Crossing the Crossings

Abstract. The subject will be very close to Manturov's represen-Back to \mathcal{H} . The "crossing the crossings" map tation of vB_n into Aut (FG_{n+1}) — I'll describe how I think about it $\mathcal{K}: vT_n \to vT_{n+1}$ is defined by the picture beloin terms of a very simple minded map \mathcal{K} from n-component v- w. Equally well, it is $\mathcal{K}: v\mathcal{B}_n \to w\mathcal{B}_{n+1}$. Better, it is tangles to (n+1)-component w-tangles. It is possible that you all $\mathcal{K}: \mathcal{N}_n \to (nv+1w)T$ or $\mathcal{K}: \mathcal{N}_n \to (nv+1w)B$. know this already. Possibly my talk will be very short — it will Claims. be as long as it is necessary to describe \mathcal{K} and say a few more 1. \mathcal{K} is well defined. words, and if this is little, so be it.



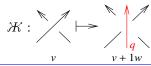
- 2. On u-links, \mathcal{K} "factors".
- **3.** \mathcal{K} does not respect OC.

4. \mathcal{K} recovers Manturov's VG and μ : $VG(K) = \pi_1(\mathcal{K}(K)), \mu = \mathcal{K} \circ \phi = \phi/\!\!/\mathcal{K}$. Even better, \mathcal{K} pulls back any invariant of 2-component w-knots to an invariant of virtual knots. in particular, there is a wheelvalued "non-commutative" invariant ω as in [BN] and DBN: Talks: Hamilton-1412 (next page).

> Likely, the various "2-variable Alexander polynomials" for virtual knots arise in this way.

All you need is $\mathcal{K}...$ • What is its domain? • What is its target?

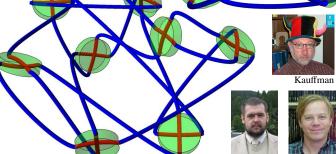
Why should one care?



Virtual Knots. Virtual knots are the algebraic structure underlying the Reidemeister presentation of ordinary knots, without the topology. Locally they are knot diagrams modulo the Reidemeister relations; globally, who cares? So,

$$\nu T = CA \langle \times, \times : R1, R2, R3 \rangle$$

Flying Pogs for $v2_1$ and for 8_{17} :



No! Note that also

$$vT = PA \langle \times, \times, \times : R1,R2,R3,VR1,VR2,VR3,M \rangle,$$

but I have a prejudice, or a deeply held belief, that this is morally off... wrong!

Manturov's $\mu: \nu B_n \to \operatorname{Aut}(F(x_1, \dots, x_n, q))$: [Ma, BGHNW] $\sigma_i = \mathbb{X}_i \mapsto \begin{cases} x_i \mapsto x_i x_{i+1} x_i^{-1} \\ x_{i+1} \mapsto x_i \end{cases} \quad \tau_i = \mathbb{X}_i \mapsto \begin{cases} x_i \mapsto q x_{i+1} q^{-1} \\ x_{i+1} \mapsto q^{-1} x_i q \end{cases}.$

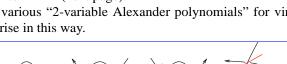
Easy resolution. Setting $y_i := q^i x_i q^{-i}$, we find that μ is equivalent to

$$\sigma_{ij} \mapsto \begin{cases} y_i \mapsto q y_i q^{-1} \\ y_i \mapsto y_i^{-1} q^{-1} y_i q y_i \end{cases}$$

w-Tangles. wT := vT/OC where "Overcrossings Commute" is:

 π_1 is defined on wT; Artin's representation ϕ is defined on wB_n .

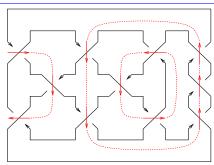
Proof of 1.



Everything slides out!

Proof of 2. The net "red flow" into every face is 0, so the red arrows can be paired. They form cycles that can hover off the picture.

No proof of 3. Well, there simply is no proof that OCis respected, and it's easy to come up with counterexamples.



Proof of 4. A simple verification, except my conventions are

References.

[BN] D. Bar-Natan, Balloons and Hoops and their Universal Finite Type Invariant, BF Theory, and an Ultimate Alexander Invariant, Acta Mathematica Vietnamica 40-2 (2015) 271–329, arXiv:1308.1721.

[BGHNW] H. U. Boden, A. I. Gaudreau, E. Harper, A. J. Ni- $\sigma_{ij} \mapsto \begin{cases} y_i \mapsto qy_i q^{-1} \\ y_j \mapsto y_i^{-1} q^{-1} y_j q y_i \end{cases}$ cas, and L. White, Virtual Knot Groups and Almost Classical Knots, arXiv:1506.01726.

But why does it exist? Especially, wherefore $vB_n \to wB_{n+1}$? [Ma] V. O. Manturov, On Invariants of Virtual Links, Acta Ap-

plicandae Mathematica **72-3** (2002) 295–309.

Prejudices should always be re-evaluated!

