

Type Theory

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6:25 AM

G grothendieck universe

$j(G)$ is also grothendieck universe
(not really, but close enough)

$G \in j(G)$

$j(G)$ is like κ $a \in \kappa \Rightarrow j(a) = a$

obvious property: φ is true in G iff
 φ is true in $j(G)$

$f: G \rightarrow G$

$j(f): j(G) \rightarrow j(G)$

$A \in G$

$j(A) \subseteq j(G)$

For all x , there is $j(x)$

for example there is $j(j(a))$

$\forall a, b, c, \dots, d \quad \varphi(a, b, c, \dots, d)$ iff

$\models \varphi(j(a), j(b), \dots, j(d))$

M is a world where σ isn't measurable

$A \subseteq G \Rightarrow j(A) \subseteq j(G)$ is $G \in j(A)$

$V = \{ A \subseteq G \mid G \in j(A) \}$

Say, you want to prove: b is so big, it has
 Q in it. Example, another grothendieck
universe. Then you look in M , and
a Q out of a inside $j(b)$,
then, by (*), there is a Q in a .

What is a natural number?

With an object $a : X$, and function
 $f : X \rightarrow X$, it is always possible to
define $f^n(x)$ for a natural number
 n .

Converse; if it is possible to
define $n(f, x) \in X$ for $x \in X$ and
 $f : X \rightarrow X$, then n corresponds to a natural
number and for any set X

Def.

A natural number is something which
turns a set X and a function $f : X \rightarrow X$
into another function $f^n : X \rightarrow X$

This is what they do in system F
System F:

Types: There type variables
 X, V, T, \dots

if $A \& B$ are types, $A \otimes B$
is a type

If A is a type and X is type
variable, $\Pi X A$ is a type
eg. $\Pi X. (X \rightarrow X) \rightarrow (X \rightarrow X)$ is the type
of all natural numbers, as above