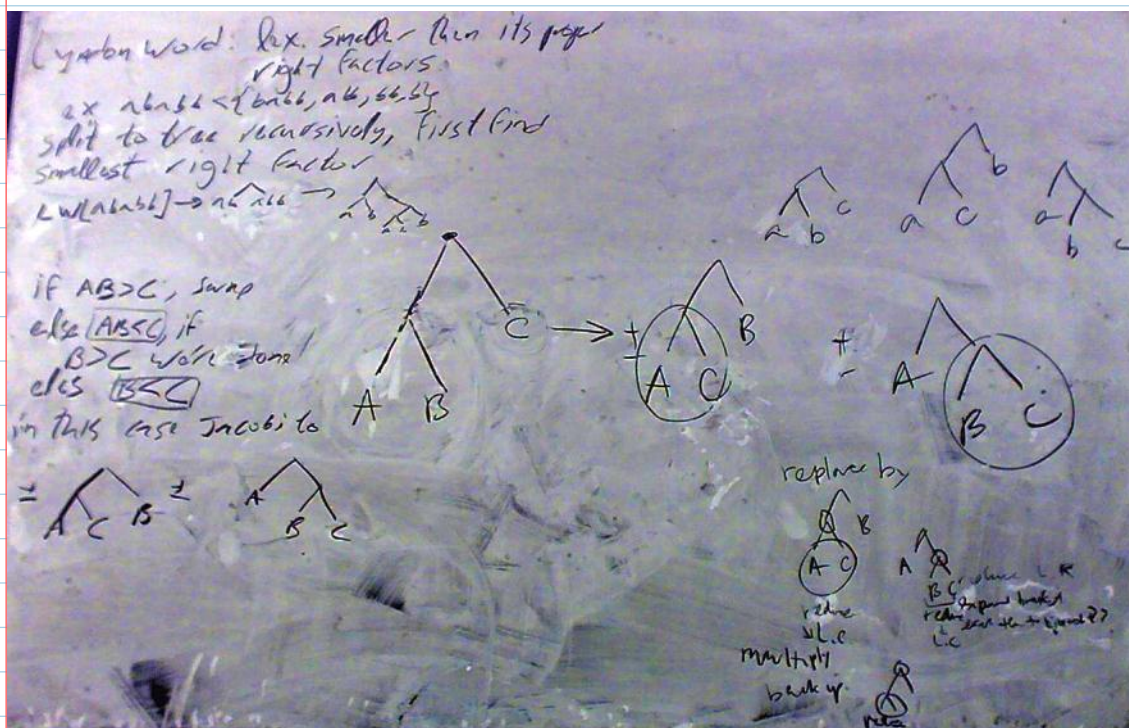


Lyndon Words

June-27-12
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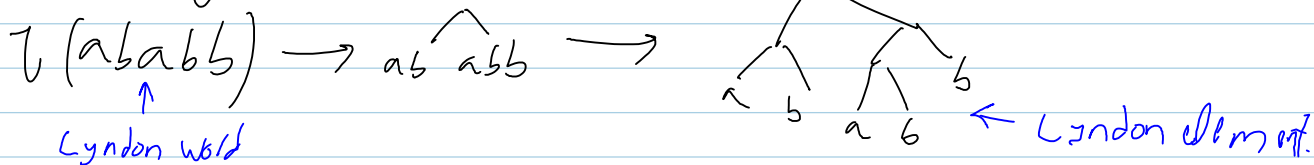
<http://katlas.math.toronto.edu/drorbn/bbs/show?shot=Chu-071214-182203.jpg>:



Lyndon word: Lexicographically smaller than its proper right subwords.

Example: $ababb < \{babb, abb, bb, b\}$

Can split to a tree recursively - First Find smallest right factor: [in itself it must be Lyndon]



Every element in Free Lie is a combination of Lyndon elements:

Take $[[A, B], C]$ where AB & C are Lyndon.



IF $C < AB$, swap

else $AB < C$ then ABC is Lyndon

if $B \geq C$ then $\tau(ABC) = [\tau(AB), \tau(C)]$

otherwise $A < B < C$; then ACB is also Lyndon

use Jacobi to write

$$[[A, B], C] = [[A, C], B] + [A, [B, C]]$$

claim. The bracket of two Lyndon elements is a combination of Lyndon elements greater or equal to the smaller of their two concatenations.

Proof. WLOG, we are looking at $[\tau(L), \tau(R)]$ where $L < R$, otherwise use AS. Clearly, LR is Lyndon.

We then argue by induction on the total degree, and then on the lexicographical ordering on LR , starting from the end and going down.

IF L is atomic then R is the minimal right factor of LR so $\tau(LR) = [\tau(L), \tau(R)]$ and all is well.

Otherwise $L = AB$, $R = C$, with $AB < C$, and with B the minimal proper right factor of AB .

IF $B \geq C$, then C is the smallest right

factor of ABC so

$$\tau(ABC) = [\tau(AB), \tau(C)] = [\tau(L), \tau(R)]$$

and all is well.

otherwise $A < B < C$ and we have

$$[\tau(L), \tau(R)] = [[\tau A, \tau B], \tau C] = \text{using Jacobi}$$

$$= [\tau A, \underbrace{[\tau B, \tau C]}_1] + \underbrace{[[\tau A, \tau C], \tau B]}_2$$

$$= \tau(ABC) + \text{higher terms} \quad (\text{from 1})$$

$$+ \text{further higher terms} \quad (\text{from 2}).$$