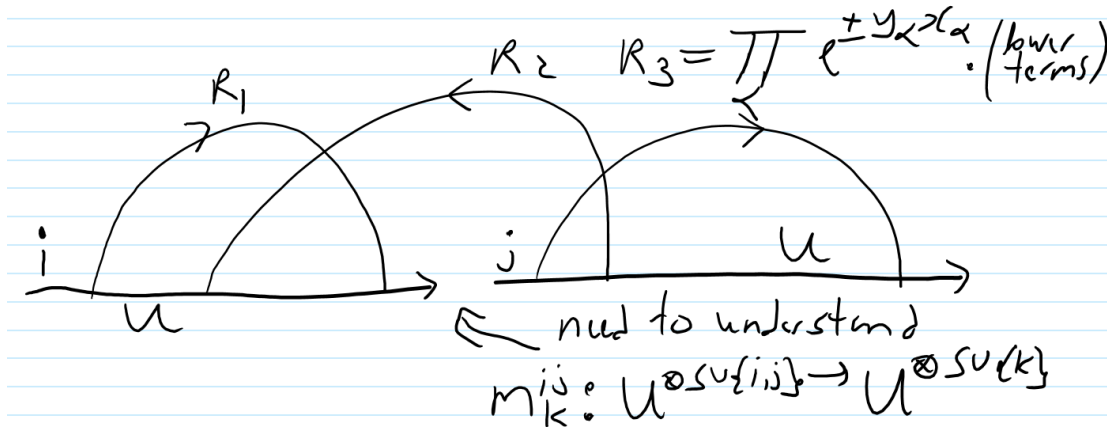


Dror Bar-Natan: Talks: Sydney-1708:



The Dogma is Wrong - Extra Details

Goal



Agenda

1. Quantizing and de-quantizing sl_2^ϵ .
2. Some understanding of sl_2^ϵ .
3. A full understanding of sl_2^ϵ at $\epsilon = 0$.
4. A full understanding of sl_2^ϵ at $\epsilon^2 = 0$.
5. Pushforwards of distributions, 0-dimensional QFT, Feynman diagrams and what had really happened here.

Some Shortcuts

`ME[x_] := MatrixExp[x]; MB[x_, y_] := x.y - y.x; MF[x_] := MatrixForm[x];`

Representing $g^\epsilon = \langle h, e, l, f \rangle / ([e, l] = -e, [f, l] = f, [e, f] = h - 2\epsilon l, [h, *] = 0)$

$$\rho_h = \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}; \rho_e = \begin{pmatrix} 0 & 0 \\ -\epsilon & 0 \end{pmatrix}; \rho_l = \begin{pmatrix} -(1+1/\epsilon)/2 & 0 \\ 0 & (1-1/\epsilon)/2 \end{pmatrix}; \rho_f = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix};$$

`Simplify@{MB[\rho_e, \rho_l] == -\rho_e, MB[\rho_f, \rho_l] == \rho_f, MB[\rho_e, \rho_f] == \rho_h - 2 \epsilon \rho_l}`

The Main $g_0 := g^\epsilon / (\epsilon = 0)$ Theorem.

The g_0 invariant of any S-component tangle T can be written in the form $Z(T) = \mathcal{O}(\omega e^{L+Q} \mid \prod_{i \in S} e_i l_i f_i)$, where ω is a scalar (meaning, a rational function in the variables h_i and their exponentials $t_i = e^{h_i}$), where $L = \sum a_{ij} h_i l_j$ is a balanced quadratic in the variables h_i and l_j with integer coefficients a_{ij} and where $Q = \sum b_{ij} e_i f_j$ is a balanced quadratic in the variables e_i and f_j with scalar coefficients b_{ij} . Furthermore, after setting $h_i = h$ and $t_i = t$ for all i , the invariant $Z(T)$ is poly-time computable.

Proof. Indeed, as we shall see, the following lemmas hold, and the rest is straight-forward.

Lemma 0. $R^S = e^{s(h \otimes l + e \otimes f)} = \mathcal{O}(\exp(s h l + \frac{e^{sh}-1}{h} e f \mid e \otimes l f)$.

Lemma 1. $\mathcal{O}(e^{\gamma l + \beta e} \mid l e) = \mathcal{O}(e^{\gamma l + e^\gamma \beta e} \mid e l)$.

Lemma 2. $\mathcal{O}(e^{\gamma l + \beta f} \mid f l) = \mathcal{O}(e^{\gamma l + e^\gamma \beta f} \mid l f)$.

Lemma 3. $\mathcal{O}(e^{\beta e + \alpha f + \delta e f} \mid f e) = \mathcal{O}(v e^{\nu(-\alpha \beta h + \beta e + \alpha f + \delta e f)} \mid e f)$, with $v = (1 + h \delta)^{-1}$.

Some g^ϵ lemmas

Lemma 1. $\mathcal{O}(e^{\gamma l + \beta e} \mid l e) = \mathcal{O}(e^{\gamma l + e^\gamma \beta e} \mid e l)$.

Lemma 2. $\mathcal{O}(e^{\gamma l + \beta f} \mid f l) = \mathcal{O}(e^{\gamma l + e^\gamma \beta f} \mid l f)$.

Proofs.

`MF /@ {ME[γ ρ l].ME[β ρ e], ME[e^γ β ρ e].ME[γ ρ l]}`

`MF /@ {ME[β ρ f].ME[γ ρ l], ME[γ ρ l].ME[e^γ β ρ f]}`

Lemma 3 at $\delta = 0$. $\mathcal{O}(e^{\alpha f + \beta e} \mid f e) = \mathcal{O}(e^{c h + a e - 2 \epsilon c l + b f} \mid e l f)$, with

$\left\{ a \rightarrow -\frac{\beta}{-1 + \alpha \beta \epsilon}, b \rightarrow -\frac{\alpha}{-1 + \alpha \beta \epsilon}, c \rightarrow \frac{\text{Log}[1 - \alpha \beta \epsilon]}{\epsilon} \right\}$.

Derivation.

`ME[α ρ f].ME[β ρ e] // Simplify // MF`

`eqn = ME[α ρ f].ME[β ρ e] == ME[a ρ e].ME[c (ρ h - 2 ε ρ l)].ME[b ρ f]`

`sol = Solve[Thread[Flatten /@ eqn], {a, b, c}] [[1]]`

`sol = sol /. C[1] → 0`

Lemma 3 for g_0 .

`Limit[{a, b, c} /. sol, ε → 0]`

And so in g_0 , $\mathcal{O}(e^{\alpha f + \beta e} \mid f e) = \mathcal{O}(e^{\alpha f + \beta e - \alpha \beta h} \mid e l f)$. Hence

$\mathcal{O}(e^{\alpha f + \beta e + \delta e f} \mid f e) = e^{\delta \partial_\alpha \partial_\beta} \mathcal{O}(e^{\alpha f + \beta e} \mid f e) = e^{\delta \partial_\alpha \partial_\beta} \mathcal{O}(e^{\alpha f + \beta e - \alpha \beta h} \mid e l f) = \mathcal{O}(\psi \mid e l f)$, where

$\psi = e^{\delta \partial_\alpha \partial_\beta} e^{\alpha f + \beta e - \alpha \beta h}$ satisfies $\psi_{\delta=0} = e^{\alpha f + \beta e - \alpha \beta h}$ and $\partial_\delta \psi = \partial_{\alpha, \beta} \psi$.

`With[{ψ = v e^ν (δ e f - α β h + α f + β e)} /. v → (1 + δ h)^-1], Simplify@{∂_δ ψ - ∂_{α, β} ψ, ψ /. δ → 0}]`

A Lemma 3 for $g_k := g^\epsilon / (\epsilon^{k+1} = 0)$.

Lemma 3_k. $\mathcal{O}(e^{\beta e + \alpha f + \delta e f} \mid f e) = \mathcal{O}(v e^{v(-\alpha \beta h + \beta e + \alpha f + \delta e f)} \wedge_k(\epsilon, e, l, f, \alpha, \beta, \delta) \mid e f)$, with $v = (1 + h \delta)^{-1}$ and where for any fixed k , $\wedge_k(\epsilon, e, l, f, \alpha, \beta, \delta)$ is a fixed polynomial of degree at most $4k$ in $e, \sqrt{l}, f, \alpha, \beta$, with scalar coefficients.

Comment. Even better, $\log(\wedge_k)$ is of degree at most $2k + 2$ in said variables.

Comment. And hence the g_k invariant is computable in polynomial time.

Proof of Lemma 3_k. We know that $\mathcal{O}(e^{\alpha f + \beta e} \mid f e) = \mathcal{O}(e^{c h + a e - 2 \epsilon c l + b f} \mid e l f)$, with

$\left\{ a \rightarrow -\frac{\beta}{-1 + \alpha \beta \epsilon}, b \rightarrow -\frac{\alpha}{-1 + \alpha \beta \epsilon}, c \rightarrow \frac{\log[1 - \alpha \beta \epsilon]}{\epsilon} \right\}$. Expanding in ϵ we get

$\mathcal{O}(e^{\alpha f + \beta e} \mid f e) = \mathcal{O}(\lambda_\epsilon(\alpha, \beta) e^{\alpha f + \beta e - \alpha \beta h} \mid e l f) = \mathcal{O}(\lambda_\epsilon(\partial_f, \partial_e) e^{\alpha f + \beta e - \alpha \beta h} \mid e l f)$ and so

$\mathcal{O}(e^{\alpha f + \beta e + \delta e f} \mid f e) = \mathcal{O}(\lambda_\epsilon(\partial_f, \partial_e) e^{\delta \partial_\alpha \partial_\beta} e^{\alpha f + \beta e - \alpha \beta h} \mid e l f) = \mathcal{O}(\lambda_\epsilon(\partial_f, \partial_e) v e^{v(-\alpha \beta h + \beta e + \alpha f + \delta e f)} \mid e l f)$.

$DP_{\alpha \rightarrow D_f, \beta \rightarrow D_e}[P_][\lambda_]:=$

Total[CoefficientRules[P, {α, β}]] /. ({m_, n_} → c_) := c D[λ, {f, m}, {e, n}]]

(* "D" for Detailed *)

$D\Delta_r[h_, e_, l_, f_, \alpha_, \beta_, \delta_] := \text{Module}$

{ρh, ρe, ρl, ρf, eqn, a, b, c, sol, λ, q, v},

$\rho h = \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}; \rho e = \begin{pmatrix} 0 & 0 \\ -\epsilon & 0 \end{pmatrix}; \rho l = \begin{pmatrix} -(1 + 1/\epsilon)/2 & 0 \\ 0 & (1 - 1/\epsilon)/2 \end{pmatrix}; \rho f = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix};$

$eqn = \text{ME}[\alpha \rho f] . \text{ME}[\beta \rho e] == \text{ME}[a \rho e] . \text{ME}[c (\rho h - 2 \epsilon \rho l)] . \text{ME}[b \rho f];$

$sol = \text{Solve}[\text{Thread}[\text{Flatten}@\text{eqn}], \{a, b, c\}] [[1]] /. C[1] \rightarrow 0;$

$\lambda = \text{Simplify}[e^{-f \alpha - e \beta + h \alpha \beta} \text{Normal}@\text{Series}[e^{c h + a e - 2 \epsilon c l + b f} /. sol, \{\epsilon, 0, k\}]]];$

$q = e^{v (f \alpha + e \beta - h \alpha \beta + e f \delta)};$

$\text{Collect}[q^{-1} DP_{\alpha \rightarrow D_f, \beta \rightarrow D_e}[\lambda][q] /. v \rightarrow (1 + h \delta)^{-1}, \epsilon, \text{Simplify}]$

];

$D\Delta_1[h, e, l, f, \alpha, \beta, \delta]$

$D\Delta_2[h, e, l, f, \alpha, \beta, \delta]$

$\Delta_k[h_, e_, l_, f_, \alpha_, \beta_, \delta_] := \Delta_k[h, e, l, f, \alpha, \beta, \delta] = \text{Module}[\{\lambda\},$

$\lambda = \text{Normal}@\text{Series}[e^{\frac{f \alpha + e \beta}{1 - \alpha \beta \epsilon}} (1 - \alpha \beta \epsilon)^{-2 L + \frac{h}{\epsilon}}, \{\epsilon, 0, k\}] /. e \rightarrow 1;$

$\text{Collect}[DP_{\alpha \rightarrow D_f, \beta \rightarrow D_e}[\lambda][e^{(f \alpha + e \beta + e f \delta) / (1 + h \delta)}] /. e \rightarrow 1, \epsilon, \text{Simplify}]]];$

$\text{Simplify}[D\Delta_2[h, e, l, f, \alpha, \beta, \delta] == \Delta_2[h, e, l, f, \alpha, \beta, \delta]]$

The Main g_k Theorem

The g_k invariant of any S-component tangle T can be written in the form $Z(T) = \mathcal{O}(\omega e^{L+Q+P} \mid \prod_{i \in S} e_i l_i f_i)$, where ω is a scalar (meaning, a rational function in the variables h_i and their exponentials $t_i = e^{h_i}$), where $L = \sum a_{ij} h_i l_j$ is a balanced quadratic in the variables h_i and l_j with integer coefficients a_{ij} , where $Q = \sum b_{ij} e_i f_j$ is a balanced quadratic in the variables e_i and f_j with scalar coefficients b_{ij} , and where P is a polynomial in $\{\epsilon, e_i, l_i, f_i\}$ (with scalar coefficients) whose ϵ^d -term is of degree at most $2d + 2$ in

$\{e_i, \sqrt{l_i}, f_i\}$. Furthermore, after setting $h_i = h$ and $t_i = t$ for all i , the invariant $Z(T)$ is poly-time computable.

Partial Proof. Indeed,

0. $R^\pm = ?$, $n^\pm = ?$.

$$1. \mathcal{O}(\mathcal{P}(l, e) e^{Yl+\beta e} \mid l e) = \mathcal{O}(\mathcal{P}(\partial_Y, \partial_\beta) e^{Yl+e^Y \beta e} \mid e l),$$

$$2. \mathcal{O}(\mathcal{P}(l, f) e^{Yl+\beta f} \mid f l) = \mathcal{O}(\mathcal{P}(\partial_Y, \partial_\beta) e^{Yl+e^Y \beta f} \mid l f),$$

$$3. \mathcal{O}(\mathcal{P}(e, f) e^{\beta e+\alpha f+\delta e f} \mid f e) = \mathcal{O}(v \mathcal{P}(\partial_\beta, \partial_\alpha) e^{v(-\alpha \beta h+\beta e+\alpha f+\delta e f)} \Lambda_k(\epsilon, e, l, f, \alpha, \beta, \delta) \mid e f),$$

with $v = (1 + h\delta)^{-1}$, and $\Lambda_k(\epsilon, e, l, f, \alpha, \beta, \delta)$ as above.

Pushforwards of distributions, 0-dimensional QFT, Feynman diagrams and what had really happened here.