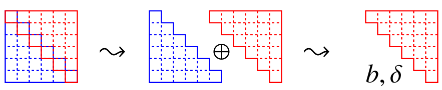
What else can you do with solvable approximations?

Thanks for the invitation!

Abstract. Recently, Roland van der Veen and myself found that there are sequences of solvable Lie algebras "converging" to any given semi-simple Lie algebra (such as sl_2 or sl_3 or $E8$). Certain computations are much easier in solvable Lie algebras; in particular, using solvable approximations we can compute in polynomial time certain projections (originally discussed by Rozansky) of the knot invariants arising from the Chern-Simons-Witten topological quantum field theory. This provides us with the first strong knot invariants that are computable for truly large knots.

But sl_2 and sl_3 and similar algebras occur in physics (and in mathematics) in many other places, beyond the Chern-Simons-Witten theory. Do solvable approximations have further applications?

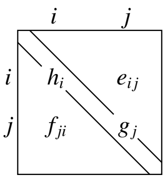
Recomposing gl_n. Half is enough! $gl_n \oplus \mathfrak{a}_n = \mathcal{D}(◹, b, \delta)$:



$b(◹) = b\colon ◹ \otimes ◹ \to ◹$

$b(◺) \rightsquigarrow \delta\colon ◹ \to ◹ \otimes ◹$

Now define $gl_n^\epsilon := \mathcal{D}(◹, b, \epsilon\delta)$. Schematically, this is $[◹, ◹] = ◹$, $[◺, ◺] = \epsilon ◺$, and $[◹, ◺] = ◺ + \epsilon ◹$. In detail, it is



$$[e_{ij}, e_{kl}] = \delta_{jk} e_{il} - \delta_{li} e_{kj} \qquad [f_{ij}, f_{kl}] = \epsilon\delta_{jk} f_{il} - \epsilon\delta_{li} f_{kj}$$

$$[e_{ij}, f_{kl}] = \delta_{jk}(\epsilon\delta_{j<k} e_{il} + \delta_{il}(h_i + \epsilon g_i)/2 + \delta_{i>l} f_{il}) - \delta_{li}(\epsilon\delta_{k<j} e_{kj} + \delta_{kj}(h_j + \epsilon g_j)/2 + \delta_{k>j} f_{kj})$$

$$[g_i, e_{jk}] = (\delta_{ij} - \delta_{ik}) e_{jk} \qquad [h_i, e_{jk}] = \epsilon(\delta_{ij} - \delta_{ik}) e_{jk}$$

$$[g_i, f_{jk}] = (\delta_{ij} - \delta_{ik}) f_{jk} \qquad [h_i, f_{jk}] = \epsilon(\delta_{ij} - \delta_{ik}) f_{jk}$$

Solvable Approximation. At $\epsilon = 1$ and modulo $h = g$, the above is just gl_n. By rescaling at $\epsilon \neq 0$, gl_n^ϵ is independent of ϵ. We let gl_n^k be gl_n^ϵ regarded as an algebra over $\mathbb{Q}[\epsilon]/\epsilon^{k+1} = 0$. It is the "$k$-smidgen solvable approximation" of gl_n!

Recall that $\mathfrak{g}$ is "solvable" if iterated commutators in it ultimately vanish: $\mathfrak{g}_2 := [\mathfrak{g}, \mathfrak{g}]$, $\mathfrak{g}_3 := [\mathfrak{g}_2, \mathfrak{g}_2]$, ..., $\mathfrak{g}_d = 0$. Equivalently, if it is a subalgebra of some large-size ◹ algebra.

Note. This whole process makes sense for arbitrary semi-simple Lie algebras.

Why are "solvable algebras" any good? Contrary to common beliefs, computations in semi-simple Lie algebras are just awful:

```
In[1]:= MatrixExp[( a b
                    c d )] // FullSimplify // MatrixForm     Enter
```

Yet in solvable algebras, exponentiation is fine and even BCH, $z = \log(e^x e^y)$, is bearable:

```
In[2]:= MatrixExp[( a b
                    0 c )] // MatrixForm
```

Out[2]//MatrixForm=
$$\begin{pmatrix} e^a & \frac{b(e^a - e^c)}{a-c} \\ 0 & e^c \end{pmatrix}$$

```
In[3]:= MatrixExp[( a1 b1
                    0  c1 )].MatrixExp[( a2 b2
                                         0  c2 )] //
          MatrixLog // PowerExpand // Simplify //
        MatrixForm                                          Enter
```

Question. What else can you do with solvable approximation? Chern-Simons-Witten theory is often "solved" using ideas from conformal field theory and using quantization of various moduli spaces. Does it make sense to use solvable approximation there too? Elsewhere in physics? Elsewhere in mathematics?

See Also. Talks at George Washington University [ωεβ/gwu], Indiana [ωεβ/ind], and Les Diablerets [ωεβ/ld], and a University of Toronto "Algebraic Knot Theory" class [ωεβ/akt].

Chern-Simons-Witten. Given a knot $\gamma(t)$ in $\mathbb{R}^3$ and a metrized Lie algebra $\mathfrak{g}$, set $Z(\gamma) :=$

$$\int_{A\in\Omega^1(\mathbb{R}^3,\mathfrak{g})} \mathcal{D}A\, e^{ik\, cs(A)} PExp_\gamma(A),$$

where $cs(A) := \frac{1}{4\pi}\int_{\mathbb{R}^3} \operatorname{tr}\left(A\,dA + \frac{2}{3}A^3\right)$ and

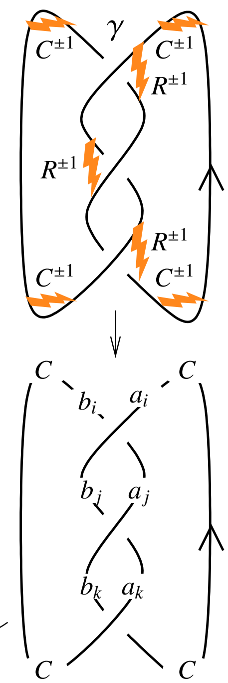
$$PExp_\gamma(A) := \prod_0^1 \exp(\gamma^* A) \in \mathcal{U} = \hat{\mathcal{U}}(\mathfrak{g}),$$

and $\mathcal{U}(\mathfrak{g}) := \langle \text{words in } \mathfrak{g}\rangle/(xy - yx = [x,y])$. In a favourable gauge, one may hope that this computation will localize near the crossings and the bends, and all will depend on just two quantities,

$$R = \sum a_i \otimes b_i \in \mathcal{U} \otimes \mathcal{U} \quad \text{and} \quad C \in \mathcal{U}.$$

This was never done formally, yet R and C can be "guessed" and all "quantum knot invariants" arise in this way. So for the trefoil,

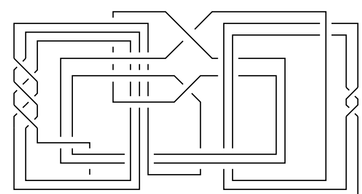
$$Z = \sum_{i,j,k} C a_i b_j a_k C^2 b_i a_j b_k C.$$



But Z lives in $\mathcal{U}$, a complicated space. How do you extract information out of it?

Solution 1, Representation Theory. Choose a finite dimensional representation ρ of $\mathfrak{g}$ in some vector space V. By luck and the wisdom of Drinfel'd and Jimbo, $\rho(R) \in V^* \otimes V^* \otimes V \otimes V$ and $\rho(C) \in V^* \otimes V$ are computable, so Z is computable too. But in exponential time!

Ribbon=Slice?



Gompf, Scharlemann, Thompson

Solution 2, Solvable Approximation. Work directly in $\hat{\mathcal{U}}(\mathfrak{g}_k)$, where $\mathfrak{g}_k = sl_2^k$ (or a similar algebra); everything is expressible using low-degree polynomials in a small number of variables, hence everything is poly-time computable!

Example 0. Take $\mathfrak{g}_0 = sl_2^0 = \mathbb{Q}\langle h, e, l, f\rangle$, with h central and $[f,l] = f$, $[e,l] = -e$, $[e,f] = h$. In it, using normal orderings,

$$R = \mathbb{O}\left(\exp\left(hl + \frac{e^h - 1}{h} ef\right) \mid e \otimes lf\right), \quad \text{and,}$$

$$\mathbb{O}\left(e^{\delta ef} \mid fe\right) = \mathbb{O}\left(\nu e^{\nu\delta ef} \mid ef\right) \quad \text{with } \nu = (1 + h\delta)^{-1}.$$

Example 1. Take $R = \mathbb{Q}[\epsilon]/(\epsilon^2 = 0)$ and $\mathfrak{g}_1 = sl_2^1 = R\langle h, e, l, f\rangle$, with h central and $[f,l] = f$, $[e,l] = -e$, $[e,f] = h - 2\epsilon l$. In it,

$$\mathbb{O}\left(e^{\delta ef} \mid fe\right) = \mathbb{O}\left(\nu(1 + \epsilon\nu\delta\Lambda/2)e^{\nu\delta ef} \mid elf\right), \quad \text{where } \Lambda \text{ is}$$

$$4\nu^3\delta^2 e^2 f^2 + 3\nu^3\delta^3 h e^2 f^2 + 8\nu^2\delta ef + 4\nu^2\delta^2 hef + 4\nu\delta elf - 2\nu\delta h + 4l.$$

Fact. Setting $h_i = h$ (for all i) and $t = e^h$, the $\mathfrak{g}_1$ invariant of any tangle T can be written in the form

$$Z_{\mathfrak{g}_1}(T) = \mathbb{O}\left(\omega^{-1} e^{hL + \omega^{-1}Q}(1 + \epsilon\omega^{-4}P) \mid \bigotimes_i e_i l_i f_i\right),$$

where L is linear, Q quadratic, and P quartic in the $\{e_i, l_i, f_i\}$ with ω and all coefficients polynomials in t. Furthermore, everything is poly-time computable.

On Elves and Invariants

Follows Rozansky [Ro1, Ro2, Ro3] and Overbay [Ov], joint with van der Veen.

Abstract. Whether or not you like the formulas on this page, they describe the strongest truly computable knot invariant we know.

Three steps to the computation of ρ_1:

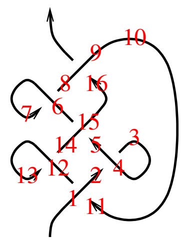
1. Preparation. Given K, results $\langle$*long word* $\|$ *simple formulas*$\rangle$.

2. Rewrite rules. Make the *word* simpler and the *formulas* more complicated, until the *word* "*elf*" is reached.

3. Readout. The invariant ρ_1 is read from the last *formulas*.

$$\begin{array}{c}\text{Knot } K\\ \downarrow \text{ preparation}\\ \langle elf\dots elf \,\|\, \omega_0; L_0; Q_0; P_0\rangle\\ \downarrow \text{ rewrite rules}\\ \langle elf \,\|\, \omega; -; -; P\rangle\\ \downarrow \text{ readout}\\ \rho_1(K)=\rho_1(\omega,P)\end{array}$$

Preparation. Draw K using a 0-framed 0-rotation planar diagram D where all crossings are pointing up. Walk along D labeling features by $1,\dots,m$ in order: over-passes, under-passes, and right-heading cups and caps ("±-cuaps"). If x is a xing, let i_x and j_x be the labels on its over/under strands, and let s_x be 0 if it right-handed and -1 otherwise. If c is a cuap, let i_c be its label and s_c be its sign. Set



$$(L;Q;P)=\sum_{x:\,(i,j,s)}(-)^s\left(l_j;\ t^s e_i f_j;\ (-t)^s e_i l_{(1+s)i-sj} f_j + l_i l_j + \frac{t^{2s}e_i^2 f_j^2}{4}\right) + \sum_{c:\,(i,s)}(0;\ 0;\ s\cdot l_i)\,.$$

This done, output $\langle e_1 l_1 f_1 e_2 l_2 f_2 \cdots e_m l_m f_m \,\|\, 1; L; Q; P\rangle$.

In *formulas*, L is always $\mathbb{Z}$-linear in $\{l_i\}$, Q is an R-linear combination of $\{e_i f_j\}$ where $R := \mathbb{Q}[t^{\pm 1}]$, and P is an R-linear combination of $\{1,\ l_i,\ l_i l_j,\ e_i f_j,\ e_i l_j f_k,\ e_i e_j f_k f_l\}$. (The key to computability!)

Rewrite Rules. Manipulate $\langle$*word* $\|$ *formulas*$\rangle$ expressions using the rewrite rules below, until you come to the form $\langle e_1 l_1 f_1 \,\|\, \omega; -; -; P\rangle$. Output (ω, P).

Rule 1, Deletions. If a letter appears in *word* but not in *formulas*, you can delete it.

Rule 2, Merges. In *word*, you can replace *adjacent* $v_i v_j$ with v_k (for $v \in \{e,l,f\}$) while making the same changes in *formulas* (provided k creates no naming clashes). E.g.,

$$\langle \dots e_i e_j \dots \,\|\, Z\rangle \to \langle \dots e_k \dots \,\|\, Z|_{e_i,e_j\to e_k}\rangle.$$

Rule 3, *le* Sorts. Provided k introduces no clashes, given $\langle \dots l_j e_i \dots \,\|\, \omega; L; Q; P\rangle$, decompose $L = \lambda l_j + L'$, $Q = \alpha e_i + Q'$, write $P = P(e_i, l_j)$ (with messy coefficients), set $q = \mathbb{e}^{\gamma}\beta e_k + \gamma l_k$, and output

$$\left\langle \dots e_k l_k \dots \,\|\, \omega; L|_{l_j\to l_k}; t^{\lambda}\alpha e_k + Q'; \mathbb{e}^{-q}P(\partial_\beta,\partial_\gamma)\mathbb{e}^{q}|_{\beta\to\alpha/\omega,\,\gamma\to\lambda\log t}\right\rangle.$$

Rule 4, *fl* Sorts. Provided k introduces no clashes, given $\langle \dots f_i l_j \dots \,\|\, \omega; L; Q; P\rangle$, decompose $L = \lambda l_j + L'$, $Q = \alpha f_i + Q'$, write $P = P(f_i, l_j)$ (with messy coefficients), set $q = \mathbb{e}^{\gamma}\beta f_k + \gamma l_k$, and output

$$\left\langle \dots l_k f_k \dots \,\|\, \omega; L|_{l_j\to l_k}; t^{\lambda}\alpha f_k + Q'; \mathbb{e}^{-q}P(\partial_\beta,\partial_\gamma)\mathbb{e}^{q}|_{\beta\to\alpha/\omega,\,\gamma\to\lambda\log t}\right\rangle.$$

Happy Birthday, Scott!

"God created the knots, all else in topology is the work of mortals."

Leopold Kronecker (modified)

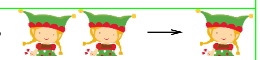
Rule 5, *fe* Sorts. Provided k introduces no clashes, given $\langle \dots f_i e_j \dots \,\|\, \omega; L; Q; P\rangle$, decompose $Q = Q_{fe} f_i e_j + Q_f f_i + Q_e e_j + Q'$ write $P = P(f_i, e_j)$ (with messy coefficients), set $\mu = 1+(t-1)\delta$ and $q = ((1-t)\alpha\beta + \beta e_k + \alpha f_k + \delta e_k f_k)/\mu$, and output

$$\left\langle \dots e_k f_k \dots \,\middle\|\, \begin{array}{c}\mu\omega; L;\ \mu\omega q + \mu Q';\\ \omega^4\Lambda_k + \mathbb{e}^{-q}P(\partial_\alpha,\partial_\beta)(\mathbb{e}^{q})\end{array}\right\rangle\Bigg|_{\substack{\alpha\to Q_f/\omega,\,\beta\to Q_e/\omega,\\ \delta\to Q_{fe}/\omega}},$$

where Λ_k is the Λόγος, "a principle of order and knowledge":

$$\begin{aligned}\Lambda_k = \frac{t+1}{4}\Big(&-\delta(\mu+1)\left(\beta^2 e_k^2 + \alpha^2 f_k^2\right) - \delta^3(3\mu+1)e_k^2 f_k^2\\ &- 2\left(\beta e_k + \alpha f_k\right)\left(\alpha\beta + 2\delta\mu + \delta^2(2\mu+1)e_k f_k + 2\delta\mu^2 l_k\right)\\ &- 4(\alpha\beta+\delta\mu)\left(\delta(\mu+1)e_k f_k + \mu^2 l_k\right) - 4\delta^2\mu^2 e_k f_k l_k\\ &+ (t-1)\left(2(\alpha\beta+\delta\mu)^2 - \alpha^2\beta^2\right)\Big).\end{aligned}$$

elf merges, m^{ij}_k, are defined as compositions



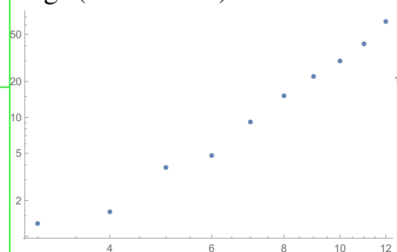
$$e_i l_i \overline{f_i e_j} l_j f_j \xrightarrow{S_x^{f_i e_j}} e_i \overline{l_i e_x}\,\overline{f_x l_j} f_j \xrightarrow{S_x^{l_i e_x} /\!\!/ S_x^{f_x l_j}} \overline{e_i e_x}\,\overline{l_x l_x}\,\overline{f_x f_j} \xrightarrow{i,j,x\to k} e_k l_k f_k$$

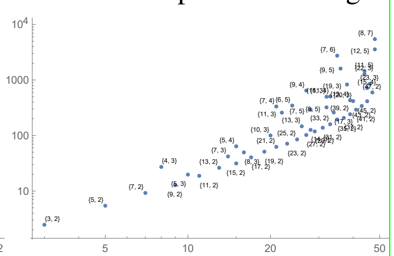
Readout. Given $\langle elf \,\|\, \omega; -; -; P\rangle$, output

$$\rho_1(K) := \frac{t(P|_{e,l,f\to 0} - t\omega'\omega^3)}{(t-1)^2\omega^2}.$$

(ω is the Alexander polynomial, L and Q are not interesting).

Experimental Analysis (ωεβ/Exp). Log-log plots of computation time (sec) vs. crossing number, for all knots with up to 12 crossings (mean times) and for all torus knots with up to 48 crossings:





Power. On the 250 knots with at most 10 crossings, the pair (ω, ρ_1) attains 250 distinct values, while (Khovanov, HOMFLY-PT) attains only 249 distinct values. To 11 crossings the numbers are (802, 788, 772) and to 12 they are (2978, 2883, 2786).

Genus. Up to 12 xings, always ρ_1 is symmetric under $t \leftrightarrow t^{-1}$. With ρ_1^+ denoting the positive-degree part of ρ_1, always $\deg \rho_1^+ \le 2g-1$, where g is the 3-genus of K (equallity for 2530 knots). This gives a lower bound on g in terms of ρ_1 (conjectural, but undoubtedly true). This bound is often weaker than the Alexander bound, yet for 10 of the 12-xing Alexander failures it does give the right answer.

Why Works? The Lie algebra $\mathfrak{g}_1$ (below) is a "solvable approximation of sl_2".

Theorem. The map (as defined below)

$$\langle w \,\|\, \omega; L; Q; P\rangle \mapsto \mathbb{O}\left(\omega^{-1}\mathbb{e}^{L\log t + \omega^{-1}Q}(1+\epsilon\omega^{-4}P) \colon\ w\right) \in \hat{\mathcal{U}}(\mathfrak{g}_1)$$

is well defined modulo the sorting rules. It maps the initial preparation to a product of "*R*-matrices" and "cuap values" satisfying the usual moves for Morse knots (R3, etc.). (And hence the result is a "quantum invariant", except computed very differently; no representation theory!).

1-Smidgen sl_2 Let $\mathfrak{g}_1$ be the 4-dimensional Lie algebra $\mathfrak{g}_1 = \langle h, e', l, f\rangle$ over the ring $R = \mathbb{Q}[\epsilon]/(\epsilon^2 = 0)$, with h central and with $[f, l] = f$, $[e', l] = -e'$, and $[e', f] = h - 2\epsilon l$. Over $\mathbb{Q}$, $\mathfrak{g}_1$ is a solvable approximation of sl_2: $\mathfrak{g}_1 \supset \langle h, e', f, \epsilon h, \epsilon e', \epsilon l, \epsilon f\rangle \supset \langle h, \epsilon h, \epsilon e', \epsilon l, \epsilon f\rangle \supset 0$. Pragmatics: declare $\deg(h, e', l, f, \epsilon) = (1, 1, 0, 0, 1)$ and set $t := \mathbb{e}^h$ and $e := (t-1)e'/h$.

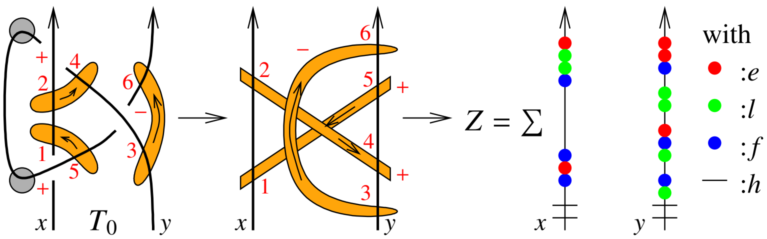
How did it arise? $sl_2 = \mathfrak{b}^+ \oplus \mathfrak{b}^-/\mathfrak{h} =: sl_2^+/\mathfrak{h}$, where $\mathfrak{b}^+ = \langle l, f\rangle/[f, l] = f$ is a Lie bialgebra with $\delta\colon \mathfrak{b}^+ \to \mathfrak{b}^+ \otimes \mathfrak{b}^+$ by $\delta\colon (l, f) \mapsto (0, l\wedge f)$. Going back, $sl_2^+ = \mathcal{D}(\mathfrak{b}^+) = (\mathfrak{b}^+)^* \oplus \mathfrak{b}^+ = \langle h', e', l, f\rangle/\cdots$. **Idea.** Replace $\delta \to \epsilon\delta$ over $\mathbb{Q}[\epsilon]/(\epsilon^{k+1} = 0)$. At $k = 1$, get $[f, l] = f$, $[f, h'] = -\epsilon f$, $[l, e'] = e'$, $[h', e'] = -\epsilon e'$, $[h', l] = 0$, and $[e', f] = h' - \epsilon l$. Now note that $h' + \epsilon l$ is central, so switch to $h := h' + \epsilon l$. This is $\mathfrak{g}_1$.

Ordering Symbols. $\mathbb{O}(poly \mid specs)$ plants the variables of *poly* in $\hat{\mathcal{S}}(\oplus_i\mathfrak{g})$ along $\hat{\mathcal{U}}(\mathfrak{g})$ according to *specs*. E.g.,

$$\mathbb{O}\left(e_1\mathbb{e}^{e_3}l_1^3 l_2 f_3^9 \mid f_3 l_1 e_1 e_3 l_2\right) = f^9 l^3 e\mathbb{e}^e l \in \hat{\mathcal{U}}(\mathfrak{g}).$$

This enables the description of elements of $\hat{\mathcal{U}}(\mathfrak{g})$ using commutative polynomials / power series. In $\mathfrak{g}_1$, no need to specify h / t.

Algebras and Invariants. Given any unital algebra A (even better if A is Hopf; typically, $A \sim \hat{\mathcal{U}}(\mathfrak{g})$), appropriate orange $R \in A\otimes A$, and appropriate cuaps $\in A$, get an $A^{\otimes S}$-valued invariant of pure S-component tangles:



What we didn't say (more, including videos, in ωεβ/Talks).

- ρ_1 is "line" in the coloured Jones polynomial; related to Melvin-Morton-Rozansky.
- ρ_1 extends to "rotational virtual tangles" and is a projection of the universal finite type invariant of such.
- ρ_1 seems to have a better chance than anything else we know to detect a counterexample to slice=ribbon.
- ρ_1 leads to many questions and a very long to-do list. Years of work, many papers ahead. Have fun!

Demo Programs. ωεβ/Demo

Formatting (prints differ ☹)

```
CF[ℰ_] := Module[{vars = Union@Cases[ℰ, e_ | l_ | f_, ∞]},
   If[vars === {}, Factor[ℰ],
    Total[CoefficientRules[ℰ, vars] /.
      (p_ → c_) :> Factor[c] Times @@ (vars^p)]]];
CF[ℰ_𝔼] := CF /@ ℰ;
```

Preparation

```
𝔼[i_, j_, s_] := 𝔼[1, (-1)^s l_j, (-t)^s e_i f_j,
   t^s e_i l_(1+s) i-s j f_j + (-1)^s l_i l_j + (-t^2)^s e_i^2 f_j^2 / 4];
𝔼[i_, s_] := 𝔼[1, 0, 0, s l_i];
𝔼 /: 𝔼[1, L1_, Q1_, P1_] 𝔼[1, L2_, Q2_, P2_] :=
  𝔼[1, L1 + L2, Q1 + Q2, P1 + P2];
```

Preparing the Trefoil

```
z1 = (𝔼[1, 11, 0] 𝔼[4, 2, -1] 𝔼[15, 5, 0]
   𝔼[6, 8, -1] 𝔼[9, 16, 0] 𝔼[12, 14, -1] 𝔼[3, -1] 𝔼[7, +1]
   𝔼[10, -1] 𝔼[13, +1])
```

$$\mathbb{E}\Big[1,\ -l_2 + l_5 - l_8 + l_{11} - l_{14} + l_{16},\ -\tfrac{e_4 f_2}{t} + e_{15} f_5 - \tfrac{e_6 f_8}{t} + e_1 f_{11} - \tfrac{e_{12} f_{14}}{t} + e_9 f_{16},\ -\tfrac{e_4^2 f_2^2}{4t^2} + \tfrac14 e_{15}^2 f_5^2 - \tfrac{e_6^2 f_8^2}{4t^2} + \tfrac14 e_1^2 f_{11}^2 - \tfrac{e_{12}^2 f_{14}^2}{4t^2} + \tfrac14 e_9^2 f_{16}^2 + e_1 f_{11} l_1 + \tfrac{e_4 f_2 l_2}{t} - l_3 - l_2 l_4 + l_7 + \tfrac{e_6 f_8 l_8}{t} - l_6 l_8 + e_9 f_{16} l_9 - l_{10} + l_1 l_{11} + l_{13} + \tfrac{e_{12} f_{14} l_{14}}{t} - l_{12} l_{14} + e_{15} f_5 l_{15} + l_5 l_{15} + l_9 l_{16}\Big]$$

Differential Polynomials

(Implementing $P(\partial_\alpha, \partial_\beta)(f)$)

```
DP_{x→D_α, y→D_β}[P_][f_] :=
 Total[CoefficientRules[P, {x, y}] /.
    ({m_, n_} → c_) :> c D[f, {α, m}, {β, n}]]
```

le and *fl* Sorts

```
S_{1_j_ (x:e|f)_i_ → k_}[𝔼[ω_, L_, Q_, P_]] :=
 With[{λ = ∂_{1_j} L, α = ∂_{x_i} Q, q = e^γ β x_k + γ l_k}, CF[
   𝔼[ω, L /. l_j → l_k, t^λ α x_k + (Q /. x_i → 0),
    e^-q DP_{1_j→D_γ, x_i→D_β}[P][e^q] /. {β → α / ω, γ → λ Log[t]}] ]];
```

The Λόγος

```
Λ[k_] := ((t - 1) (2 (α β + δ μ)^2 - α^2 β^2) - 4 e_k l_k f_k δ^2 μ^2 -
    δ (1 + μ) (f_k^2 α^2 + e_k^2 β^2) - e_k^2 f_k^2 δ^3 (1 + 3 μ) -
    2 (α β + 2 δ μ + e_k f_k δ^2 (1 + 2 μ) + 2 l_k δ μ^2) (f_k α + e_k β) -
    4 (l_k μ^2 + e_k f_k δ (1 + μ)) (α β + δ μ)) (1 + t) / 4;
```

fe Sorts

```
S_{f_i_ e_j_ → k_}[𝔼[ω_, L_, Q_, P_]] :=
 With[{q = ((1 - t) α β + β e_k + α f_k + δ e_k f_k) / μ}, CF[
   𝔼[μ ω, L, μ ω q + μ (Q /. f_i | e_j → 0),
     μ^4 e^-q DP_{f_i→D_α, e_j→D_β}[P][e^q] + ω^4 Λ[k]] /. μ → 1 + (t - 1) δ /.
    {α → ω^-1 (∂_{f_i} Q /. e_j → 0), β → ω^-1 (∂_{e_j} Q /. f_i → 0),
     δ → ω^-1 ∂_{f_i,e_j} Q}]];
```

Elf Merges

```
m_{i_,j_→k_}[Z_𝔼] := Module[{x, z},
  CF[(Z // S_{f_i e_j→x} // S_{l_i e_x→x} // S_{f_x l_j→x}) /. z_(i|j|x) → z_k]]
```

Rewriting the Trefoil

(by merging 16 elves)

```
(Do[z1 = z1 // m_{1,k→1}, {k, 2, 16}]; z1)
```

$$\mathbb{E}\Big[\tfrac{1-t+t^2}{t},\ 0,\ 0,\ \tfrac{(-1+t)(1-t+t^2)^2(1-t+2t^2)}{t^3} - \tfrac{2(1+t)(1-t+t^2)^3 e_1 f_1}{t^4} - \tfrac{2(-1+t)(1+t)(1-t+t^2)^3 l_1}{t^4}\Big]$$

Readout

```
ρ1[𝔼[ω_, _, _, P_]] := CF[ t ((P /. e_ | l_ | f_ → 0) - t ω^3 (∂_t ω)) / ((t - 1)^2 ω^2) ]
```

$\rho_1(3_1)$

```
ρ1[z1] // Expand
```

$\frac{1}{t} + t$

References.

[Ov] A. Overbay, *Perturbative Expansion of the Colored Jones Polynomial,* University of North Carolina PhD thesis, ωεβ/Ov.

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[Ro2] L. Rozansky, *The Universal R-Matrix, Burau Representation and the Melvin-Morton Expansion of the Colored Jones Polynomial,* Adv. Math. **134-1** (1998) 1–31, arXiv:q-alg/9604005.

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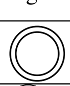
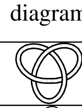
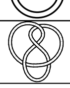
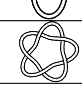
diagram	n_k^t / Today's / Rozansky's ρ_1^+	Alexander's ω^+	genus / ribbon; unknotting number / amphicheiral	diagram	n_k^t / Today's / Rozansky's ρ_1^+	Alexander's ω^+	genus / ribbon; unknotting number / amphicheiral
	0_1^a	1	0 / ✔		3_1^a	$t-1$	1 / ✘
	0		0 / ✔		t		1 / ✘
	4_1^a	$3-t$	1 / ✘		5_1^a	t^2-t+1	2 / ✘
	0		1 / ✔		$2t^3+3t$		2 / ✘

diagram	n_k^t	Alexander's ω^+ Today's / Rozansky's ρ_1^+	genus / ribbon unknotting number / amphicheiral	diagram	n_k^t	Alexander's ω^+ Today's / Rozansky's ρ_1^+	genus / ribbon unknotting number / amphicheiral
	5_2^a	$2t-3$ $5t-4$	1 / ✘ 1 / ✘		6_1^a	$5-2t$ $t-4$	1 / ✔ 1 / ✘
	6_2^a	$-t^2+3t-3$ t^3-4t^2+4t-4	2 / ✘ 1 / ✘		6_3^a	t^2-3t+5 0	2 / ✘ 1 / ✔
	7_1^a	t^3-t^2+t-1 $3t^5+5t^3+6t$	3 / ✘ 3 / ✘		7_2^a	$3t-5$ $14t-16$	1 / ✘ 1 / ✘
	7_3^a	$2t^2-3t+3$ $-9t^3+8t^2-16t+12$	2 / ✘ 2 / ✘		7_4^a	$4t-7$ $32-24t$	1 / ✘ 2 / ✘
	7_5^a	$2t^2-4t+5$ $9t^3-16t^2+29t-28$	2 / ✘ 2 / ✘		7_6^a	$-t^2+5t-7$ $t^3-8t^2+19t-20$	2 / ✘ 1 / ✘
	7_7^a	t^2-5t+9 $8-3t$	2 / ✘ 1 / ✘		8_1^a	$7-3t$ $5t-16$	1 / ✘ 1 / ✘
	8_2^a	$-t^3+3t^2-3t+3$ $2t^5-8t^4+10t^3-12t^2+13t-12$	3 / ✘ 2 / ✘		8_3^a	$9-4t$ 0	1 / ✘ 2 / ✔
	8_4^a	$-2t^2+5t-5$ $3t^3-8t^2+6t-4$	2 / ✘ 2 / ✘		8_5^a	$-t^3+3t^2-4t+5$ $-2t^5+8t^4-13t^3+20t^2-22t+24$	3 / ✘ 2 / ✘
	8_6^a	$-2t^2+6t-7$ $5t^3-20t^2+28t-32$	2 / ✘ 2 / ✘		8_7^a	t^3-3t^2+5t-5 $-t^5+4t^4-10t^3+12t^2-13t+12$	3 / ✘ 1 / ✘
	8_8^a	$2t^2-6t+9$ $-t^3+4t^2-12t+16$	2 / ✔ 2 / ✘		8_9^a	$-t^3+3t^2-5t+7$ 0	3 / ✔ 1 / ✔
	8_{10}^a	t^3-3t^2+6t-7 $-t^5+4t^4-11t^3+16t^2-21t+20$	3 / ✘ 2 / ✘		8_{11}^a	$-2t^2+7t-9$ $5t^3-24t^2+39t-44$	2 / ✘ 1 / ✘
	8_{12}^a	$t^2-7t+13$ 0	2 / ✘ 2 / ✔		8_{13}^a	$2t^2-7t+11$ $-t^3+4t^2-14t+20$	2 / ✘ 1 / ✘
	8_{14}^a	$-2t^2+8t-11$ $5t^3-28t^2+57t-68$	2 / ✘ 1 / ✘		8_{15}^a	$3t^2-8t+11$ $21t^3-64t^2+120t-140$	2 / ✘ 2 / ✘
	8_{16}^a	t^3-4t^2+8t-9 $t^5-6t^4+17t^3-28t^2+35t-36$	3 / ✘ 2 / ✘		8_{17}^a	$-t^3+4t^2-8t+11$ 0	3 / ✘ 1 / ✔
	8_{18}^a	$-t^3+5t^2-10t+13$ 0	3 / ✘ 2 / ✔		8_{19}^n	t^3-t^2+1 $-3t^5-4t^2-3t$	3 / ✘ 3 / ✘
	8_{20}^n	t^2-2t+3 $4t-4$	2 / ✔ 1 / ✘		8_{21}^n	$-t^2+4t-5$ $t^3-8t^2+16t-20$	2 / ✘ 1 / ✘
	9_1^a	$t^4-t^3+t^2-t+1$ $4t^7+7t^5+9t^3+10t$	4 / ✘ 4 / ✘		9_2^a	$4t-7$ $30t-40$	1 / ✘ 1 / ✘
	9_3^a	$2t^3-3t^2+3t-3$ $-13t^5+12t^4-25t^3+20t^2-32t+24$	3 / ✘ 3 / ✘		9_4^a	$3t^2-5t+5$ $23t^3-28t^2+46t-44$	2 / ✘ 2 / ✘
	9_5^a	$6t-11$ $100-65t$	1 / ✘ 2 / ✘		9_6^a	$2t^3-4t^2+5t-5$ $13t^5-24t^4+45t^3-52t^2+68t-64$	3 / ✘ 3 / ✘
	9_7^a	$3t^2-7t+9$ $23t^3-56t^2+99t-108$	2 / ✘ 2 / ✘		9_8^a	$-2t^2+8t-11$ $3t^3-16t^2+29t-28$	2 / ✘ 2 / ✘
	9_9^a	$2t^3-4t^2+6t-7$ $13t^5-24t^4+55t^3-72t^2+98t-96$	3 / ✘ 3 / ✘		9_{10}^a	$4t^2-8t+9$ $-40t^3+72t^2-114t+120$	2 / ✘ 2, 3 / ✘
	9_{11}^a	$-t^3+5t^2-7t+7$ $-2t^5+16t^4-41t^3+52t^2-66t+64$	3 / ✘ 2 / ✘		9_{12}^a	$-2t^2+9t-13$ $5t^3-36t^2+84t-100$	2 / ✘ 1 / ✘
	9_{13}^a	$4t^2-9t+11$ $-40t^3+92t^2-154t+168$	2 / ✘ 2, 3 / ✘		9_{14}^a	$2t^2-9t+15$ $-t^3+8t^2-35t+60$	2 / ✘ 1 / ✘
	9_{15}^a	$-2t^2+10t-15$ $-5t^3+40t^2-108t+136$	2 / ✘ 2 / ✘		9_{16}^a	$2t^3-5t^2+8t-9$ $-13t^5+36t^4-80t^3+120t^2-161t+168$	3 / ✘ 3 / ✘
	9_{17}^a	t^3-5t^2+9t-9 $t^5-8t^4+23t^3-32t^2+28t-24$	3 / ✘ 2 / ✘		9_{18}^a	$4t^2-10t+13$ $40t^3-108t^2+193t-220$	2 / ✘ 2 / ✘
	9_{19}^a	$2t^2-10t+17$ $t^3-8t^2+20t-24$	2 / ✘ 1 / ✘		9_{20}^a	$-t^3+5t^2-9t+11$ $2t^5-16t^4+47t^3-84t^2+117t-124$	3 / ✘ 2 / ✘
	9_{21}^a	$-2t^2+11t-17$ $-5t^3+44t^2-127t+164$	2 / ✘ 1 / ✘		9_{22}^a	$t^3-5t^2+10t-11$ $-t^5+8t^4-24t^3+38t^2-40t+36$	3 / ✘ 1 / ✘
	9_{23}^a	$4t^2-11t+15$ $40t^3-128t^2+243t-288$	2 / ✘ 2 / ✘		9_{24}^a	$-t^3+5t^2-10t+13$ $-4t^2+16t-20$	3 / ✘ 1 / ✘
	9_{25}^a	$-3t^2+12t-17$ $12t^3-70t^2+153t-188$	2 / ✘ 2 / ✘		9_{26}^a	$t^3-5t^2+11t-13$ $-t^5+8t^4-31t^3+64t^2-85t+92$	3 / ✘ 1 / ✘
	9_{27}^a	$-t^3+5t^2-11t+15$ $t^3-8t^2+24t-32$	3 / ✔ 1 / ✘		9_{28}^a	$t^3-5t^2+12t-15$ $t^5-8t^4+30t^3-68t^2+105t-120$	3 / ✘ 1 / ✘
	9_{29}^a	$t^3-5t^2+12t-15$ $t^5-8t^4+26t^3-48t^2+59t-56$	3 / ✘ 2 / ✘		9_{30}^a	$-t^3+5t^2-12t+17$ $2t^3-10t^2+25t-32$	3 / ✘ 1 / ✘