

Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 1

Dror Bar-Natan at Knots in Washington XXXIV

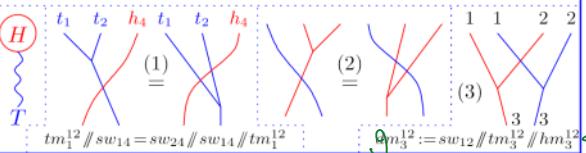
<http://www.math.toronto.edu/~drorbn/Talks/GWU-1203/>

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Abstract. A straightforward proposal for a group-theoretic invariant of knots fails if one really means groups, but works once generalized to meta-groups (to be defined). We will construct one complicated but elementary meta-group as a meta-bicrossed-product (to be defined), and explain how the resulting invariant is a not-yet-understood generalization of the Alexander polynomial, while at the same time being a specialization of a somewhat-understood “universal finite type invariant of w-knots” and of an elusive “universal finite type invariant of v-knots”.

Bicrossed Products. If $G = HT$ is a group presented as a product of two of its subgroups, with $H \cap T = \{e\}$, then also $G = TH$ and G is determined by H, T , and the “swap” map $sw^{th} : (t, h) \mapsto (h', t')$ defined by $th = h't'$. The map sw satisfies (1) and (2) below; conversely, if $sw : T \times H \rightarrow H \times T$ satisfies (1) and (2) (+ lesser conditions), then (3) defines a group structure on $H \times T$, the “bicrossed product”.



A **Meta-Bicrossed-Product** is a collection of sets $\beta(H, T)$ and operations tm_z^{xy}, hm_z^{xy} and sw_{xy}^{th} (and lesser ones), such that tm and hm are “associative” and (1) and (2) hold (+ lesser conditions). A meta-bicrossed-product defines a meta-group with $G_X := \beta(X, X)$ and dm as in (3).

β Calculus. Let $\beta(H, T)$ be

$$\left\{ \begin{array}{c|ccc} \omega & h_1 & h_2 & \dots \\ \hline t_1 & \alpha_{11} & \alpha_{12} & \cdot \\ t_2 & \alpha_{21} & \alpha_{22} & \cdot \\ \vdots & \cdot & \cdot & \cdot \end{array} \right. \quad h_j \in H, t_i \in T, \text{ and } \omega \text{ and the } \alpha_{ij} \text{ are Laurent polynomials in variables } T_i, \text{ in a bijection with the } t_i \text{'s} \right\},$$

$$\text{with operations } tm_z^{xy} : \left(\begin{array}{c|cc} \omega & \dots \\ \hline t_x & \alpha \\ t_y & \beta \\ \vdots & \gamma \end{array} \right) \mapsto \left(\begin{array}{c|cc} \omega & \dots \\ \hline t_z & \alpha + \beta \\ \vdots & \gamma \end{array} \right),$$

$$hm_z^{xy} : \left(\begin{array}{c|cc} \omega & h_x & h_y & \dots \\ \hline \alpha & \beta & \gamma & \cdot \end{array} \right) \mapsto \left(\begin{array}{c|cc} \omega & h_z & \dots \\ \hline \alpha + \beta + \langle \alpha \rangle \beta & \gamma & \cdot \end{array} \right),$$

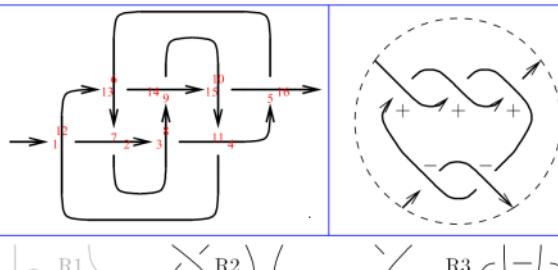
$$sw_{xy}^{th} : \left(\begin{array}{c|cc} \omega & h_y & \dots \\ \hline t_x & \alpha & \beta \\ \vdots & \gamma & \delta \end{array} \right) \mapsto \left(\begin{array}{c|cc} \omega \epsilon & h_y & \dots \\ \hline t_x & \alpha(1 + \langle \gamma \rangle / \epsilon) & \beta(1 + \langle \gamma \rangle / \epsilon) \\ \vdots & \gamma / \epsilon & \delta - \gamma \beta / \epsilon \end{array} \right),$$

where $\epsilon := 1 + \alpha$, $\langle \alpha \rangle := \sum_i \alpha_i$, and $\langle \gamma \rangle := \sum_{i \neq x} \gamma_i$, and let

$$R_{xy}^p := \frac{1}{t_x} \left| \begin{array}{cc} h_x & h_y \\ 0 & T_x - 1 \\ t_y & 0 \end{array} \right| \quad R_{xy}^m := \frac{1}{t_x} \left| \begin{array}{cc} h_x & h_y \\ 0 & T_x^{-1} - 1 \\ t_y & 0 \end{array} \right|.$$

Theorem. Z^β is a tangle invariant (and much more). Restricted to knots, the ω part is the Alexander polynomial. Restricted to links, it contains the multivariable Alexander polynomial. Restricted to braids, it is equivalent to the Burau representation.

Why happy? Besides applications to w-knots, everything that I know about the Alexander polynomial can be expressed clearly in this language (even if w/o proof), except HF but including genus, ribbonness, cabling, V-knots, knotted graphs, etc., and there's potential for vast generalizations.

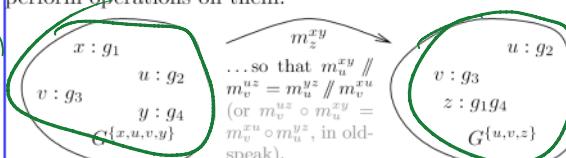


Idea. Given a group G and two pairs $R^\pm = (g_o^\pm, g_u^\pm) \in G^2$, map them to xings and “multiply along”, so that

$$\left(\begin{array}{c} \text{trefoil knot} \\ \text{with crossings} \\ \text{labeled 1-16} \end{array} \right) \xrightarrow{Z} \left(\begin{array}{c} \text{two pairs of xings} \\ \text{labeled } g_o^\pm, g_u^\pm \end{array} \right)$$

This Fails! R2 implies that $g_o^\pm g_u^\mp = e$ and then R3 implies that g_o^+ and g_u^+ commute, to teh result is a simple counting invariant.

A Group Computer. Given G , can store group elements and perform operations on them:



(Also has S_x for inversion, e_x for unit insertion, d_x for register deletion, Δ_{xy}^z for element cloning, and $(D_1, D_2) \mapsto D_1 \cup D_2$ for merging, and very many obvious composition axioms relating these.)

A Meta-Group. Is a similar “computer”, only its internal structure is unknown to us. Namely it is a collection of sets $\{G_X\}$ indexed by all finite sets X , and a collection of operations $m_z^{xy}, S_x, e_x, d_x, \Delta_{xy}^z$, and \cup , satisfying the exact same properties.

Example 1. The non-meta example, $G_X := G^X$.

Example 2. $G_X := M_{X \times X}(\mathbb{Z})$, with simultaneous row and column operations.

global dm → gm

global dm → gm

Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 2

Consider applying with Utilities.m

```
BSimp = Factor; SetAttributes[βCollect, Listable];
βCollect[B[ω_, A_]] := B[βSimp[ω],
  Collect[A, h_, Collect[#, t_, βSimp] &]];
βForm[B[ω_, A_]] := Module[{ts, hs, M},
  ts = Union[Cases[B[ω, A], {t | T}_s_ ↪ s, Infinity]];
  hs = Union[Cases[B[ω, A], h_s_ ↪ s, Infinity]];
  M = Outer[βSimp[Coefficient[A, h_{s1} t_{s2}]] &, hs, ts];
  PrependTo[M, t_ & /@ ts];
  M = Prepend[Transpose[M], Prepend[h_s & /@ hs, ω]];
  MatrixForm[M]];
βForm[else_] := else /. β_B ↪ βForm[β];
Format[β_B, StandardForm] := βForm[β];
```

The key implementation trick is the bijection

$$\frac{\omega + h_j}{t_i} \underset{\alpha_{ij}}{\longleftrightarrow} B(\omega, \sum_{i,j} \alpha_{ij} t_i h_j) :$$

```
μ_ := μ /. t_ → 1;
tmx_y→z_[β_] := β /. {tx|y → tx, Tx|y → Tz};
hmx_y→z_[B[ω_, A_]] := Module[
  {α = D[A, hx], β = D[A, hy], γ = A /. hx|y → 0},
  B[ω, (α + (1 + α) β) hx + γ] // βCollect];
swx_y→z_[B[ω_, A_]] := Module[{α, β, γ, δ, ε},
  α = Coefficient[A, hy tx]; β = D[A, tx] /. hy → 0;
  γ = D[A, hy] /. tx → 0; δ = A /. hy | tx → 0;
  ε = 1 + α;
  B[ω * ε, α (1 + (γ / ε) hy tx + β (1 + (γ / ε) tx
    + γ / ε hy) + δ - γ * β / ε
  ] // βCollect];
dmx_y→z_[β_] := β // swx_y // hm_x,y→z // tmx_y→z;
B /: B[a1_, A1_] B[a2_, A2_] := B[a1 * a2, A1 + A2];
Rp_x,y_ := B[1, (Tx - 1) tx hy];
Rm_x,y_ := B[1, (Tx^-1 - 1) tx hy];
```

```
{β = B[ω, Sum[a10 i+j t1 hj, {i, {1, 2, 3}}, {j, {4, 5}}]],
β // tm1,2→1 // sw1,4,
β // sw2,4 // sw1,4 // tm2,2→1
} // ColumnForm
```

Some testing...

Put a copy of (1) here.

$\begin{pmatrix} \omega & h_4 & h_5 \\ t_1 & \alpha_{14} & \alpha_{15} \\ t_2 & \alpha_{24} & \alpha_{25} \\ t_3 & \alpha_{34} & \alpha_{35} \end{pmatrix}$	$\begin{pmatrix} h_4 \\ h_5 \end{pmatrix}$
$\begin{pmatrix} \omega (1 + \alpha_{14} + \alpha_{24}) \\ t_1 \\ t_2 \\ t_3 \end{pmatrix}$	$\begin{pmatrix} \frac{(\alpha_{14}+\alpha_{24})(1-\alpha_{14}-\alpha_{24}+\alpha_{34})}{1-\alpha_{14}-\alpha_{24}} & \frac{(\alpha_{15}+\alpha_{25})(1-\alpha_{14}-\alpha_{24}+\alpha_{34})}{1-\alpha_{14}-\alpha_{24}} \\ \frac{\alpha_{34}}{1-\alpha_{14}-\alpha_{24}} & -\alpha_{15}\alpha_{34}-\alpha_{25}\alpha_{34}+\alpha_{35}+\alpha_{14}\alpha_{35}+\alpha_{24}\alpha_{35} \\ \frac{1}{1-\alpha_{14}-\alpha_{24}} & 1-\alpha_{14}-\alpha_{24} \end{pmatrix}$
$\begin{pmatrix} \omega (1 + \alpha_{14} + \alpha_{24}) \\ t_1 \\ t_2 \\ t_3 \end{pmatrix}$	$\begin{pmatrix} h_4 \\ h_5 \end{pmatrix}$

$\{Rm_{5,1} Rm_{6,2} Rp_{3,4} // dm_{1,4→1} // dm_{2,5→2} // dm_{3,6→3},$
 $Rp_{6,1} Rm_{2,4} Rm_{3,5} // dm_{1,4→1} // dm_{2,5→2} // dm_{3,6→3}\}$

$\left\{ \begin{pmatrix} 1 & h_1 & h_2 \\ t_2 & -\frac{1-T_2}{T_2} & 0 \\ t_3 & \frac{-1-T_3}{T_2} & -\frac{1-T_3}{T_3} \end{pmatrix}, \begin{pmatrix} 1 & h_1 & h_2 \\ t_2 & -\frac{1-T_2}{T_2} & 0 \\ t_3 & \frac{-1-T_3}{T_2} & -\frac{1-T_3}{T_3} \end{pmatrix} \right\}$

$\beta = Rm_{12,1} Rm_{2,7} Rm_{8,3} Rm_{4,11} Rp_{16,5} Rp_{6,13} Rp_{14,9} Rp_{10,15}$ [End]

1	h ₁	h ₃	h ₅	h ₇	h ₉	h ₁₁	h ₁₃	h ₁₅
t ₂	0	0	0	$-\frac{1-T_2}{T_2}$	0	0	0	0
t ₄	0	0	0	0	0	$-\frac{1-T_4}{T_4}$	0	0
t ₆	0	0	0	0	0	0	$-1+T_6$	0
t ₈	0	$-\frac{1-T_8}{T_8}$	0	0	0	0	0	0
t ₁₀	0	0	0	0	0	0	0	$-1+T_{10}$
t ₁₂	$-\frac{1-T_{12}}{T_{12}}$	0	0	0	0	0	0	0
t ₁₄	0	0	0	0	$-1+T_{14}$	0	0	0
t ₁₆	0	0	$-1+T_{16}$	0	0	0	0	0

$\frac{T_2-T_{14}-T_1-T_{16}}{T_1^2}$	h ₁	h ₁₁	h ₁₃	h ₁₅
t ₁	$-\frac{(-1-T_1)(T_{14}-(T_1^2-T_2^2))}{T_1^2 T_{12}(T_1^2-T_{14}-T_1-T_{16})}$	$-\frac{(-1-T_1)(1-T_1-T_1^2)T_{14}T_{16}}{T_1(T_1^2-T_{14}-T_1-T_{16})}$	$-\frac{(-1-T_1)(1-T_1-T_1^2)T_{14}}{T_1^2-T_{14}-T_1-T_{16}}$	$-1+T_1$
t ₁₂	$-\frac{1-T_{12}}{T_{12}}$	0	0	0
t ₁₄	$-\frac{(-1-T_{14})(-T_1-T_1^2-T_{16})}{T_{12}(T_1^2-T_{16}-T_1-T_{14})}$	$-\frac{(-1-T_1)(1-T_1-T_1^2)(-1-T_{14})T_{16}}{T_1(T_1^2-T_{16}-T_1-T_{14})}$	$-\frac{(-1-T_1)(1-T_1-T_1^2)(-1-T_{14})}{T_1^2-T_{16}-T_1-T_{14}}$	0
t ₁₆	$\frac{T_1(-T_2)T_6}{T_{12}(T_1^2-T_{16}-T_1-T_{14})}$	$-\frac{(-1-T_1)T_7(-T_1-T_1^2)T_{16}}{T_1^2-T_{16}-T_1-T_{14}}$	$-\frac{(-1-T_1)T_8(-T_1-T_1^2)T_{16}}{T_1^2-T_{16}-T_1-T_{14}}$	0

Do[β = β // dm _{1,k+1} , {k, 11, 16}]; β	$\frac{1-4 T_1+8 T_1^2-11 T_1^3+8 T_1^4-4 T_1^5+T_1^6}{T_1^3}$
t ₁	0

<< KnotTheory`
Alexander[Knot[8, 17]][T1] // Factor

Loading KnotTheory` version of August 22, 2010, 13:36:57.55.
Read more at <http://katlas.org/wiki/KnotTheory>.

KnotTheory`loading: Loading precomputed data in PD4Knots`
 $\frac{1-4 T_1+8 T_1^2-11 T_1^3+8 T_1^4-4 T_1^5+T_1^6}{T_1^3}$

Where does it come from?
My to do list.

Some testing...

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Footnotes

1. Test.

References

- [BND] D. Bar-Natan and Zsuzsanna Dancso, *Finite Type Invariants of w-Knotted Objects: From Alexander to Kashiwara and Vergne*, in preparation, online at <http://www.math.toronto.edu/~drorbn/papers/WKO/>.

Plan

1. (? minutes) ??