

Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 1

Dror Bar-Natan in Reheja, June 2012



~~Abstract~~ The a priori expectation of first year elementary school students who were just introduced to the natural numbers, if they would be ready to verbalize it, must be that soon, perhaps by second grade, they'd master the theory and know all there is to know about those numbers. But they would be wrong, for number theory remains a thriving subject, well-connected to practically anything there is out there in mathematics.

I was a bit more sophisticated when I first heard of knot theory. My first thought was that it was either trivial or intractable, and most definitely, I wasn't going to learn it is interesting. But it is, and I was wrong, for the reader of knot theory is often lead to the most interesting and beautiful structures in topology, geometry, quantum field theory, and algebra.

Today I will talk about just one minor example, mostly having to do with the link to algebra: A straightforward proposal for a group-theoretic invariant of knots fails if one really means groups, but works once generalized to meta-groups (to be defined). We will construct one complicated but elementary meta-group as a meta-bicrossed-product (to be defined), and explain how the resulting invariant is a not-yet-understood yet potentially significant generalization of the Alexander polynomial, while at the same time being a specialization of a somewhat-understood “universal finite type invariant of w-knots” and of an elusive “universal finite type invariant of v-knots”.

**Closely related** to work by Le Dimet (Comment. Math. Helv. **67** (1992) 306–315), Kirk, Livingston and Wang (arXiv:math/9806035) and Cimasoni and Turaev (arXiv:math.GT/0406269).

ol Alexander Issues.

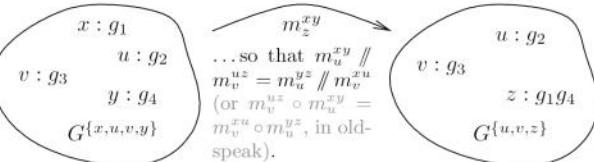
- Quick to compute, but computation departs from topology.
  - Extends to tangles, but at an exponential cost.
  - Hard to categorify.

**Idea.** Given a group  $G$  and two “YB” pairs  $R^\pm = (g_o^\pm, g_u^\pm) \in G^2$ , map them to xings and “multiply along”, so that

$$\begin{array}{c} \diagup^{\pm} \diagdown \rightarrow \diagup g_o^{\pm} \\ \diagdown \quad \diagup \quad \diagdown q_u^{\pm} \\ g_u^+ g_o^+ g_u^+ \end{array} \Big)$$

**This Fails!** R2 implies that  $g_o^\pm g_o^\mp = e = g_u^\pm g_u^\mp$  and then R3 implies that  $g_o^+$  and  $g_u^+$  commute, so the result is a simple counting invariant.

**A Group Computer.** Given  $G$ , can store group elements and perform operations on them:



Also has  $S_x$  for inversion,  $e_x$  for unit insertion,  $d_x$  for register deletion,  $\Delta_{xy}^z$  for element cloning,  $\rho_y^x$  for renamings, and  $(D_1, D_2) \mapsto D_1 \cup D_2$  for merging, and many obvious composition axioms relating those.

$$P = \{x : q_1, y : q_2\} \Rightarrow P = \{d_y P\} \cup \{d_x P\}$$

$$P = \{x : g_1, y : g_2\} \Rightarrow P = \{d_y P\} \cup \{d_x P\}$$

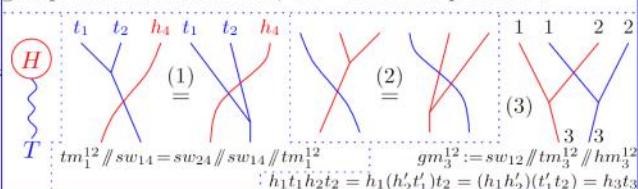
**A Meta-Group.** Is a similar “computer”, only its internal structure is unknown to us. Namely it is a collection of sets  $\{G_\gamma\}$  indexed by all finite sets  $\gamma$ , and a collection of operations  $m_z^{xy}, S_x, e_x, d_x, \Delta_{xy}^z$  (sometimes),  $\rho_y^x$ , and  $\cup$ , satisfying the exact same *linear* properties.

**Example 1.** The non-meta example,  $G_\gamma := G^\gamma$ .

**Example 2.**  $G_\gamma := M_{\gamma \times \gamma}(\mathbb{Z})$ , with simultaneous row and column operations, and “block diagonal” merges. Here if  $P = \begin{pmatrix} x & a & b \\ y & c & d \end{pmatrix}$  then  $d_y P = (x : a)$  and  $d_x P = (y : d)$  so  $\{d_x P\} \cup \{d_y P\} = \{x : a \ 0\} \neq P$ . So this  $G$  is truly meta.

**Claim.** From a meta-group  $G$  and YB elements  $R^\pm \in G_2$  we can construct a knot/tangle invariant.

**Bicrossed Products.** If  $G = HT$  is a group presented as a product of two of its subgroups, with  $H \cap T = \{e\}$ , then also  $G = TH$  and  $G$  is determined by  $H, T$ , and the “swap” map  $sw^{th} : (t, h) \mapsto (h', t')$  defined by  $th = h't'$ . The map  $sw$  satisfies (1) and (2) below; conversely, if  $sw : T \times H \rightarrow H \times T$  satisfies (1) and (2) (+ lesser conditions), then (3) defines a group structure on  $H \times T$ , the “bicrossed product”.



**A Standard Alexander Formula.** Label the arcs 1 through  $(n+1) = 1$ , make an  $n \times n$  matrix as below, delete one row and one column, and compute the determinant:

$$\begin{array}{c} \text{Diagram 1: } \begin{array}{l} \text{Top row: } b \text{ (up), } c \text{ (down-left), } a \text{ (down-right)} \\ \text{Bottom row: } + \text{ (up-right), } c \text{ (down)} \end{array} \rightarrow \begin{array}{c} \text{Top row: } a \text{ (up), } b \text{ (down-left), } c \text{ (down-right)} \\ \text{Bottom row: } c \text{ (up-right), } -1 \text{ (down)} \end{array} \\[10pt] \text{Diagram 2: } \begin{array}{l} \text{Top row: } b \text{ (up), } a \text{ (down-left), } c \text{ (down-right)} \\ \text{Bottom row: } a \text{ (up-right), } c \text{ (down)} \end{array} \rightarrow \begin{array}{c} \text{Top row: } a \text{ (up), } b \text{ (down-left), } c \text{ (down-right)} \\ \text{Bottom row: } c \text{ (up-right), } -X \text{ (down)} \end{array} \end{array}$$

```

1   0   0   0   0   X-1   0   -X
-1   X   0   0   0   0   1-X   0
0   -1   X   0   1-X   0   0   0
X-1   0   -X   1   0   0   0   0
0   1-X   0   -1   X   0   0   0
0   0   0   0   -X   1   0   X-1
0   0   1-X   0   0   -1   X   0
0   0   0   X-1   0   0   -X   1

```

$\left[ [1 :: 7, 1 :: 7] \right] // \text{Det}$

## Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 2

A **Meta-Bicrossed-Product** is a collection of sets  $\beta(\eta, \tau)$  and operations  $tm_z^{xy}$ ,  $hm_z^{xy}$  and  $sw_{xy}^{th}$  (and lesser ones), such that  $tm$  and  $hm$  are “associative” and (1) and (2) hold (+ lesser conditions). A meta-bicrossed-product defines a meta-group with  $G_\gamma := \beta(\gamma, \gamma)$  and  $gm$  as in (3).

**Example.** Take  $\beta(\eta, \tau) = M_{\tau \times \eta}(\mathbb{Z})$  with row operations for the tails, column operations for the heads, and a trivial swap.

**$\beta$  Calculus.** Let  $\beta(\eta, \tau)$  be

$$\left\{ \begin{array}{c|ccc} \omega & h_1 & h_2 & \dots \\ \hline t_1 & \alpha_{11} & \alpha_{12} & \dots \\ t_2 & \alpha_{21} & \alpha_{22} & \dots \\ \vdots & \ddots & \ddots & \ddots \end{array} \right| \begin{array}{l} h_j \in \eta, t_i \in \tau, \text{ and } \omega \text{ and} \\ \text{the } \alpha_{ij} \text{ are rational functions in a variable } X \end{array},$$

$$tm_z^{xy}: \frac{\omega}{t_x} \frac{\dots}{t_y} \mapsto \frac{\omega}{t_z} \frac{\dots}{\gamma}, \quad \frac{\omega_1}{t_1} \frac{\eta_1}{\alpha_1} \cup \frac{\omega_2}{t_2} \frac{\eta_2}{\alpha_2} = \frac{\omega_1 \omega_2}{t_1} \frac{\eta_1 \eta_2}{\alpha_1 \alpha_2},$$

$$hm_z^{xy}: \frac{\omega}{t_x} \frac{h_x}{\alpha} \frac{h_y}{\beta} \frac{\dots}{\gamma} \mapsto \frac{\omega}{t_x} \frac{h_z}{\alpha + \beta + \langle \alpha \rangle \beta} \frac{\dots}{\gamma},$$

$$sw_{xy}^{th}: \frac{\omega}{t_x} \frac{h_y}{\alpha} \frac{\dots}{\gamma} \mapsto \frac{\omega \epsilon}{t_x} \frac{h_y}{\alpha(1 + \langle \gamma \rangle / \epsilon)} \frac{\dots}{\gamma / \epsilon},$$

where  $\epsilon := 1 + \alpha$  and  $\langle c \rangle := \sum_i c_i$ , and let

$$R_{xy}^p := \frac{1}{t_x} \frac{h_x}{0} \frac{h_y}{X-1} \quad R_{xy}^m := \frac{1}{t_x} \frac{h_x}{0} \frac{h_y}{X^{-1}-1}.$$

**Theorem.**  $Z^\beta$  is a tangle invariant (and more). Restricted to knots, the  $\omega$  part is the Alexander polynomial. On braids, it is equivalent to the Burau representation. A variant for links contains the multivariable Alexander polynomial.

**Why Happy?** • Applications to w-knots.

- Everything that I know about the Alexander polynomial can be expressed cleanly in this language (even if without proof), except HF, but including genus, ribbonness, cabling, v-knots, knotted graphs, etc., and there's potential for vast generalizations.
- The least wasteful “Alexander for tangles” I'm aware of.
- Every step along the computation is the invariant of something.
- Fits on one sheet, including implementation & propagation.



I mean business!

```

$Simp = Factor; SetAttributes[$Collect, Listable];
$Collect[B[a_, b_]] := B[$Simp[a], 
  Collect[a, b, $Simp[b]]];
Form[B[w_, d_]] := Module[{ts, hs, M},
  ts = Union[Cases[B[w_, d_], t_, _?ModuleQ]];
  hs = Union[Cases[B[w_, d_], h_, _?ModuleQ]];
  ts = Outer[$Simp[Coefficient[a, h1] tag1] &, hs, ts];
  PrependTo[M, ts & /@ ts];
  M = Prepend[Transpose[M], Prepend[bw & /@ hs, w]];
  MatrixForm[M]];
Form[_else_] := else /.  $\beta_B \mapsto \beta$ Form[\beta];
Format[\beta_B, StandardForm] :=  $\beta$ Form[\beta];

```

$\{\beta = B[\omega, \text{Sum}[\alpha_{10 i+j} t_i h_j, \{i, \{1, 2, 3\}\}, \{j, \{4, 5\}\}]],$   
 $(\beta // tm_{12 \rightarrow 1} // sw_{14}) = (\beta // sw_{24} // sw_{14} // tm_{12 \rightarrow 1})$

$\left\{ \begin{array}{c|ccc} \omega & h_4 & h_5 \\ \hline t_1 & \alpha_{14} & \alpha_{15} \\ t_2 & \alpha_{24} & \alpha_{25} \\ t_3 & \alpha_{34} & \alpha_{35} \end{array} \right|, \text{True} \right\} \stackrel{(1)}{=} \text{Some testing}$

$\{Rm_{51} Rm_{62} Rp_{34} // gm_{14 \rightarrow 1} // gm_{25 \rightarrow 2} // gm_{36 \rightarrow 3},$   
 $Rp_{61} Rm_{24} Rm_{35} // gm_{14 \rightarrow 1} // gm_{25 \rightarrow 2} // gm_{36 \rightarrow 3}\}$

$\left\{ \begin{array}{c|cc} 1 & h_1 & h_2 \\ \hline t_2 & -\frac{1+x}{x} & 0 \\ t_3 & \frac{-1-x}{x} & -\frac{1-x}{x} \end{array} \right\}, \left\{ \begin{array}{c|cc} 1 & h_1 & h_2 \\ \hline t_2 & -\frac{1+x}{x} & 0 \\ t_3 & \frac{-1-x}{x} & -\frac{1+x}{x} \end{array} \right\}$

... divide and conquer!

$\beta = Rm_{12,1} Rm_{27} Rm_{83} Rm_{4,11} Rp_{16,5} Rp_{6,13} Rp_{14,9} Rp_{10,15}$  817

1	h <sub>1</sub>	h <sub>3</sub>	h <sub>5</sub>	h <sub>7</sub>	h <sub>9</sub>	h <sub>11</sub>	h <sub>13</sub>	h <sub>15</sub>
t <sub>2</sub>	0	0	0	- $\frac{1-x}{x}$	0	0	0	0
t <sub>4</sub>	0	0	0	0	0	- $\frac{1+x}{x}$	0	0
t <sub>6</sub>	0	0	0	0	0	0	-1+x	0
t <sub>8</sub>	0	- $\frac{1-x}{x}$	0	0	0	0	0	0
t <sub>10</sub>	0	0	0	0	0	0	0	-1+x
t <sub>12</sub>	- $\frac{1+x}{x}$	0	0	0	0	0	0	0
t <sub>14</sub>	0	0	0	0	-1+x	0	0	0
t <sub>16</sub>	0	0	-1+x	0	0	0	0	0

Do[\beta = \beta // gm\_{1k \rightarrow 1}, \{k, 2, 10\}]; \beta 817, cont.

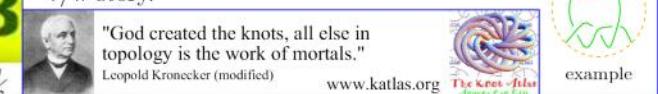
$\frac{1}{x}$	$h_1$	$h_{11}$	$h_{13}$	$h_{15}$
$t_{12}$	$-\frac{(-1-x)(1-x)}{x}$	$-(1+x)(1-X+X^2)$	$(-1+x)(1-X+X^2)$	$-1+x$
$t_{14}$	$-1+x$	$\frac{(-1-x)^2(1-x-x^2)}{x}$	$\frac{(-1-x)^2(1-x-x^2)}{x}$	0
$t_{16}$	$\frac{-1-x}{x}$	$(-1+x)^2$	$\frac{(-1-x)^3}{x}$	0

Do[\beta = \beta // gm\_{1k \rightarrow 1}, \{k, 11, 16\}]; \beta

$$\left( -\frac{1-4x+8x^2-11x^3+8x^4-4x^5+x^6}{x^3} \right)$$

**A Partial To Do List.** 1. Where does it *more simply* come from?

2. Remove all the denominators.
3. How do determinants arise in this context?
4. Understand links.
5. Find the “reality condition”.
6. Do some “Algebraic Knot Theory”.
7. Categorify.
8. Do the same in other natural quotients of the v/w-story.



Further examples of meta-structures.

Meta-monoids:  $\text{TT}_1$ ,  $\mathbb{A}$ ,  $\sqrt{\text{T}}$

(& quotients)

(& quotients)

Meta-bicrossed-products:  $\text{TT}_1 \times \mathbb{A}^r$ ,  $M_0$ ,  $M_1$ ,  $K^{\leq h}$ ,  
 $K^{rbh}$

$\rightarrow$   
and variants  
 $\downarrow$

Meta-Lie-objects.  $\mathbb{A}$  (& many quotients),  $S$

Meta-Lie-Sialgebras:  $\vec{\mathbb{A}}$  (& quotients)