

Abstract. Reporting on joint work with Roland van der Veen, I’ll tell you some stories about $\rho_1$, an easy to define, strong, fast to compute, homomorphic, and well-connected knot invariant. $\rho_1$ was first studied by Rozansky and Overbay [Ro1, Ro2, Ro3, Ov], it has far-reaching generalizations, it is dominated by the coloured Jones polynomial, and I wish I understood it. Common misconception. “Dominated” ⇒ “less”.

Invariants.


Homomorphic. Extends to tangles and behaves under tangle operations; especially gluings and doublings:

Why care for “Homomorphic”? Theorem. A knot $K$ is ribbon iff there exists a 2n-component tangle $T$ with skeleton as below such that $\tau(T) = K$ and where $\delta(T) = U$ is the untangle:

Hear more at oefi/AKT.

Formulas. Draw an $n$-crossing knot $K$ as on the right: all crossings face up, and the edges are marked with a running index $k \in \{1, \ldots, 2n+1\}$ and with rotation numbers $\varphi_k$. Let $A$ be the $(2n+1) \times (2n+1)$ matrix constructed by starting with the identity matrix $I$, and adding a $2 \times 2$ block for each crossing:

\[
\begin{bmatrix}
0 & 1 & 0 & 0 & 0 \\
1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0
\end{bmatrix}
\]

Let $G = (g_{ij}) = A^{-1}$. For the trefoil example, it is:

\[
A = \begin{bmatrix}
1 & T & 0 & 0 & 0 \\
0 & 1 & -1 & 0 & 0 \\
0 & 0 & 0 & 1 & -1 \\
0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 1
\end{bmatrix}
\]

[Note. The Alexander polynomial $\Delta$ is given by $\Delta = T^{(w-w)/2} \det(A)$, with $w = \sum \varphi_k$, $w = \sum c$.]

Classical Topologists: This is boring. Yawn.

Formulas, continued. Finally, set

\[
R_1(c) := s(g_{ij} (g_{j+1} - g_{i}) - g_{ij}(g_{j+1} - 1) - 1/2)
\]

\[
\rho_1 := \Delta^2 \left( \sum_c R_1(c) - \sum_k \varphi_k (g_{kk} - 1/2) \right).
\]

In our example $\rho_1 = -T^2 + 2T - 2 + 2T^{-1} - T^{-2}$.

Theorem. $\rho_1$ is a knot invariant. Proof: later.

Classical Topologists: Whiskey Tango Foxtrot?

Cars, Interchanges, and Traffic Counters. Cars always drive forward. When a car crosses over a bridge it goes through with (algebraic) probability $T^s \sim 1$, but falls off with probability $1 - T^s \sim 0$. See cars. Thanks for listening!

References.


Preliminaries

This is Rho1.nb of http://drorbn.net/gro22/ap.

Once[<<KnotTheory`;<<Rot.m];

Loading KnotTheory version of February 2, 2020, 10:53:45.2097.
Read more at http://katlas.org/wiki/KnotTheory.
Loading Rot.m from http://drorbn.net/gro22/ap to compute rotation numbers.

The Program

\[
R_1[s_, i_, j_] := \left( \frac{1 - T + T^2}{T}, \left\{ \begin{array}{cl} \frac{1}{2} (1 + T^2) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right)
\]

\[
\rho[K_\text{a, b, c} := \text{Inverse}[\text{IdentityMatrix}[2 n + 1]]; \quad \text{Cases}[Cs, {s_, i_, j_} \rightarrow \left( A[[i, j]], (i + 1, j + 1) \right)] = \left( \begin{array}{cc} -T^5 & T^2 - 1 \\ 0 & -1 \end{array} \right) \right)
\]

\[
\Delta = \left( -\text{Total}[Cs] - \text{Total}[Cs[[11, 1]]] \right)/2 \text{Det}[A];
\]

\[
G = \text{Inverse}[A];
\]

\[
\rho_1 = \sum_{k=1}^{n} R_1[Cs[k]] - \sum_{k=1}^{n} \phi[k](g_{kk} - 1/2);
\]

\[
\text{Factor}[(\Delta, \Delta^2, \rho_1) / (g_{\alpha, \beta}, G[\alpha, \beta])]
\]

The First Few Knots

Table[K \rightarrow \rho[K], \{K, AllKnots[[3, 6]]\}]

\[
\begin{align*}
\text{Knot}[3, 1] & \rightarrow \left( \frac{1 - T + T^2}{T}, \left\{ \begin{array}{cl} \frac{1}{2} (1 + T^2) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right), \\
\text{Knot}[4, 1] & \rightarrow \left( \frac{1 - 3 T + T^2}{T}, 0 \right), \quad \text{Knot}[5, 1] \rightarrow \left( \frac{1 - T + T^2 - T^3 + T^4}{T^2}, \left\{ \begin{array}{cl} \frac{1}{2} (1 + T^2) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right), \\
\text{Knot}[5, 2] & \rightarrow \left( \frac{2 - 3 T + 2 T^2}{T}, \left\{ \begin{array}{cl} \frac{1}{2} (1 + T^2) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right), \\
\text{Knot}[6, 1] & \rightarrow \left( \frac{(-2 + T) (-1 + 2 T)}{T}, \left\{ \begin{array}{cl} \frac{1}{2} (1 - 4 T + T^2) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right), \\
\text{Knot}[6, 2] & \rightarrow \left( \frac{-1 - 3 T + 3 T^2 - 3 T^3 + T^4}{T^2}, \left\{ \begin{array}{cl} \frac{1}{2} (1 + 4 T^2 - 4 T^3 + 4 T^4 - 4 T^5 + 6 T^6) & \text{if } s = 1, i = j, j + 1 \text{ or } i = j, j + 2 \\ 1 & \text{otherwise} \end{array} \right. \right), \\
\text{Knot}[6, 3] & \rightarrow \left( \frac{1 - 3 T + 5 T^2 - 3 T^3 + T^4}{T^2}, 0 \right)
\end{align*}
\]

\[ p = 1 - T^8 \]
**Theorem.** $g_{\alpha \beta}$ is the reading of a traffic counter at $\beta$, if car traffic is injected at $\alpha$ (if $\alpha = \beta$, the counter is *after* the injection point).

**Example.**

$$\sum_{p=0}^{\infty}(1-T)^p = T^{-1} \quad \sum_{j,k} = G = \begin{pmatrix} 1 & 0 & 0 \\ 0 & T^{-1} & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

**Proof.** Near a crossing $c$ with sign $s$, incoming upper edge $i$ and incoming lower edge $j$, both sides satisfy the $g$-rules:

$$g_{ij} = \delta_{ij} + T^s g_{i+1,j} + g_{j+1,i} \quad \text{and always, } g_{\alpha 2 \alpha+1} = 1: \text{ use common sense and } AG = I (= GA).$$

**Bonus.** Near $c$, both sides satisfy the further $r$-rules:

$$g_{ij} = T^{-i} (g_{i+1,j-1} - \delta_{i,j+1}) \quad g_{ij} = g_{i+1,j+1} = (1-T^s) g_{i,j} = \delta_{i,j+1}.$$
The Most Important Missing Infrastructure Project in Knot Theory

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10:12 AM

An "infrastructure project" is hard (and sometimes non-glorious) work that's done now and pays off later.

An example, and the most important one within knot theory, is the tabulation of knots up to 10 crossings: I think it precedes Rolfsen, yet the result is often called "the Rolfsen Table of Knots", as it is famously printed as an appendix to the famous book by Rolfsen. There is no doubt the production of the Rolfsen table was hard and non-glorious. Yet its impact was and is tremendous. Every new thought in knot theory is tested against the Rolfsen table, and it is hard to find a paper in knot theory that doesn't refer to the Rolfsen table in one way or another.

A second example is the Hoste-Thistlethwaite tabulation of knots with up to 17 crossings. Perhaps more fun to do as the real hard work was delegated to a machine, yet hard it certainly was: a proof is in the fact that nobody so far had tried to replicate their work, not even to a smaller crossing number. Yet again, it is hard to overestimate the value of that project: in many ways the Rolfsen table is "not yet generic", and many phenomena that appear to be rare when looking at the Rolfsen table become the rule when the view is expanded. Likewise, other phenomena only appear for the first time when looking at higher crossing numbers.

But as I like to say, knots are the wrong object to study in knot theory. Let me quote (with some variation) my own (with Dancso) "WKO" paper:

Studying knots on their own is the parallel of studying cakes and pastries as they come out of the bakery - we sure want to make them our own, but the theory of desserts is more about the ingredients and how they are put together than about the end products. In algebraic knot theory this reflects through the fact that knots are not finitely generated in any sense (hence they must be made of some more basic ingredients), and through the fact that there are very few operations defined on knots (connected sums and satellite operations being the main exceptions), and thus most interesting properties of knots are transcendental, or non-algebraic, when viewed from within the algebra of knots and operations on knots (see [AKT-CFA]).

The right objects for study in knot theory are thus the ingredients that make up knots and that permit a richer algebraic structure. These are braids (which are already well-studied and tabulated) and even more so tangles and tangled graphs.

Thus in my mind the most important missing infrastructure project in knot theory is the tabulation of tangles to as high a crossing number as practical. This will enable a great amount of testing and experimentation for which the grounds are now still missing. The existence of such a tabulation will greatly impact the direction of knot theory, as many tangle theories and issues that are now ignored for the lack of scope, will suddenly become alive and relevant. The overall influence of such a tabulation, if done right, will be comparable to the influence of the Rolfsen table.

Aside. What are tangles? Are they embedded in a disk? A ball? Do they have an "up side" and a "down side"? Are the strands oriented? Do we mod out by some symmetries or figure out the action of some symmetries? Shouldn't we also calculate the affect of various tangle operations (strand doubling and deletion, juxtapositions, etc.)? Shouldn't we also enumerate virtual tangles? Tw-tangles? Tangled graphs?

In my mind it would be better to leave these questions to the tabulator. Anything is better than nothing, yet good tabulators would try to tabulate the more general things from which the more special ones can be sieved relatively easily, and would see that their programs already contain all that would be easy to implement within their frameworks. Counting legs is easy and can be left to the end user. Determining symmetries is better done along with the enumeration itself, and so it should:

An even better tabulation should come with a modern front-end - a set of programs for basic manipulations of tangles, and a web-based "tangle atlas" for an even easier access.

Overall this would be a major project, well worthy of your time.

(Source: http://katlas.math.toronto.edu/drorbn/AcademicPensieve/2012-01/)

The interchange of I-95 and I-695, northeast of Baltimore. (more)

From [AKT-CFA]