $\omega/AM$ , elucidated Alekseev-Torossian 2008-2012  $\omega/AT$ ), a problem about convolutions on Lie groups and Lie algebras.

**The Kashiwara-Vergne Conjecture.** There exist two series  $F$  and  $G$  in the completed free Lie algebra  $FL$  in generators  $x$  and  $y$  so that  $x+y - \log e^y e^x = (1 - e^{-\text{ad } x})F + (e^{\text{ad } y} - 1)G$  in  $FL$  and so that with  $z = \log e^x e^y$ ,

$$\text{tr}(\text{ad } x)\partial_x F + \text{tr}(\text{ad } y)\partial_y G \text{ in cyclic words} \\ = \frac{1}{2} \text{tr} \left( \frac{\text{ad } x}{e^{\text{ad } x} - 1} + \frac{\text{ad } y}{e^{\text{ad } y} - 1} - \frac{\text{ad } z}{e^{\text{ad } z} - 1} - 1 \right)$$

Implies the loosely-stated **convolutions statement**: Convolutions of invariant functions on a Lie group agree with convolutions of invariant functions on its Lie algebra.

**The Machine.** Let  $G$  be a group,  $\mathcal{K} = \mathbb{Q}G = \{\sum a_i g_i : a_i \in \mathbb{Q}, g_i \in G\}$  its group-ring,  $\mathcal{I} = \{\sum a_i g_i : \sum a_i = 0\} \subset \mathcal{K}$  its augmentation ideal. Let

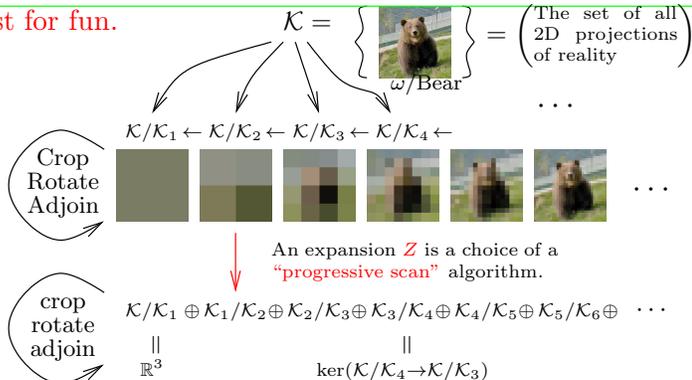
$$\mathcal{A} = \text{gr } \mathcal{K} := \bigoplus_{m \geq 0} \mathcal{I}^m / \mathcal{I}^{m+1}.$$

Note that  $\mathcal{A}$  inherits a product from  $G$ .

**Definition.** A linear  $Z: \mathcal{K} \rightarrow \mathcal{A}$  is an "expansion" if for any  $\gamma \in \mathcal{I}^m$ ,  $Z(\gamma) = (0, \dots, 0, \gamma / \mathcal{I}^{m+1}, *, \dots)$ , and a "homomorphic expansion" if in addition it preserves the product.

**Example.** Let  $\mathcal{K} = C^\infty(\mathbb{R}^n)$  and  $\mathcal{I} = \{f : f(0) = 0\}$ . Then  $\mathcal{I}^m = \{f : f \text{ vanishes like } |x|^m\}$  so  $\mathcal{I}^m / \mathcal{I}^{m+1}$  degree  $m$  homogeneous polynomials and  $\mathcal{A} = \{\text{power series}\}$ . The Taylor series is a homomorphic expansion!

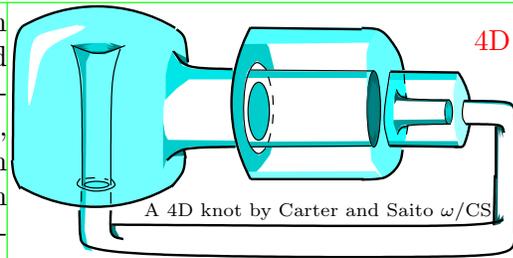
Just for fun.



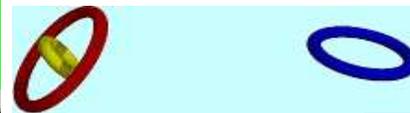
In the finitely presented case, finding  $Z$  amounts to solving a system of equations in a graded space.

**Theorem** (with Zsuzsanna Dancso,  $\omega/WKO$ ). There is a bijection between the set of homomorphic expansions for  $w\mathcal{K}$  and the set of solutions of the Kashiwara-Vergne problem. **This is the tip of a major iceberg.**

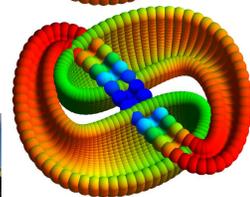
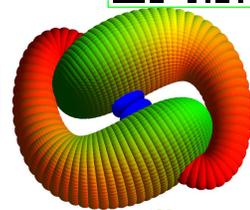
Dancso,  $\omega/ZD$



4D Knots.



$\omega/F$

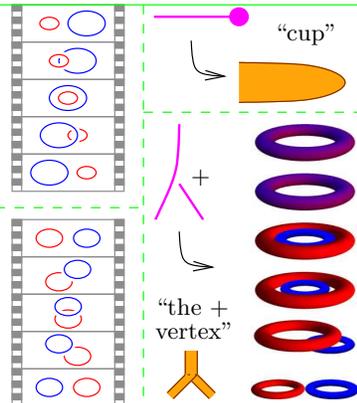
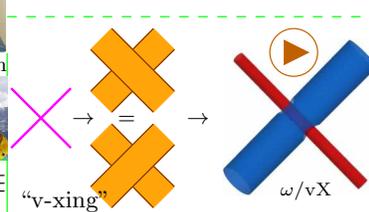
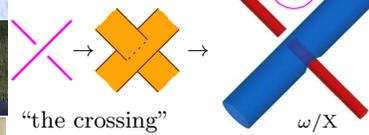


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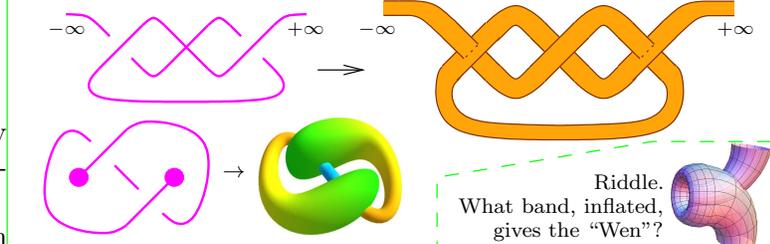


Dalvit  $\omega/Dal$

The Generators

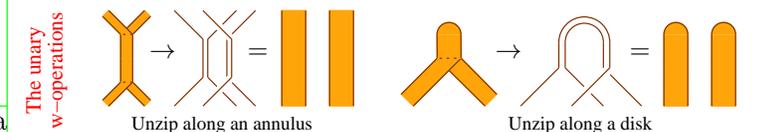
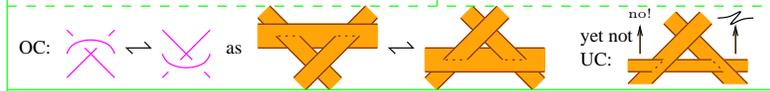
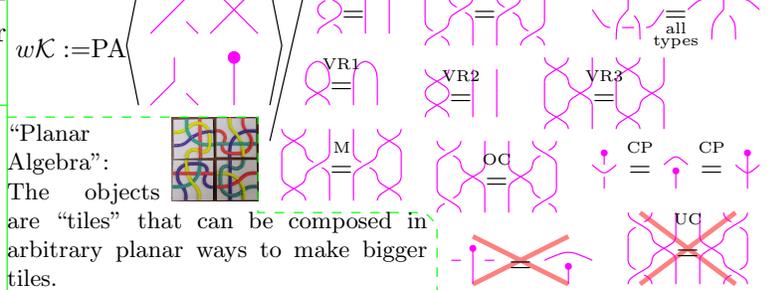


The Double Inflation Procedure.



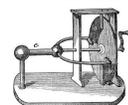
Riddle. What band, inflated, gives the "Wen"?

wKO.



The Machine generalizes to arbitrary algebraic structures!

$\omega/mac$



"God created the knots, all else in topology is the work of mortals."



Leopold Kronecker (modified)

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# 1. $\text{proj } \mathcal{K}^w(\uparrow_n) \cong_j \mathcal{U}((\mathfrak{a}_n \oplus \mathfrak{tder}_n) \times \mathfrak{tr}_n)$

— All Signs Are Wrong! —

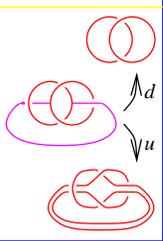
Dror Bar-Natan, Montpellier, June 2010, <http://www.math.toronto.edu/~drorbn/Talks/Montpellier-1006/>

I understand Drinfel'd and Alekseev-Torossian, I don't understand Etingof-Kazhdan yet, and I'm clueless about Kontsevich

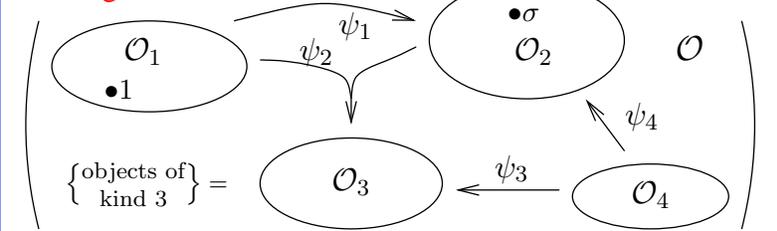
## Cans and Can't Yets.

(arbitrary algebraic structure)  $\xrightarrow[\text{machine}]{\text{projectivization}}$  (a problem in graded algebra)

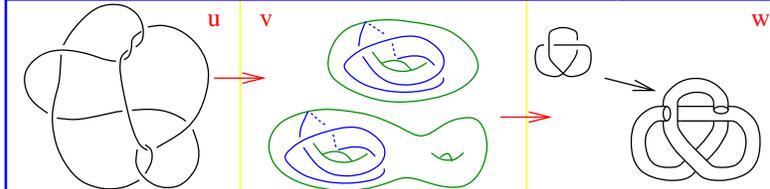
- Feed knot-things, get Lie algebra things.
- (u-knots)  $\rightarrow$  (Drinfel'd associators).
- (w-knots)  $\rightarrow$  (K-V-A-E-T).
- Dream: (v-knots)  $\rightarrow$  (Etingof-Kazhdan).
- Clueless: (???)  $\rightarrow$  (Kontsevich)?
- Goals: add to the Knot Atlas, produce a working AKT and touch ribbon 1-knots, rip benefits from *truly* understanding quantum groups.



## "An Algebraic Structure"



- Has kinds, objects, operations, and maybe constants.
- Perhaps subject to some axioms.
- We always allow formal linear combinations.



u-Knots (PA := Planar Algebra)

$$\{\text{knots \& links}\} = \text{PA} \langle \text{R123: } \begin{matrix} \text{ } \\ \text{ } \end{matrix} \Big| \begin{matrix} \text{ } \\ \text{ } \end{matrix} \rangle_{0 \text{ legs}}$$

## Circuit Algebras



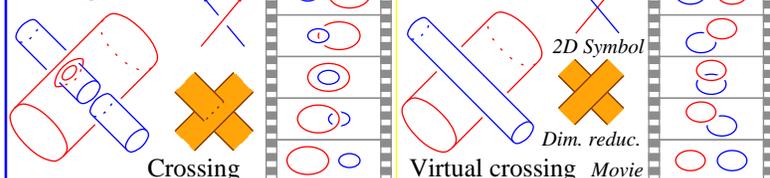
v-Tangles and w-Tangles (CA := Circuit Algebra)

$$\{\text{v-knots \& links}\} = \text{CA} \langle \text{R23: } \begin{matrix} \text{ } \\ \text{ } \end{matrix} \rangle$$

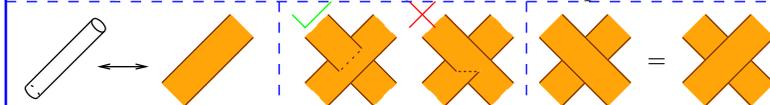
$$= \text{PA} \langle \text{VR123: } \begin{matrix} \text{ } \\ \text{ } \end{matrix} \Big| \text{R23} \rangle$$

$$\{\text{w-Tangles}\} = \text{v-Tangles} / \text{OC: } \begin{matrix} \text{ } \\ \text{ } \end{matrix} = \begin{matrix} \text{ } \\ \text{ } \end{matrix}$$

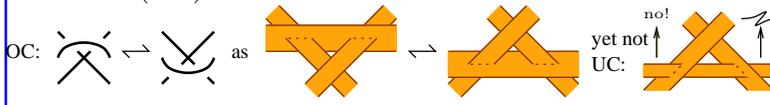
## The w-generators.



A Ribbon 2-Knot is a surface  $S$  embedded in  $\mathbb{R}^4$  that bounds an immersed handlebody  $B$ , with only "ribbon singularities"; a ribbon singularity is a disk  $D$  of transverse double points, whose preimages in  $B$  are a disk  $D_1$  in the interior of  $B$  and a disk  $D_2$  with  $D_2 \cap \partial B = \partial D_2$ , modulo isotopies of  $S$  alone.



The w-relations include R234, VR1234, D, Overcrossings Commute (OC) but not UC:



## Homomorphic expansions for a filtered algebraic structure $\mathcal{K}$ :

$$\text{ops} \subset \mathcal{K} = \mathcal{K}_0 \supset \mathcal{K}_1 \supset \mathcal{K}_2 \supset \mathcal{K}_3 \supset \dots$$

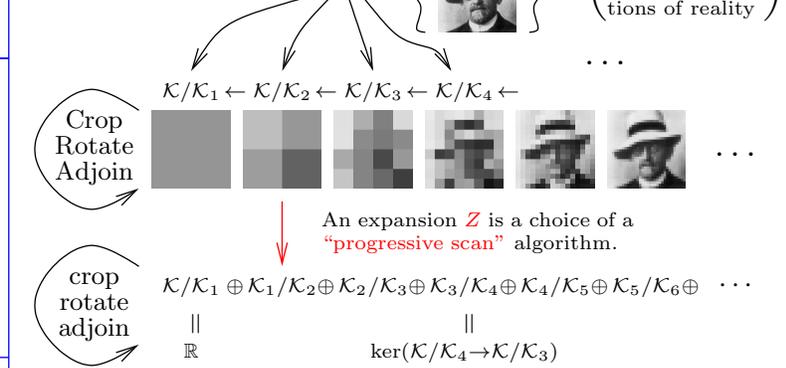
$$\downarrow \qquad \qquad \qquad \downarrow Z$$

$$\text{ops} \subset \text{gr } \mathcal{K} := \mathcal{K}_0/\mathcal{K}_1 \oplus \mathcal{K}_1/\mathcal{K}_2 \oplus \mathcal{K}_2/\mathcal{K}_3 \oplus \mathcal{K}_3/\mathcal{K}_4 \oplus \dots$$

An expansion is a filtered  $Z : \mathcal{K} \rightarrow \text{gr } \mathcal{K}$  that "covers" the identity on  $\text{gr } \mathcal{K}$ . A homomorphic expansion is an expansion that respects all relevant "extra" operations.

Reality.  $\text{gr } \mathcal{K}$  is often too hard. An  $\mathcal{A}$ -expansion is a graded "guess"  $\mathcal{A}$  with a surjection  $\tau : \mathcal{A} \rightarrow \text{gr } \mathcal{K}$  and a filtered  $Z : \mathcal{K} \rightarrow \mathcal{A}$  for which  $(\text{gr } Z) \circ \tau = I_{\mathcal{A}}$ . An  $\mathcal{A}$ -expansion confirms  $\mathcal{A}$  and yields an ordinary expansion. Same for "homomorphic".

Just for fun.



## Filtered algebraic structures are cheap and plenty.

In any  $\mathcal{K}$ , allow formal linear combinations, let  $\mathcal{K}_1 = \mathcal{I}$  be the ideal generated by differences (the "augmentation ideal"), and let  $\mathcal{K}_m := \langle (\mathcal{K}_1)^m \rangle$  (using all available "products"). In this case, set  $\text{proj } \mathcal{K} := \text{gr } \mathcal{K}$ .

Examples. 1. The projectivization of a group is a graded associative algebra.

2. Pure braids —  $PB_n$  is generated by  $x_{ij}$ , "strand  $i$  goes around strand  $j$  once", modulo "Reidemeister moves".  $A_n := \text{gr } PB_n$  is generated by  $t_{ij} := x_{ij} - 1$ , modulo the  $4T$  relations  $[t_{ij}, t_{ik} + t_{jk}] = 0$  (and some lesser ones too). Much happens in  $A_n$ , including the Drinfel'd theory of associators.

3. Quandle: a set  $Q$  with an op  $\wedge$  s.t.

$$1 \wedge x = 1, \quad x \wedge 1 = x, \quad (\text{appetizers})$$

$$(x \wedge y) \wedge z = (x \wedge z) \wedge (y \wedge z). \quad (\text{main})$$

$\text{proj } Q$  is a graded Leibniz algebra: Roughly, set  $\bar{v} := (v - 1)$  (these generate  $I$ !), feed  $1 + \bar{x}, 1 + \bar{y}, 1 + \bar{z}$  in (main), collect the surviving terms of lowest degree:

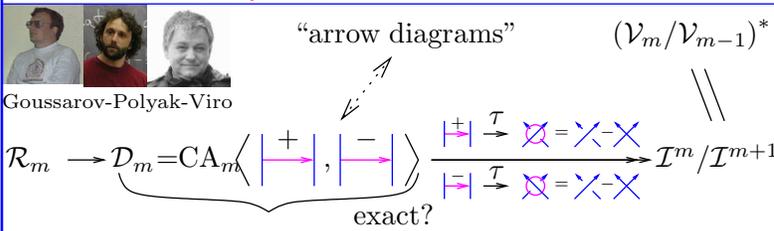
$$(\bar{x} \wedge \bar{y}) \wedge \bar{z} = (\bar{x} \wedge \bar{z}) \wedge \bar{y} + \bar{x} \wedge (\bar{y} \wedge \bar{z}).$$

"God created the knots, all else in topology is the work of mortals."  
Leopold Kronecker (modified)

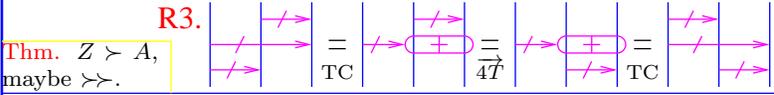
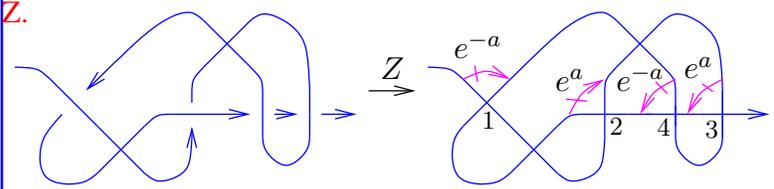
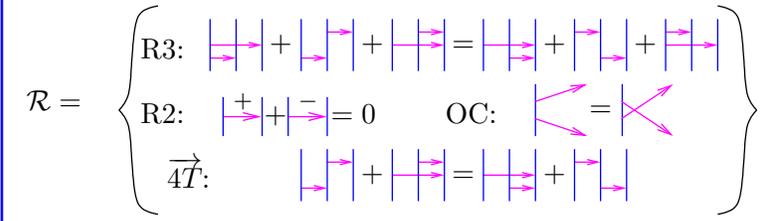
Also see <http://www.math.toronto.edu/~drorbn/papers/WKO/>



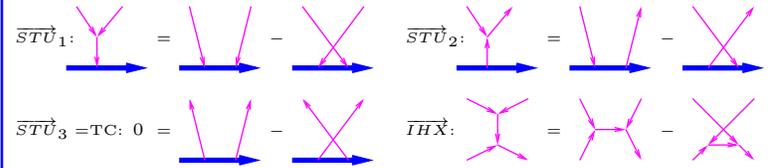
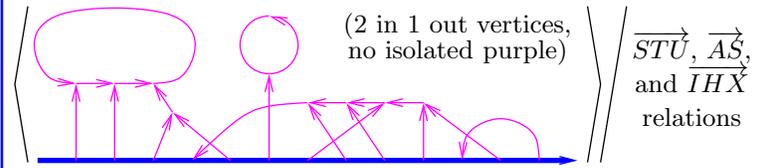
1.  $\text{proj } \mathcal{K}^w(\uparrow_n) \cong_j \mathcal{U}((\mathfrak{a}_n \oplus \mathfrak{tder}_n) \times \mathfrak{tr}_n)$ , continued.



**Imperfect Thumb-Rule.** Take R3 (say), substitute  $\times \rightarrow \times +$ , keep the lowest degree terms that don't immediately die:



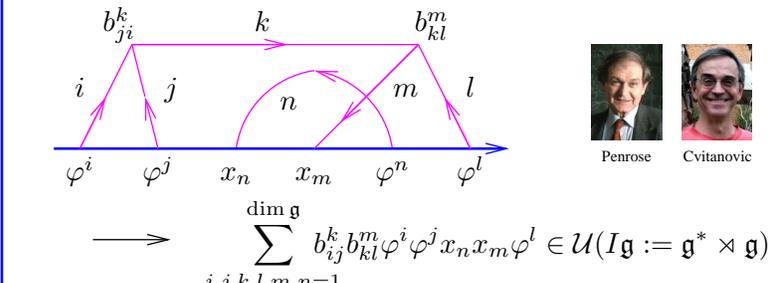
**The Bracket-Rise Theorem.**  $\mathcal{A}^w(\uparrow_1)$  is isomorphic to



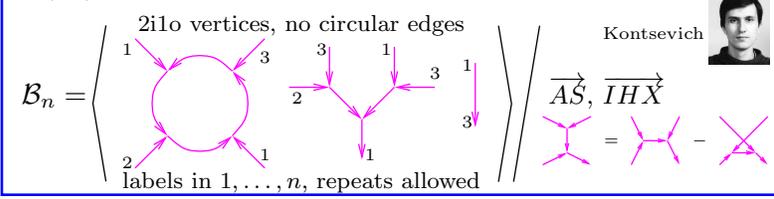
**Proof.**

**Corollaries.** (1) Related to Lie algebras! (2) Only wheels and isolated arrows persist.

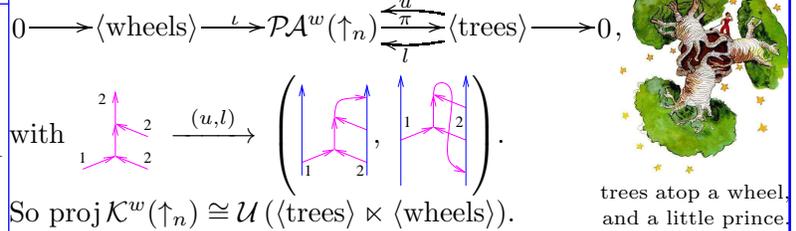
**To Lie Algebras.** With  $(x_i)$  and  $(\varphi^j)$  dual bases of  $\mathfrak{g}$  and  $\mathfrak{g}^*$  and with  $[x_i, x_j] = \sum b_{ij}^k x_k$ , we have  $\mathcal{A}^w \rightarrow \mathcal{U}$  via



**Theorem (PBW, “ $\mathcal{U}(\text{Ig})^{\otimes n} \cong \mathcal{S}(\text{Ig})^{\otimes n}$ ”).** As vector spaces,  $\mathcal{A}^w(\uparrow_n) \cong \mathcal{B}_n$ , where



**Wheels and Trees.** With  $\mathcal{P}$  for Primitives,



**Some A-T Notions.**  $\mathfrak{a}_n$  is the vector space with basis  $x_1, \dots, x_n$ ,  $\text{lie}_n = \text{lie}(\mathfrak{a}_n)$  is the free Lie algebra,  $\text{Ass}_n = \mathcal{U}(\text{lie}_n)$  is the free associative algebra “of words”,  $\text{tr} : \text{Ass}_n^+ \rightarrow \mathfrak{tr}_n = \text{Ass}_n^+ / (x_{i_1} x_{i_2} \dots x_{i_m} = x_{i_2} \dots x_{i_m} x_{i_1})$  is the “trace” into “cyclic words”,  $\mathfrak{der}_n = \mathfrak{der}(\text{lie}_n)$  are all the derivations, and  $\mathfrak{tder}_n = \{D \in \mathfrak{der}_n : \forall i \exists a_i \text{ s.t. } D(x_i) = [x_i, a_i]\}$

are “tangential derivations”, so  $D \leftrightarrow (a_1, \dots, a_n)$  is a vector space isomorphism  $\mathfrak{a}_n \oplus \mathfrak{tder}_n \cong \bigoplus_n \text{lie}_n$ . Finally,  $\text{div} : \mathfrak{tder}_n \rightarrow \mathfrak{tr}_n$  is  $(a_1, \dots, a_n) \mapsto \sum_k \text{tr}(x_k(\partial_k a_k))$ , where for  $a \in \text{Ass}_n^+$ ,  $\partial_k a \in \text{Ass}_n$  is determined by  $a = \sum_k (\partial_k a) x_k$ , and  $j : \text{TAut}_n = \exp(\mathfrak{tder}_n) \rightarrow \mathfrak{tr}_n$  is  $j(e^D) = \frac{e^D - 1}{D} \cdot \text{div } D$ .

**Theorem.** Everything matches.  $\langle \text{trees} \rangle$  is  $\mathfrak{a}_n \oplus \mathfrak{tder}_n$  as Lie algebras,  $\langle \text{wheels} \rangle$  is  $\mathfrak{tr}_n$  as  $\langle \text{trees} \rangle / \mathfrak{tder}_n$ -modules,  $\text{div } D = \iota^{-1}(u-l)(D)$ , and  $e^{uD} e^{-lD} = e^{jD}$ .

**Differential Operators.** Interpret  $\hat{\mathcal{U}}(\text{Ig})$  as tangential differential operators on  $\text{Fun}(\mathfrak{g})$ :

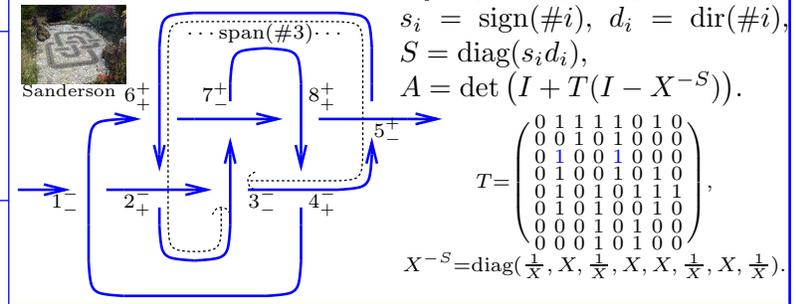
- $\varphi \in \mathfrak{g}^*$  becomes a multiplication operator.
  - $x \in \mathfrak{g}$  becomes a tangential derivation, in the direction of the action of  $\text{ad } x$ :  $(x\varphi)(y) := \varphi([x, y])$ .
- Trees become vector fields and  $uD \mapsto lD$  is  $D \mapsto D^*$ . So  $\text{div } D$  is  $D - D^*$  and  $jD = \log(e^D(e^D)^*) = \int_0^1 dt e^{tD} \text{div } D$ .

**Special Derivations.** Let  $\mathfrak{sder}_n = \{D \in \mathfrak{tder}_n : D(\sum x_i) = 0\}$ .

**Theorem.**  $\mathfrak{sder}_n = \pi\alpha(\text{proj u-tangles})$ , where  $\alpha$  is the obvious map  $\text{proj u-tangles} \rightarrow \text{proj w-tangles}$ .

**Proof.** After decoding, this becomes Lemma 6.1 of Drinfel'd’s amazing  $\text{Gal}(\mathbb{Q}/\mathbb{Q})$  paper.

**The Alexander Theorem.**



**Conjecture.** For u-knots,  $A$  is the Alexander polynomial.

**Theorem.** With  $w : x^k \mapsto w_k =$  (the  $k$ -wheel),  $Z = N \exp_{\mathcal{A}^w}(-w(\log_{\mathbb{Q}[[x]]} A(e^x))) \pmod{w_k w_l = w_{k+l}, Z = N \cdot A^{-1}(e^x)}$

This is the **ultimate Alexander invariant!** computable in polynomial time, local, composes well, behaves under cabling. Seems to significantly generalize the multi-variable Alexander polynomial and the theory of Milnor linking numbers. But it’s ugly, and much work remains.



## 2. w-Knots, Alekseev-Torossian, and baby Etingof-Kazhdan

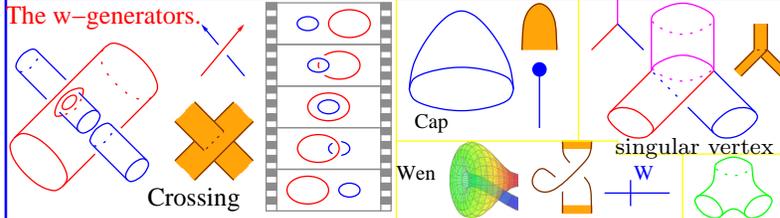
Dror Bar-Natan, Montpellier, June 2010, <http://www.math.toronto.edu/~drorbn/Talks/Montpellier-1006/>

I understand Drinfel'd and Alekseev-Torossian, I don't understand Etingof-Kazhdan yet, and I'm clueless about Kontsevich

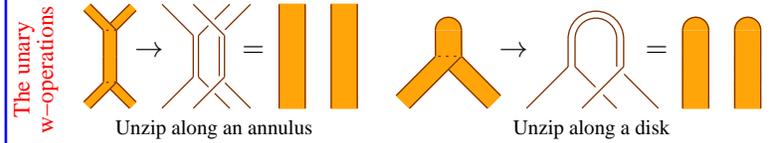
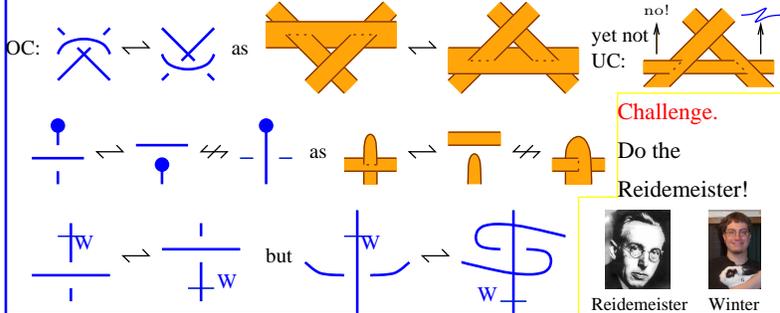
### Trivalent w-Tangles.

$$wTT = CA \left\langle \begin{array}{c|c|c} \text{w-} & \text{w-} & \text{unary w-} \\ \text{generators} & \text{relations} & \text{operations} \end{array} \right\rangle$$

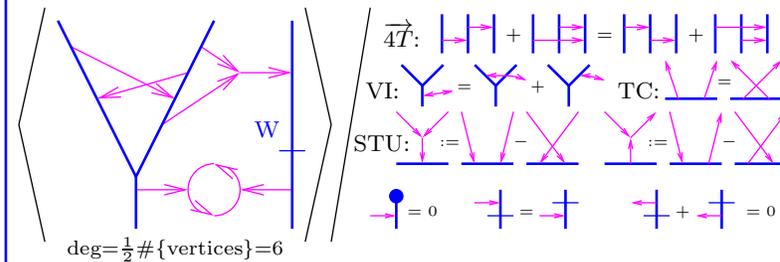
### The w-generators.



The w-relations include R234, VR1234, D, Overcrossings Commute (OC) but not UC,  $W^2 = 1$ , and funny interactions between the wen and the cap and over- and under-crossings:



w-Jacobi diagrams and  $\mathcal{A}$ .  $\mathcal{A}^w(Y \uparrow) \cong \mathcal{A}^w(\uparrow\uparrow\uparrow)$  is



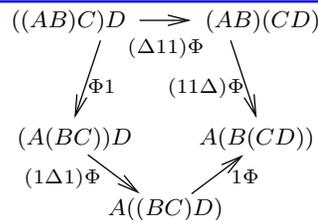
### An Associator:

$$(AB)C \xrightarrow{\Phi \in \mathcal{U}(\mathfrak{g})^{\otimes 3}} A(BC)$$

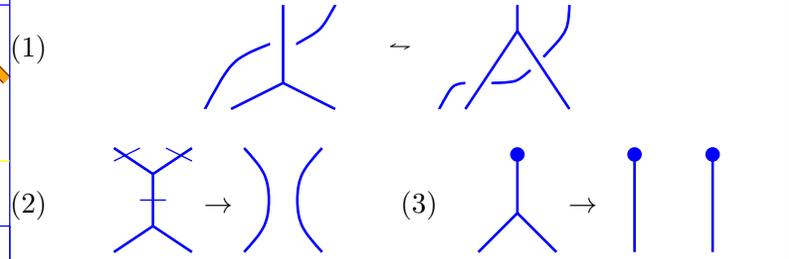
satisfying the "pentagon",

$$\Phi \cdot (1\Delta 1) \Phi \cdot 1\Phi = (\Delta 11) \Phi \cdot (11\Delta) \Phi$$

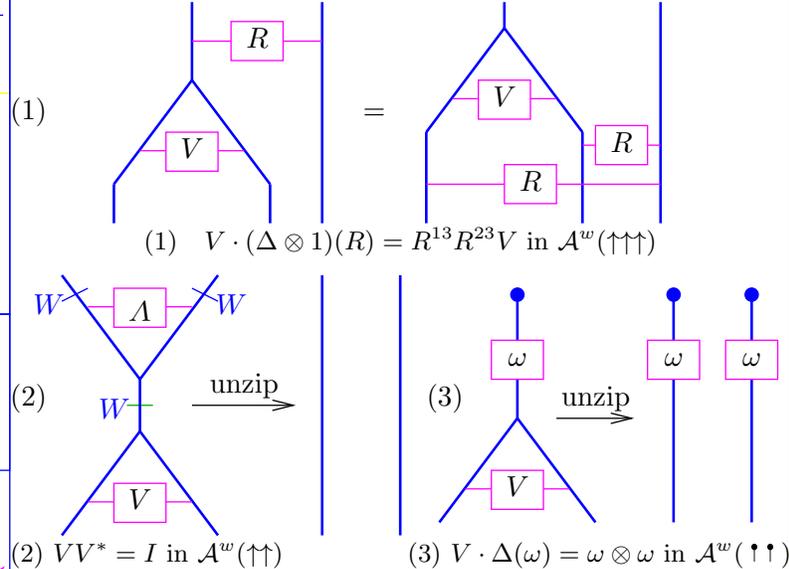
The hexagon? Never heard of it.



**Knot-Theoretic statement.** There exists a homomorphic expansion  $Z$  for trivalent w-tangles. In particular,  $Z$  should respect R4 and intertwine annulus and disk unzips:



**Diagrammatic statement.** Let  $R = \exp \uparrow \uparrow \in \mathcal{A}^w(\uparrow\uparrow)$ . There exist  $\omega \in \mathcal{A}^w(\uparrow)$  and  $V \in \mathcal{A}^w(\uparrow\uparrow)$  so that



**Alekseev-Torossian statement.** There are elements  $F \in \text{TAut}_2$  and  $a \in \mathfrak{tr}_1$  such that

$$F(x+y) = \log e^x e^y \quad \text{and} \quad jF = a(x) + a(y) - a(\log e^x e^y).$$

**Theorem.** The Alekseev-Torossian statement is equivalent to the knot-theoretic statement.

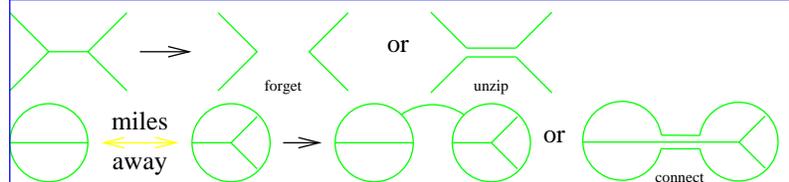
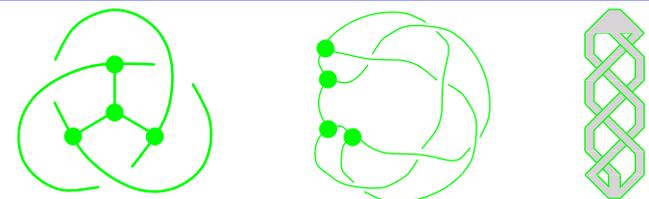
**Proof.** Write  $V = e^c e^{uD}$  with  $c \in \mathfrak{tr}_2$ ,  $D \in \mathfrak{tder}_2$ , and  $\omega = e^b$  with  $b \in \mathfrak{tr}_1$ . Then (1)  $\Leftrightarrow e^{uD}(x+y)e^{-uD} = \log e^x e^y$ , (2)  $\Leftrightarrow I = e^c e^{uD}(e^{uD})^* e^c = e^{2c} e^{jD}$ , and (3)  $\Leftrightarrow e^c e^{uD} e^{b(x+y)} = e^{b(x)+b(y)} \Leftrightarrow e^c e^{b(\log e^x e^y)} = e^{b(x)+b(y)} \Leftrightarrow c = b(x) + b(y) - b(\log e^x e^y)$ .

### The Alekseev-Torossian Correspondence.

$$\{\text{Drinfel'd Associators}\} \simeq \{\text{Solutions of KV}\}.$$

We need an even bigger algebraic structure!

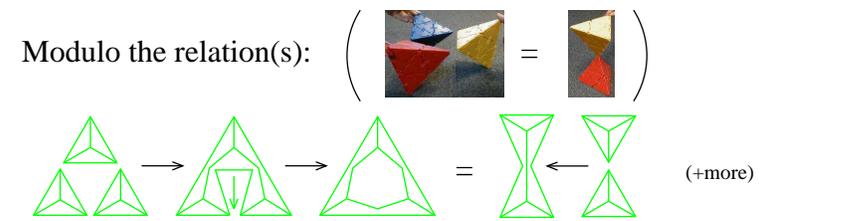
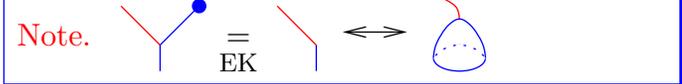
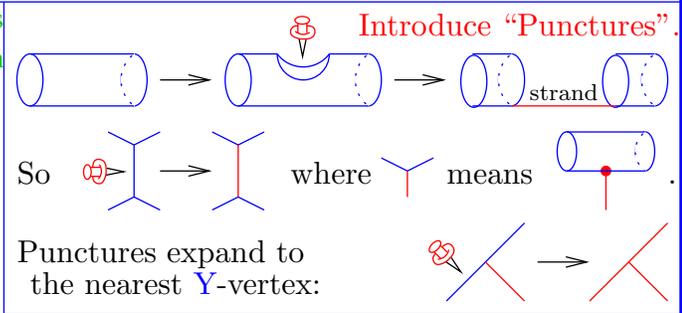
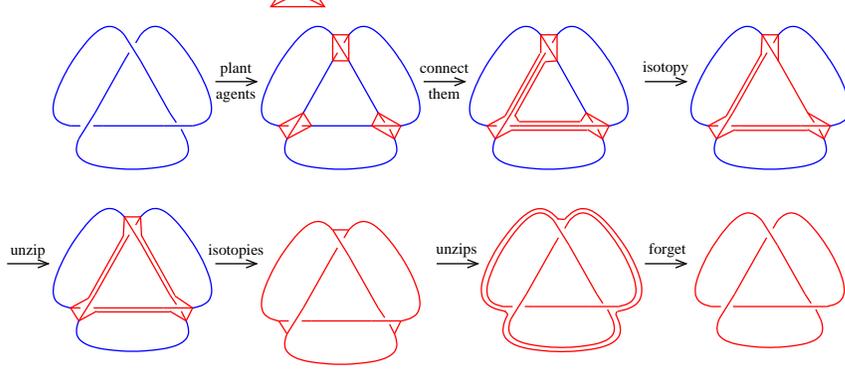
$$\left( \begin{array}{c} \text{green knotted trivalent} \\ \text{graphs in } \mathbb{R}^3 (u) \end{array} \right) \xrightarrow{\alpha_{\bar{w}}} \left( \begin{array}{c} \text{blue tubes and red} \\ \text{strings in } \mathbb{R}^4 (\bar{w}) \end{array} \right)$$



2. w-Knots, Alekseev-Torossian, and baby Etingof-Kazhdan, continued.

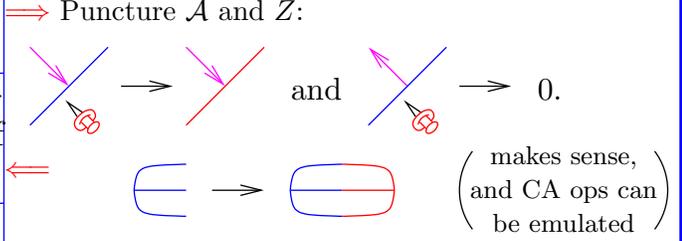
Using moves, KTG is generated by ribbon twists and the tetrahedron

All strands here are green

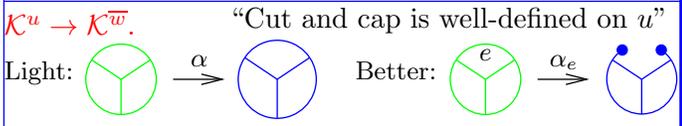
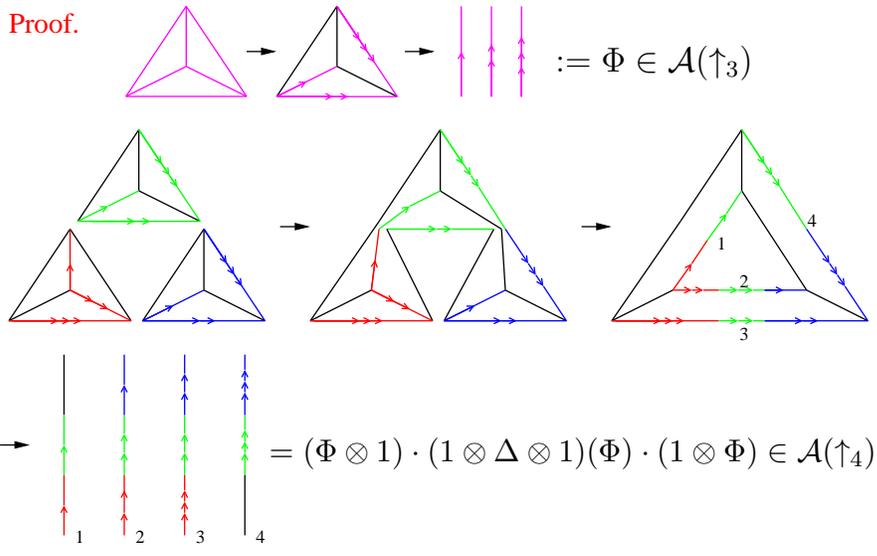


$\mathcal{K}^w$ . Allow tubes and strands and tube-strand vertices as above, yet allow only "compact" knots — nothing runs to  $\infty$ .

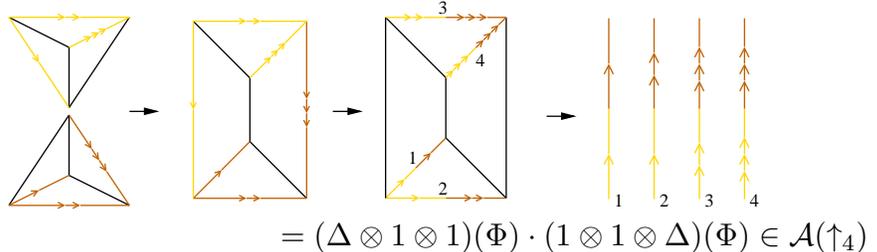
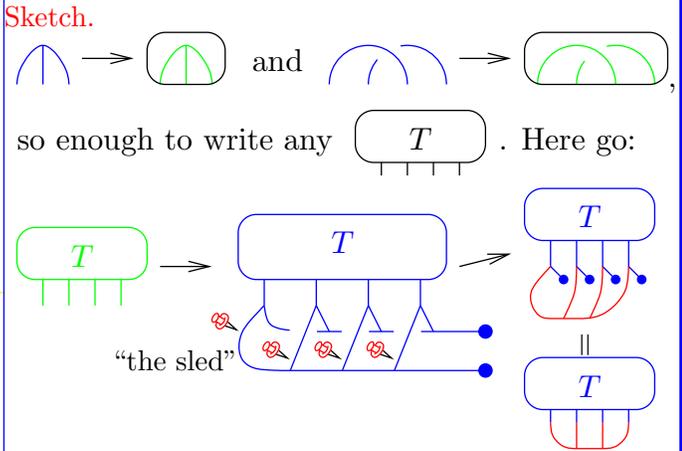
$\mathcal{K}^w \leftrightarrow \mathcal{K}^{\bar{w}}$  equivalence.  $\mathcal{K}^w$  has a homomorphic expansion iff  $\mathcal{K}^{\bar{w}}$  has a homomorphic expansion.



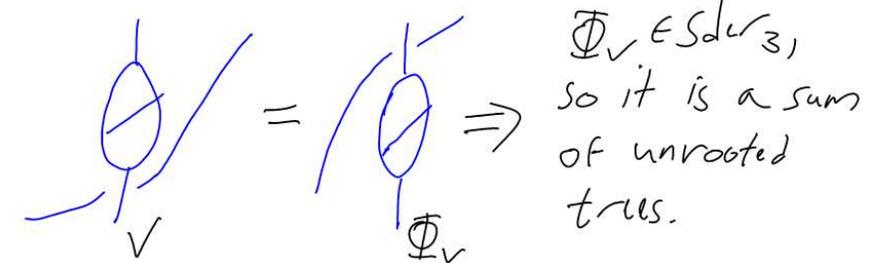
Claim. With  $\Phi := Z(\Delta)$ , the above relation becomes equivalent to the Drinfel'd's pentagon of the theory of quasi-Hopf algebras.

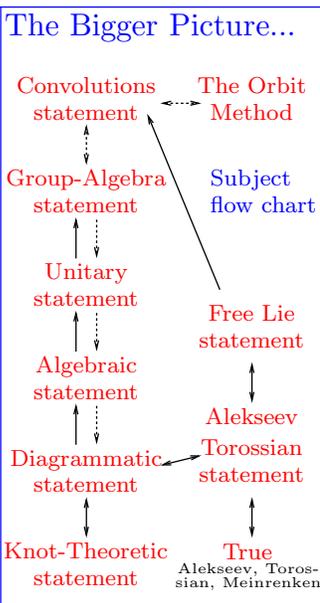


Theorem. The generators of  $\mathcal{K}^{\bar{w}}$  can be written in terms of the generators of  $\mathcal{K}^u$  (i.e., given  $\Phi$ , can write a formula for  $V$ ).



{SolvK}  $\rightarrow$  {Associators}: Trivial - a tetrahedron has 4 vertices.





**What are w-Trivalent Tangles?** (PA := Planar Algebra)

$\{ \text{knots} \} = \text{PA} \langle \text{R123} : \text{R123} \rangle$   
 $\{ \text{trivalent tangles} \} = \text{PA} \langle \text{R23, R4} \rangle$   
 $w\text{TT} = \{ \text{trivalent w-tangles} \} = \text{PA} \langle \text{w-generators} \mid \text{w-relations} \mid \text{unary w-operations} \rangle$

**The w-generators.**

**Broken surface**  
 Crossing:  $\text{Dim. reduc.}$   
 Virtual crossing:  $\text{Movie}$



**Cap Wen w Vertices**

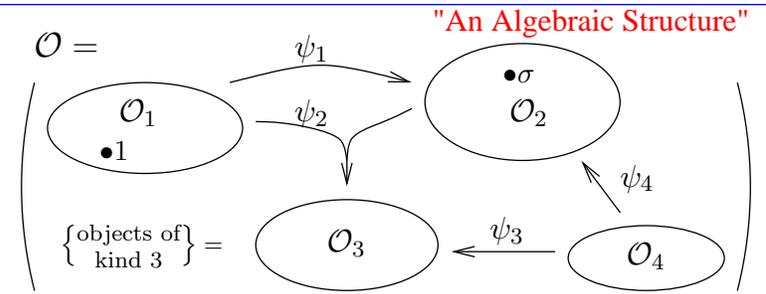
**Homomorphic expansions** for a filtered algebraic structure  $\mathcal{K}$ :

$$\text{ops}^{\leftarrow} \mathcal{K} = \mathcal{K}_0 \supset \mathcal{K}_1 \supset \mathcal{K}_2 \supset \mathcal{K}_3 \supset \dots$$

$$\text{ops}^{\leftarrow} \text{gr } \mathcal{K} := \mathcal{K}_0/\mathcal{K}_1 \oplus \mathcal{K}_1/\mathcal{K}_2 \oplus \mathcal{K}_2/\mathcal{K}_3 \oplus \mathcal{K}_3/\mathcal{K}_4 \oplus \dots$$

An **expansion** is a filtration  $Z : \mathcal{K} \rightarrow \text{gr } \mathcal{K}$  that "covers" the identity on  $\text{gr } \mathcal{K}$ . A **homomorphic expansion** is an expansion that respects all relevant "extra" operations.

**Filtered algebraic structures are cheap and plenty.** In any  $\mathcal{K}$ , allow formal linear combinations, let  $\mathcal{K}_1$  be the ideal generated by differences (the "augmentation ideal"), and let  $\mathcal{K}_m := \langle (\mathcal{K}_1)^m \rangle$  (using all available "products").



- Has kinds, objects, operations, and maybe constants.
- Perhaps subject to some axioms.
- We always allow formal linear combinations.

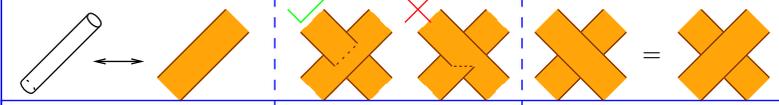
**Example: Pure Braids.**  $PB_n$  is generated by  $x_{ij}$ , "strand  $i$  goes around strand  $j$  once", modulo "Reidemeister moves".  $A_n := \text{gr } PB_n$  is generated by  $t_{ij} := x_{ij} - 1$ , modulo the 4T relations  $[t_{ij}, t_{ik} + t_{jk}] = 0$  (and some lesser ones too). Much happens in  $A_n$ , including the Drinfel'd theory of associators.

**Our case(s).**

$\mathcal{K} \xrightarrow{Z: \text{high algebra}} \mathcal{A} := \text{gr } \mathcal{K} \xrightarrow{\text{given a "Lie" algebra } \mathfrak{g}} \mathcal{U}(\mathfrak{g})$   
 solving finitely many equations in finitely many unknowns      low algebra: pictures represent formulas

$\mathcal{K}$  is knot theory or **topology**;  $\text{gr } \mathcal{K}$  is finite **combinatorics**: bounded-complexity diagrams modulo simple relations.

A **Ribbon 2-Knot** is a surface  $S$  embedded in  $\mathbb{R}^4$  that bounds an immersed handlebody  $B$ , with only "ribbon singularities"; a ribbon singularity is a disk  $D$  of transverse double points, whose preimages in  $B$  are a disk  $D_1$  in the interior of  $B$  and a disk  $D_2$  with  $D_2 \cap \partial B = \partial D_2$ , modulo isotopies of  $S$  alone.



The **w-relations** include R234, VR1234, M, Overcrossings Commute (OC) but not UC,  $W^2 = 1$ , and **funny interactions** between the wen and the cap and over- and under-crossings:

OC:  $\text{OC} : \text{Crossing} \leftrightarrow \text{Crossing}$  as  $\text{Crossing} \leftrightarrow \text{Crossing}$  yet not UC:  $\text{Crossing} \not\leftrightarrow \text{Crossing}$

$\text{Wen} \leftrightarrow \text{Cap}$  as  $\text{Wen} \leftrightarrow \text{Cap}$   
 $\text{Wen} \leftrightarrow \text{Crossing}$  but  $\text{Wen} \not\leftrightarrow \text{Crossing}$

**Challenge.** Do the Reidemeister!

**The unary w-operations**

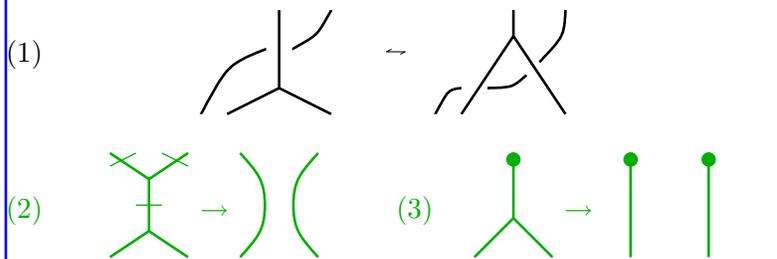
Unzip along an annulus:  $\text{Y-junction} \rightarrow \text{Two parallel strands}$   
 Unzip along a disk:  $\text{Y-junction} \rightarrow \text{Two parallel strands}$

**Just for fun.**

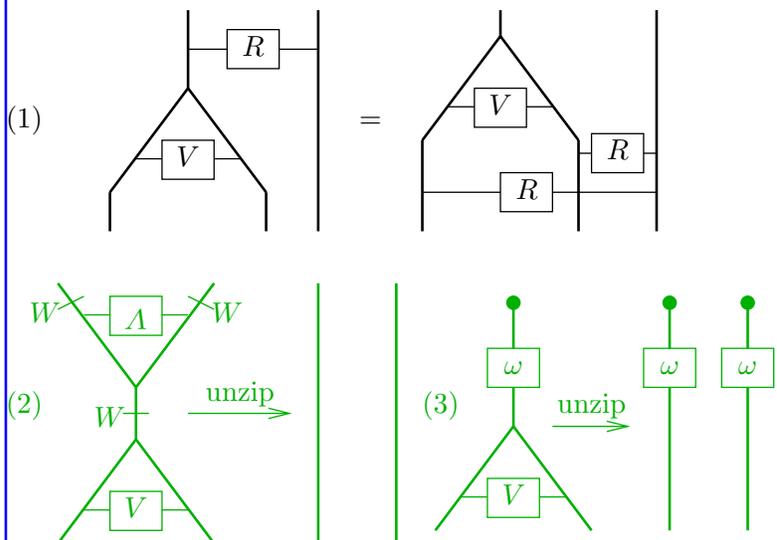
$\mathcal{K} = \{ \text{Diagrams} \} = \left( \text{The set of all b/w 2D projections of reality} \right)$   
 $\mathcal{K}/\mathcal{K}_1 \leftarrow \mathcal{K}/\mathcal{K}_2 \leftarrow \mathcal{K}/\mathcal{K}_3 \leftarrow \mathcal{K}/\mathcal{K}_4 \leftarrow \dots$   
 Crop Rotate Adjoin      An expansion  $Z$  is a choice of a "progressive scan" algorithm.  
 $\mathcal{K}/\mathcal{K}_1 \oplus \mathcal{K}_1/\mathcal{K}_2 \oplus \mathcal{K}_2/\mathcal{K}_3 \oplus \mathcal{K}_3/\mathcal{K}_4 \oplus \mathcal{K}_4/\mathcal{K}_5 \oplus \mathcal{K}_5/\mathcal{K}_6 \oplus \dots$   
 $\mathbb{R} \parallel \ker(\mathcal{K}/\mathcal{K}_4 \rightarrow \mathcal{K}/\mathcal{K}_3)$

# Convolutions on Lie Groups and Lie Algebras and Ribbon 2-Knots, Page 2

**Knot-Theoretic statement.** There exists a homomorphic expansion  $Z$  for trivalent w-tangles. In particular,  $Z$  should respect  $R4$  and intertwine annulus and disk unzips:



**Diagrammatic statement.** Let  $R = \exp \uparrow \in \mathcal{A}^w(\uparrow\uparrow)$ . There exist  $\omega \in \mathcal{A}^w(\uparrow)$  and  $V \in \mathcal{A}^w(\uparrow\uparrow)$  so that



**Algebraic statement.** With  $I\mathfrak{g} := \mathfrak{g}^* \rtimes \mathfrak{g}$ , with  $c : \hat{U}(I\mathfrak{g}) \rightarrow \hat{U}(\mathfrak{g})/\hat{U}(\mathfrak{g}) = \hat{S}(\mathfrak{g}^*)$  the obvious projection, with  $S$  the antipode of  $\hat{U}(I\mathfrak{g})$ , with  $W$  the automorphism of  $\hat{U}(I\mathfrak{g})$  induced by flipping the sign of  $\mathfrak{g}^*$ , with  $r \in \mathfrak{g}^* \otimes \mathfrak{g}$  the identity element and with  $R = e^r \in \hat{U}(I\mathfrak{g}) \otimes \hat{U}(\mathfrak{g})$  there exist  $\omega \in \hat{S}(\mathfrak{g}^*)$  and  $V \in \hat{U}(I\mathfrak{g})^{\otimes 2}$  so that

- (1)  $V(\Delta \otimes 1)(R) = R^{13}R^{23}V$  in  $\hat{U}(I\mathfrak{g})^{\otimes 2} \otimes \hat{U}(\mathfrak{g})$
- (2)  $V \cdot SWV = 1$
- (3)  $(c \otimes c)(V\Delta(\omega)) = \omega \otimes \omega$

**Unitary statement.** There exists  $\omega \in \text{Fun}(\mathfrak{g})^G$  and an (infinite order) tangential differential operator  $V$  defined on  $\text{Fun}(\mathfrak{g}_x \times \mathfrak{g}_y)$  so that

- (1)  $V\widehat{e^{x+y}} = \widehat{e^x e^y} V$  (allowing  $\hat{U}(\mathfrak{g})$ -valued functions)
- (2)  $VV^* = I$
- (3)  $V\omega_{x+y} = \omega_x \omega_y$

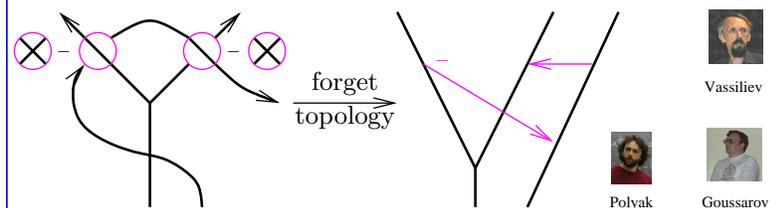
**Group-Algebra statement.** There exists  $\omega^2 \in \text{Fun}(\mathfrak{g})^G$  so that for every  $\phi, \psi \in \text{Fun}(\mathfrak{g})^G$  (with small support), the following holds in  $\hat{U}(\mathfrak{g})$ :

$$\iint_{\mathfrak{g} \times \mathfrak{g}} \phi(x)\psi(y)\omega_{x+y}^2 e^{x+y} = \iint_{\mathfrak{g} \times \mathfrak{g}} \phi(x)\psi(y)\omega_x^2 \omega_y^2 e^x e^y. \quad (\text{shhh, this is Duflo})$$

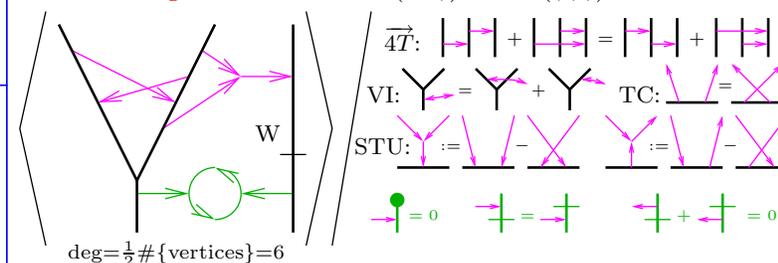
**Convolutions statement** (Kashiwara-Vergne). Convolutions of invariant functions on a Lie group agree with convolutions of invariant functions on its Lie algebra. More accurately, let  $G$  be a finite dimensional Lie group and let  $\mathfrak{g}$  be its Lie algebra, let  $j : \mathfrak{g} \rightarrow \mathbb{R}$  be the Jacobian of the exponential map  $\exp : \mathfrak{g} \rightarrow G$ , and let  $\Phi : \text{Fun}(G) \rightarrow \text{Fun}(\mathfrak{g})$  be given by  $\Phi(f)(x) := j^{1/2}(x)f(\exp x)$ . Then if  $f, g \in \text{Fun}(G)$  are Ad-invariant and supported near the identity, then

$$\Phi(f) \star \Phi(g) = \Phi(f \star g).$$

**From wTT to  $\mathcal{A}^w$ .**  $\text{gr}_m \text{wTT} := \{m\text{-cubes}\} / \{(m+1)\text{-cubes}\}$ :



**w-Jacobi diagrams and  $\mathcal{A}$ .**  $\mathcal{A}^w(Y \uparrow) \cong \mathcal{A}^w(\uparrow\uparrow\uparrow)$  is



**Diagrammatic to Algebraic.** With  $(x_i)$  and  $(\varphi^j)$  dual bases of  $\mathfrak{g}$  and  $\mathfrak{g}^*$  and with  $[x_i, x_j] = \sum b_{ij}^k x_k$ , we have  $\mathcal{A}^w \rightarrow \mathcal{U}$  via

$$\sum_{i,j,k,l,m,n=1}^{\dim \mathfrak{g}} b_{ij}^k b_{kl}^m \varphi^i \varphi^j x_n x_m \varphi^n \in \mathcal{U}(I\mathfrak{g})$$

**Unitary  $\iff$  Algebraic.** The key is to interpret  $\hat{U}(I\mathfrak{g})$  as tangential differential operators on  $\text{Fun}(\mathfrak{g})$ :

- $\varphi \in \mathfrak{g}^*$  becomes a multiplication operator.
- $x \in \mathfrak{g}$  becomes a tangential derivation, in the direction of the action of  $\text{ad } x$ :  $(x\varphi)(y) := \varphi([x, y])$ .
- $c : \hat{U}(I\mathfrak{g}) \rightarrow \hat{U}(I\mathfrak{g})/\hat{U}(\mathfrak{g}) = \hat{S}(\mathfrak{g}^*)$  is “the constant term”.

**Unitary  $\implies$  Group-Algebra.**

$$\iint \omega_{x+y}^2 e^{x+y} \phi(x)\psi(y) = \langle \omega_{x+y}, \omega_{x+y} e^{x+y} \phi(x)\psi(y) \rangle = \langle V\omega_{x+y}, V e^{x+y} \phi(x)\psi(y)\omega_{x+y} \rangle = \langle \omega_x \omega_y, e^x e^y V \phi(x)\psi(y)\omega_{x+y} \rangle = \langle \omega_x \omega_y, e^x e^y \phi(x)\psi(y)\omega_x \omega_y \rangle = \iint \omega_x^2 \omega_y^2 e^x e^y \phi(x)\psi(y).$$

**Convolutions and Group Algebras** (ignoring all Jacobians). If  $G$  is finite,  $A$  is an algebra,  $\tau : G \rightarrow A$  is multiplicative then  $(\text{Fun}(G), \star) \cong (A, \cdot)$  via  $L : f \mapsto \sum f(a)\tau(a)$ . For Lie  $(G, \mathfrak{g})$ ,

$$\begin{array}{ccc} (\mathfrak{g}, +) \ni x \xrightarrow{\tau_0 = \exp_S} e^x \in \hat{S}(\mathfrak{g}) & & \text{Fun}(\mathfrak{g}) \xrightarrow{L_0} \hat{S}(\mathfrak{g}) \\ \downarrow \exp_G & \searrow \exp_U & \downarrow \chi \\ (G, \cdot) \ni e^x \xrightarrow{\tau_1} e^x \in \hat{U}(\mathfrak{g}) & & \text{Fun}(G) \xrightarrow{L_1} \hat{U}(\mathfrak{g}) \end{array} \quad \text{so} \quad \begin{array}{ccc} & & \downarrow \Phi^{-1} \\ & & \downarrow \chi \end{array}$$

with  $L_0\psi = \int \psi(x)e^x dx \in \hat{S}(\mathfrak{g})$  and  $L_1\Phi^{-1}\psi = \int \psi(x)e^x \in \hat{U}(\mathfrak{g})$ . Given  $\psi_i \in \text{Fun}(\mathfrak{g})$  compare  $\Phi^{-1}(\psi_1) \star \Phi^{-1}(\psi_2)$  and  $\Phi^{-1}(\psi_1 \star \psi_2)$  in  $\hat{U}(\mathfrak{g})$ : (shhh,  $L_{0/1}$  are “Laplace transforms”)

$$\star \text{ in } G : \iint \psi_1(x)\psi_2(y)e^x e^y \quad \star \text{ in } \mathfrak{g} : \iint \psi_1(x)\psi_2(y)e^{x+y}$$

- We skipped...**
- The Alexander polynomial and Milnor numbers.
  - v-Knots, quantum groups and Etingof-Kazhdan.
  - u-Knots, Alekseev-Torossian, and BF theory and the successful and Drinfel'd associators.
  - The simplest problem hyperbolic geometry solves.