

Dror Bar-Natan: Classes: 2014-15: Math 1100 Algebra I:

JCF Tricks and Programs

Row and Column Operations

Row operations are performed by left-multiplying N by some properly-positioned 2×2 matrix and at the same time left-multiplying the “tracking matrix” P by the same 2×2 matrix. Column operations are similar, with left replaced by right and P by Q .

```
RowOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[TT.NN]; PP = Simplify[TT.PP];
];
ColOp[i_, j_, mat_] := Module[{TT = II},
  TT[[{i, j}, {i, j}]] = mat;
  NN = Simplify[NN.TT]; QQ = Simplify[QQ.TT];
];
```

Swapping Rows and Columns

```
SwapRows[i_, j_] := RowOp[i, j,  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ];
SwapColumns[i_, j_] := ColOp[i, j,  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ];
SwapBoth[i_, j_] := (SwapRows[i, j]; SwapColumns[i, j];)
```

The “GCD” Trick

If $q = \gcd(a, b) = sa + tb$, the equality $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} q \\ 0 \end{pmatrix}$ allows us to replace pairs of entries in the same column by their greatest common divisor (and a zero!), using invertible row operations. A similar trick works for rows.

? PolynomialExtendedGCD

PolynomialExtendedGCD[poly₁, poly₂, x] gives the extended GCD of poly₁ and poly₂ treated as univariate polynomials in x .
 PolynomialExtendedGCD[poly₁, poly₂, x, Modulus → p] gives the extended GCD over the integers mod prime p . >>

```
GCDTrick[{i_, j_}, k_] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[i, k]], b = NN[[j, k]], x];
  RowOp[i, j,  $\begin{pmatrix} s & t \\ -b/q & a/q \end{pmatrix}$ ]
];
GCDTrick[k_, {i_, j_}] := Module[{a, b, q, s, t},
  {q, {s, t}} = PolynomialExtendedGCD[a = NN[[k, i]], b = NN[[k, j]], x];
  ColOp[i, j,  $\begin{pmatrix} s & -b/q \\ t & a/q \end{pmatrix}$ ]
];
```

Factoring Diagonal Entries

If $1 = \gcd(a, b) = sa + tb$, the equality $\begin{pmatrix} s & a & 1 \\ -t & b & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & a & b \end{pmatrix} \begin{pmatrix} a & -b \\ t & s \end{pmatrix} = \begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$ is an invertible row-column-operations proof of the isomorphism $\frac{R}{\langle a \rangle} \oplus \frac{R}{\langle b \rangle} \simeq \frac{R}{\langle ab \rangle}$.

```
SplitToSum[i_, j_, a_, b_] := Module[
  {q, s, t, T1, T2},
  {q, {s, t}} = PolynomialExtendedGCD[a, b, x];
  If[q == 1,
    RowOp[i, j,  $\begin{pmatrix} s & a & 1 \\ -t & b & 1 \end{pmatrix}$ ]; ColOp[i, j,  $\begin{pmatrix} a & -b \\ t & s \end{pmatrix}$ ];
  ]
];
```

The Jordan Trick

A repeated application of the identity $\begin{pmatrix} p^{k-1} & -1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & p^k \end{pmatrix} \cdot \begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} p^{-1+k} & 0 \\ 1 & p \end{pmatrix}$ will bring a matrix

like $\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & p^4 \end{pmatrix}$ to the “Jordan” form of $\begin{pmatrix} p & 0 & 0 & 0 \\ 1 & p & 0 & 0 \\ 0 & 1 & p & 0 \\ 0 & 0 & 1 & p \end{pmatrix}$, using invertible row and column operations.

```
JordanTrick[i_, j_, p_, s_] := (RowOp[i, j,  $\begin{pmatrix} p^{s-1} & -1 \\ 1 & 0 \end{pmatrix}$ ]; ColOp[i, j,  $\begin{pmatrix} 1 & p \\ 0 & 1 \end{pmatrix}$ ]);
```