

claim $L(\sum_{i=1}^n d_i x_i) = \sum d_i L(x_i)$

Check 3 things
 point or:
 check 1 property

prop Given function $L: V \rightarrow W$
 $(L \in \Omega \text{ lin Tr}) \iff (\forall x, y, c, L(cx+y) = cL(x) + L(y))$

Pf \implies ✓
 \Leftarrow Suppose $L(cx+y) = cL(x) + L(y)$

prop 1 \implies if $c=1 \implies$ prop #1.
 if $y=0, x=0, c=1 \implies L(0) = L(0) + L(0)$

prop 2 \implies if $y=0, L(cx) = cL(x) \implies 0 = L(0)$

example 1. $L: \mathbb{R}^2 \rightarrow \mathbb{R}^3$ by $L \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} 2a-b \\ 0 \\ a+7b \end{pmatrix}$

□

2. $L \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} a^3 \\ |b| \\ a+1 \end{pmatrix}$ not linear $\forall: L(0) = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$

$L(cx+cy) = L \begin{pmatrix} ca+cb \\ b \end{pmatrix} = L \begin{pmatrix} ca+cb \\ b \end{pmatrix} = \begin{pmatrix} (ca+cb)^3 \\ |cb+b| \\ ca+cb+1 \end{pmatrix}$

$c \begin{pmatrix} a^3 \\ |b| \\ a+1 \end{pmatrix} + \begin{pmatrix} a'^3 \\ |b'| \\ a'+1 \end{pmatrix} = cL(x) + L(y)$

$(ca+cb)^3 \stackrel{?}{=} ca^3 + a'^3$ ✗
 $|cb+b| = c|b| + |b'|$ ✗
 $ca+cb+1 = ca+a'+1$ ✓

3. $D: P_n \rightarrow P_{n-1} \quad X: P_n \rightarrow P_{n+1}$
 $DF := \frac{dF}{dx}$ | multiplication by x

* in quantum mechanics

$D(x^3 - 7x) = 3x^2 - 7$ | $\hat{x} \cdot \hat{p} = x \cdot p$
 momentum | $\hat{p}(7x) = 7x^2$ postn

3a is linear $D(F+g) = D(F) + D(g)$
 $(F+g)' = F'+g'$

(Uniqueness)

10/22/14

\Rightarrow Every $f \in P_n(F)$ can be written in a unique way as a

lin combination of the P_i 's.

Claim Suppose $g \in P_n(F)$ satisfies $g(x_i) = y_i$ then $g = f$

Proof By prior claim, $\exists d_i \in F$ satisfies

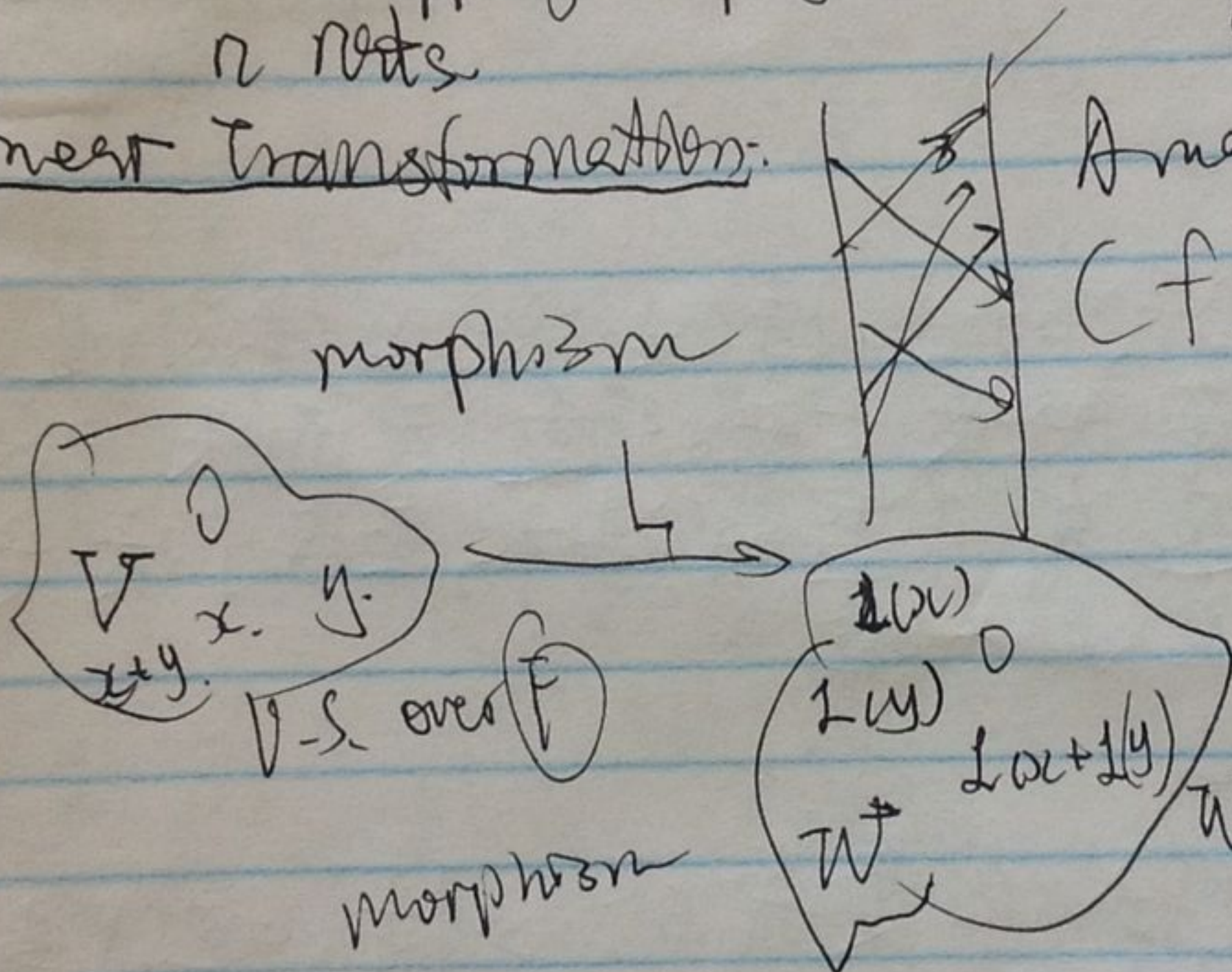
$$g = \sum d_i P_i \text{ but } y_k = f(x_k) = d_k$$

$$\text{So } \forall k, y_k = d_k \text{ so } g = \sum d_i P_i = \sum y_i P_i = f$$

Aside If $\forall P(x_i) = 0$ then P itself is 0

I.e. If P is a degree polynomial not equal to 0, then P has at most n roots.

Linear Transformation:



A map between 2 objects of the (f) kind that respect their structure.

$$0. L(0_V) = 0_W$$

$$1. \forall x, y \in V, L(x+y)$$

$$= L(x) + L(y)$$

$$2. \forall x \in V, \forall c \in F$$

$$L(cx) = c \cdot L(x)$$

Def Let V & W be V.S. over some field F . A Linear Transformation from $V \rightarrow W$ is a map (function) $L: V \rightarrow W$ satisfies

$$1. \forall x, y \in V, L(x+y) = L(x) + L(y)$$

$$2. \forall x \in V, c \in F, L(cx) = c \cdot L(x)$$

Claim If L is a L.T. then $L(0) = 0$

$$\text{Proof } L(0_V) = L(0_F \cdot 0_V) = 0_F L(0_V) = 0_W \quad \square$$

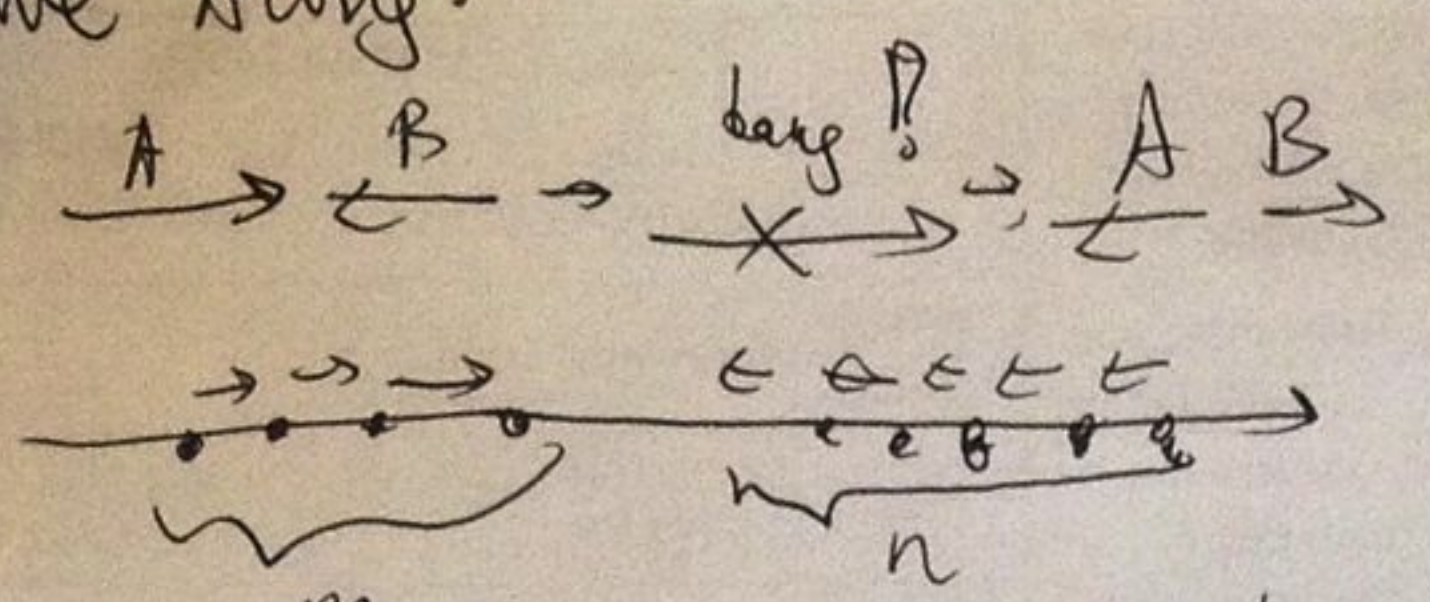
Claim $L(x-y) = L(x) - L(y)$

$$\text{Proof } L(x-y) = L(x + (-1)y) = L(x) + L((-1)y)$$

$$L(x) + L((-1)y) = L(x) + (-1)L(y) = L(x) - L(y)$$

Reading Along 1.6, 2.1 Hw 5 on Web/Print

Riddle Along:



Counting collisions.

To try Lagrange Interpolation Linear Algebra

Reminder $x_1, \dots, x_{n+1} \in F$ distinct
 $y_1, \dots, y_{n+1} \in F$ any

$\exists \! \! \! \int \! \! \! \int P \in P_n(F)$ s.t. $P(x_i) = y_i$

$$P_i(x) = \prod_{j \neq i} (x - x_j)$$

$$P_i(x_k) = \begin{cases} 0, & k \neq i \\ \neq 0, & k = i \end{cases}$$

Now set

$$P = 5P_1 + 2P_2 + 2P_3 = x^2 - 4x + 5$$

Example $x_{1,2,3} = 0, 1, 3$ $y_{1,2,3} = 5, 2, 2$

$$P(x) = 5 \quad P(1) = 2 = P(3)$$

$$P_1 = (x - x_2)(x - x_3) = (x - 1)(x - 3) = x^2 - 4x + 3$$

$$P_2 = (x - x_1)(x - x_3) = x(x - 3) = x^2 - 3x$$

$$P_3 = (x - x_1)(x - x_2) = x(x - 1) = x^2 - x$$

$$\text{Set } P_i(x) = \frac{P_i(x)}{P_i(x_i)} = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$$

$$P_i(x_k) = \begin{cases} 0, & k \neq i \\ 1, & k = i \end{cases}$$

Now set
 Indeed $P(x) = \sum_{i=1}^{n+1} y_i P_i(x)$

$$\begin{aligned} P(x_k) &= \sum y_i P_i(x_k) \\ &= 0 + 0 + \dots + y_k P_k(x_k) \\ &= y_k \cdot 1 = y_k \end{aligned}$$

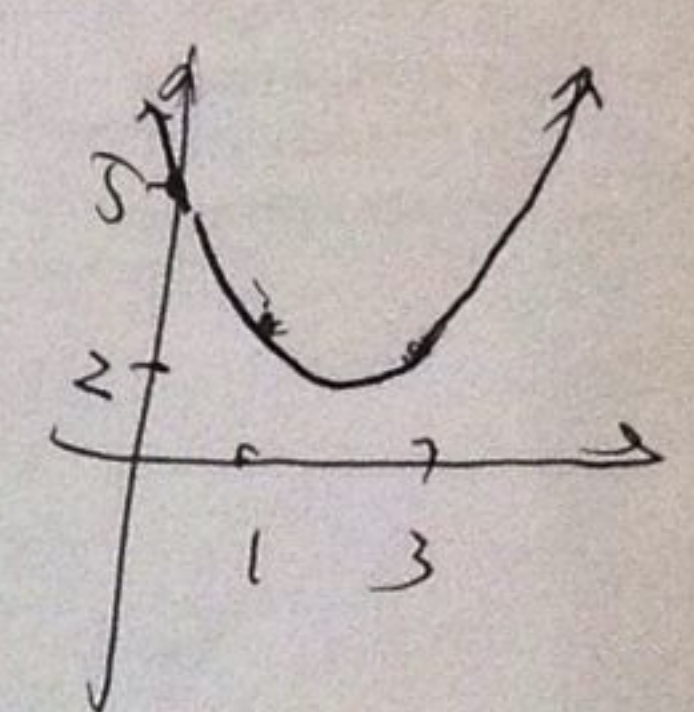
Uniqueness: $\beta = \{P_1, P_2, P_3, \dots, P_{n+1}\}$ is linearly independent.

Indeed, suppose $\sum d_i P_i = 0$
 compute at x_k :

$$0 + 0 + \dots + 0 + d_k P_k(x_k) = 0$$

$$\Rightarrow d_k = 0 \Rightarrow \forall k \, d_k = 0$$

$\dim P_n(F) = n+1 \Rightarrow \beta$ is a basis



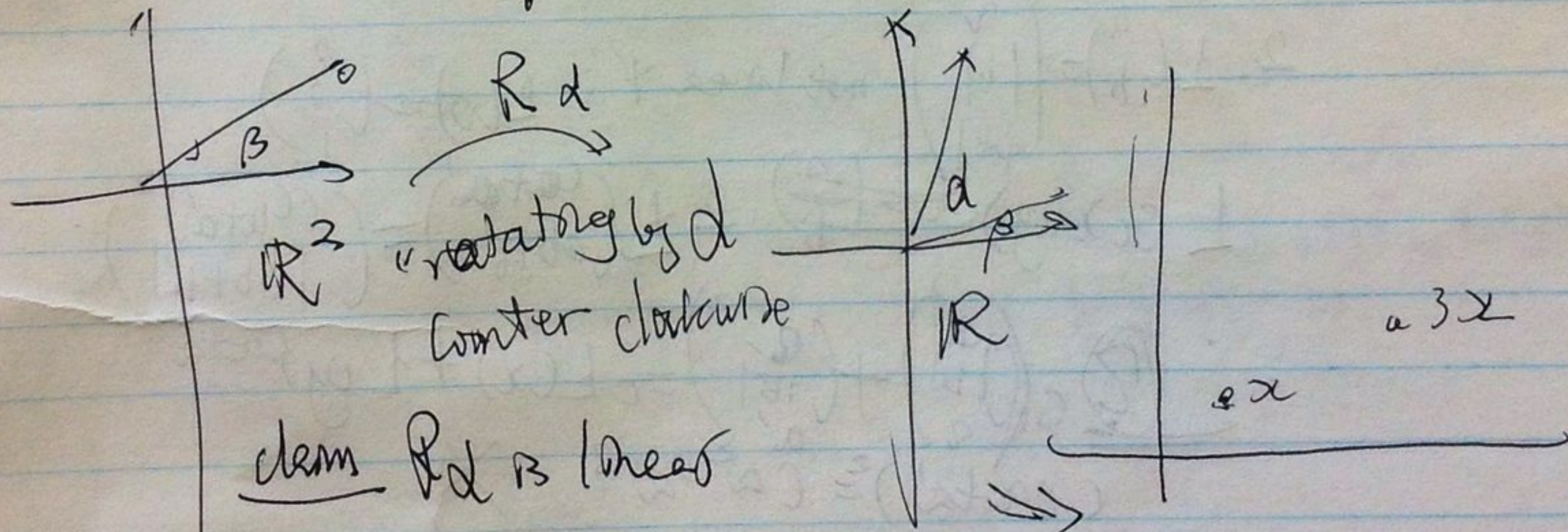
$(F+g)' \xrightarrow{f} F'+g'$
 true by basic calculus Prop 1 holds
Prop 2 $(c f)' = c \cdot f'$

$A \in M_{m \times n}(F)$ $A = (a_{ij}) = \begin{pmatrix} a_{11} & \dots & a_{1n} \\ \vdots & & \vdots \\ a_{m1} & \dots & a_{mn} \end{pmatrix}$
 define $L_A: F^n \rightarrow F^m$ by $x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \mapsto \begin{pmatrix} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n \end{pmatrix}$
 $y = \begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix}$ $y_j = \sum_{i=1}^n a_{ji} x_i$

LHS $L(cx+z)_j = \sum_{i=1}^n a_{ji} (cx_i + z_i) = \sum_{i=1}^n a_{ji} (cx_i + z_i)$

RHS: $[cL(x) + L(z)]_j = c \sum_{i=1}^n a_{ji} x_i + \sum_{i=1}^n a_{ji} z_i$

if LHS $\stackrel{!}{=}$ RHS \downarrow Distributive Law



Prop in pictures $L(R_d(x)) = R_d(L(x))$
 cef

